

# CSE211: Compiler Design

Oct. 6, 2021

- **Topic:** Finish PLY overview, go over symbol tables.
- **Questions:**
  - *Has anyone started on the homework? Any issues?*



from: <https://en.wikipedia.org/wiki/Yak>

# Announcements

- Homework 1 is out
  - Due on the 18<sup>th</sup>
  - Get started early!
- Office hours tomorrow (2-3pm, E2-233)
- if you have ideas for projects, we can start discussing!
- Keep an eye out for homework questions/clarifications on slack

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from: <https://en.wikipedia.org/wiki/Yak>

# Review: Parser generators

- Specify:
  - Tokens
  - Production Rules
  - Production Actions
- Parser generator gives you a function in which you can pass strings
  - Executes production actions
  - Error reporting

# Review: PLY:

- How did we specify tokens?
- What are token actions?
- How did we specify production rules?
  - Are you allowed to in your homework?
- How did we specify precedence and associativity?

# Review: PLY:

- Catch-up on the calculator example

# Simplifying binary operations with Lambdas

```
def p_expr_bin(p):  
    """  
    expr : expr PLUS expr  
          | expr MINUS expr  
          | expr MULT expr  
    """  
    if p[2] == '+':  
        p[0] = p[1] + p[3]  
    elif p[2] == '-':  
        p[0] = p[1] - p[3]  
    elif p[2] == '*':  
        p[0] = p[1] * p[3]  
    else:  
        assert(False)
```

Can be changed to (next slide)

# Simplifying binary operations with Lambdas

```
def p_plus(p):  
    "plusp : PLUS"  
    p[0] = lambda x,y: x+y
```

```
def p_multp(p):  
    "multp : MULT"  
    p[0] = lambda x,y: x*y
```

```
def p_minusp(p):  
    "minusp : MINUS"  
    p[0] = lambda x,y: x-y
```

```
def p_expr_bin(p):  
    ""  
    expr : expr plusp expr  
          | expr minusp expr  
          | expr multp expr  
    ""  
    p[0] = p[2](p[1], p[3])
```

Can be changed to (next slide)



# Multiline calculator example

- *A sequence of expressions?*

```
to_print = []
```

```
def p_expression_list(p):  
    "expr_list : expr SEMI"  
    to_print.append(p[1])
```

```
def p_expression_list_rec(p):  
    "expr_list : expr_list expr SEMI"  
    to_print.append(p[2])
```

*Is this order important?*

# Multiline calculator example

- *A better error function?*

```
def p_error(p):  
    print("Syntax error in input on line: %d" % p.lineno)  
    exit(1)
```

What are other options? try to recover?

# Multiline calculator example

- *Attempting to recover:*

```
def p_error(p):
    print("Syntax error in input on line: %d" % p.lineno)
    print("trying to recover")
    while True:
        tok = parser.token()
        if tok.type == 'SEMI': break
    print("trying restart after the ; on line %d" % p.lineno)
    to_print.append("ERROR")
    parser.restart()
```

# How to handle keywords and ids

- How to differentiate keywords from ids:
  - e.g. “if”, from “x”
  - token for id is “[a-zA-Z]+”
  - it will also match keywords...

# How to handle keywords and ids

```
tokens = ["IF", "ELSE", "ID"]
```

```
t_ID = "[a-zA-Z]+"
```

```
t_IF = "if"
```

```
t_ELSE = "else"
```

```
t_ignore = ' '
```

```
def t_error(t):
```

```
    print("Illegal character '%s'" % t.value[0])
```

```
    print("line number: %d" % t.lexer.lineno)
```

```
    exit(1)
```

```
lexer = lex.lex()
```

```
lexer.input("if")
```

parses "if" as an ID!

# How to handle keywords and ids

```
reserved = {  
    'if'      : 'IF',  
    'else'    : 'ELSE'  
}  
  
tokens = ["ID"] + list(reserved.values())  
  
def t_ID(t):  
    "[a-zA-Z]+"  
    t.type = reserved.get(t.value, 'ID')  
    return t
```

This will work!

# Conclusion: lots of interesting features

- Modern parser generators are really great!
- I highly suggest reading the PLY readme
  - Even more examples and interesting functionality
- PLY was largely developed for educational purposes, but it's been reliable for me for several projects, especially other parts of your project are in Python.
- While I have never used it, Antlr is highly recommended. If anyone is interested in doing any of homework in Antlr let me know!

# Back to presentation mode

- To discuss symbol tables!



# One consideration: Scope

- What is scope?
- Can it be determined at compile time? Can it be determined at runtime?
- C vs. Python
- Anyone have any interesting scoping rules they know of?

# One consideration: Scope

- Lexical scope example

```
int x = 0;
int y = 0;
{
    int y = 0;
    x+=1;
    y+=1;
}
x+=1;
y+=1;
```

What are the final values in x and y?

# How to track scope?

- Symbol table
- Global object, accessible (and mutable) by all production actions
- two methods:
  - **lookup(id)** : lookup an id in the symbol table.  
Returns None if the id is not in the symbol table.
  - **insert(id,info)** : insert a new id (or overwrite an existing id) into the symbol table along with a set of information about the id.

What information might we store about an id?

# a very simple programming language

VARIABLE\_NAME = [a-z]+

INCREMENT = "\+\+"

TYPE = "int"

LB = "{"

RB = "}"

SEMI = ";"

```
int x;  
x++;  
int y;  
y++;
```

statements are either a declaration or an increment

# a very simple programming language

VARIABLE\_NAME = [a-z]+

INCREMENT = "\+\+"

TYPE = "int"

LB = "{"

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```
int x;  
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    x++;  
    y++;  
}  
y++;
```

statements are either a declaration or an increment

# a very simple programming language

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```
int x;  
{  
    int y;  
    x++;  
    y++;  
}  
y++;
```

statements are either a declaration or an increment

# How to track scope?

- `SymbolTable ST;`

```
declare_variable: TYPE VARIABLE_NAME SEMI  
{ }
```

Say we are matched string:  
`int x;`

**lookup(id)** : lookup an id in the symbol table. Returns None if the id is not in the symbol table.

**insert(id,info)** : insert a new id (or overwrite an existing id) into the symbol table along with a set of information about the id.

# How to track scope?

- `SymbolTable ST;`

```
declare_variable: TYPE VARIABLE_NAME SEMI  
{ST.insert(C[1],C[0])}
```

Say we are matched string:  
`int x;`

In this example we are storing a type



# How to track scope?

- `SymbolTable ST;`

Say we are matched string:  
`x++;`

```
variable_inc: VARIABLE_NAME INCREMENT SEMI  
{ }
```

**lookup(id)** : lookup an id in the symbol table. Returns None if the id is not in the symbol table.

**insert(id,info)** : insert a new id (or overwrite an existing id) into the symbol table along with a set of information about the id.

# How to track scope?

- `SymbolTable ST;`

```
variable_inc: VARIABLE_NAME INCREMENT SEMI
{if not ST.lookup(x):
    raise SymbolTableException;
else:
    ... // continue}
```

Say we are matched string:  
`x++;`

# How to track scope?

- `SymbolTable ST;`

```
statement : variable_inc  
          | declare_variable
```

```
statement_list : statement_list statement  
              | statement
```

*why do we have the statement list declared like this?*

# How to track scope?

- `SymbolTable ST;`

```
statement : variable_inc  
          | declare_variable
```

```
statement_list : statement_list statement  
              | statement
```

*adding in scope*

# How to track scope?

- `SymbolTable ST;`

```
statement : variable_inc  
          | declare_variable  
          | LBAR statement_list RBAR
```

```
statement_list : statement_list statement  
              | statement
```

# How to track scope?

- `SymbolTable ST;`

statement : **LBAR** statement\_list **RBAR**

start a new scope S

remove the scope S

# How to track scope?

- Symbol table
- **four** methods:
  - **lookup(id)** : lookup an id in the symbol table.  
Returns None if the id is not in the symbol table.
  - **insert(id,info)** : insert a new id into the symbol table along with a set of information about the id.
  - **push\_scope()** : push a new scope to the symbol table
  - **pop\_scope()** : pop a scope from the symbol table

# How to track scope?

- `SymbolTable ST;`

statement : **LBAR** statement\_list **RBAR**

start a new scope S

remove the scope S



# How to track scope?

- `SymbolTable ST;`

statement : **LBAR** statement\_list **RBAR**

start a new scope S

remove the scope S

*How to write a production action here?*

# How to track scope?

- `SymbolTable ST;`

`statement : start_scope statement_list RBAR`

`start_scope : LBAR`

*add a new production rule!*

# How to track scope?

- SymbolTable ST;

```
statement : start_scope statement_list RBAR  
{
```

```
start_scope : LBAR  
{
```

# How to track scope?

- `SymbolTable ST;`

```
statement : start_scope statement_list RBAR  
{ST.pop_scope( )}
```

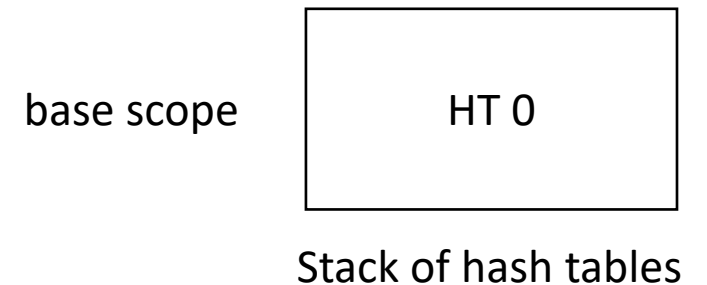
```
start_scope : LBAR  
{ST.push_scope( )}
```

# How to implement a symbol table?

- Thoughts? What data structures are good at mapping strings?
- Symbol table
- **four** methods:
  - **lookup(id)** : lookup an id in the symbol table.  
Returns None if the id is not in the symbol table.
  - **insert(id, info)** : insert a new id into the symbol table along with a set of information about the id.
  - **push\_scope()** : push a new scope to the symbol table
  - **pop\_scope()** : pop a scope from the symbol table

# How to implement a symbol table?

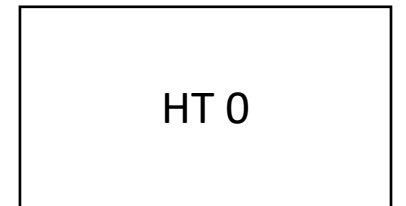
- Many ways to implement:
- A good way is a stack of hash tables:



# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:

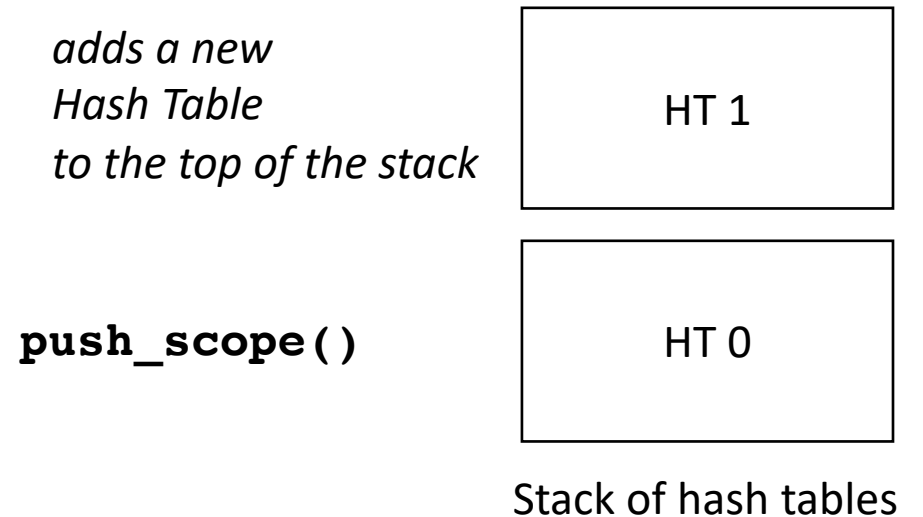
**push\_scope()**



Stack of hash tables

# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:

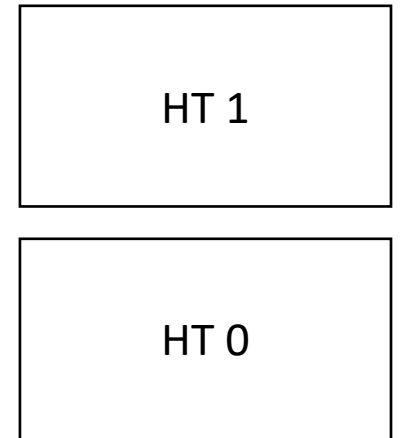




# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:

**`insert(id, data)`**



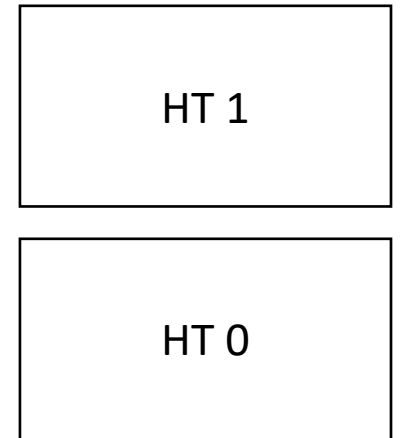
Stack of hash tables

# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:

**`insert(id, data)`**

`insert(id -> data)` at  
top hash table

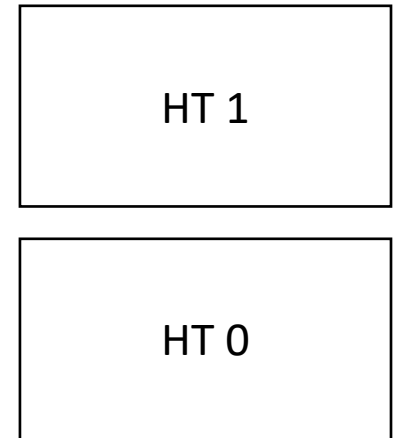


Stack of hash tables

# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:

**lookup(id)**



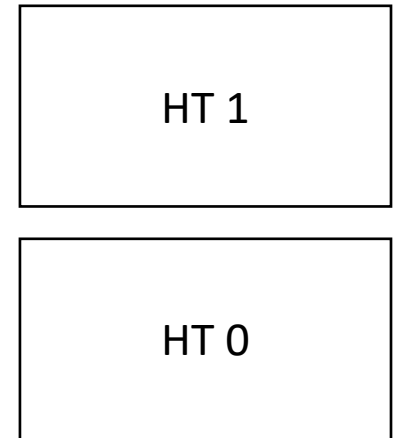
Stack of hash tables

# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:

**lookup(id)**

check here  
first



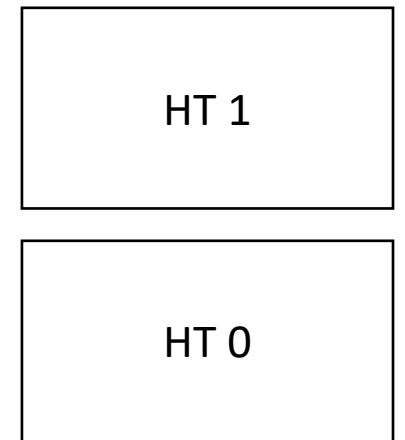
Stack of hash tables

# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:

**lookup(id)**

then check  
here

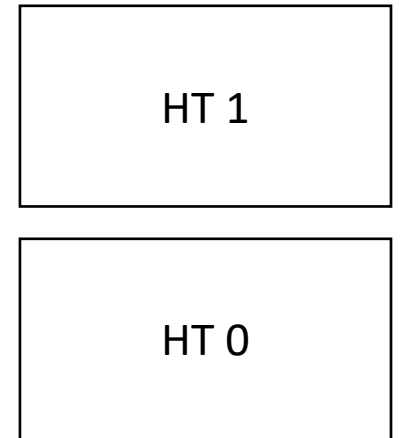


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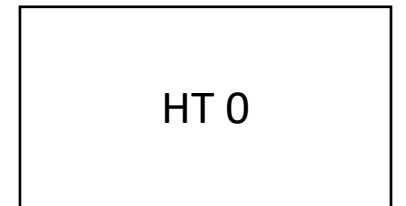
**pop\_scope ( )**



Stack of hash tables

# How to implement a symbol table?

- Many ways to implement:
- A good way is a stack of hash tables:



Stack of hash tables

# How to implement a symbol table?

- Example

```
int x = 0;
int y = 0;
{
    int y = 0;
    x++;
    y++;
}
x++;
y++;
```



x = 2  
y = 1

Stack of hash tables



# See you on Friday!

- You should have everything you need to know to work on Homework part 1!
- Next class: Parsing regular expressions with derivatives
- Office hours tomorrow: (2 - 3 pm)