CSE211: Compiler Design Nov. 5, 2021

• **Topic**: restructuring loops



Announcements

- Homework 3 is due Nov. 17
 - 1 more office hour before then (next Thursday)
 - part 1 and 2: generating c code from python
 - part 3: creating and checking z3 constraints

Paper/Project proposals

- Please start thinking about these.
 - Message me for recommendations
 - Tell me what you're interested in so we can find a good fit!
- Proposals due on Nov. 14 (less than 2 weeks)
 - Please be pro-active about this. If you don't have one in mind, please send me an email with some of your interests ASAP
- Midterm is a good indicator for how the final will be.

CSE211: Compiler Design Nov. 3, 2021

• **Topic**: restructuring loops



Review

- Compiler approach for checking if DOALL loops are safe to do in parallel
 - What is a DOALL loop?
 - What conditions are required for safety?

Review

• Creating constraints

```
for (i = 0; i < 128; i++) {
    a[i]= a[i]*2;
}
two integers: i<sub>x</sub> != i<sub>y</sub>
    i<sub>x</sub> >= 0
    i<sub>x</sub> < 128
    i<sub>y</sub> >= 0
    i<sub>y</sub> < 128
    write-write conflict    i<sub>x</sub> == i<sub>y</sub>
    read-write conflict    i<sub>x</sub> == i<sub>y</sub>
```

Ask if these constraints are satisfiable (if so, it is not safe to parallelize)

```
for (i = 0; i < 128; i++) {
    a[i%64]= a[i+64]*2;
}</pre>
```

```
\begin{array}{rll} \text{two integers: } i_x \; !=\; i_y \\ i_x \; >=\; 0 \\ i_x \; <\; 128 \\ i_y \; >=\; 0 \\ push \; bounds & i_y \; <\; 128 \\ constraints \end{array}
```

```
for (i = 0; i < 128; i++) {
    a[i%64]= a[i+64]*2;
}</pre>
```

```
two integers: i_x != i_y

i_x \ge 0

i_x < 128

i_y \ge 0

push bounds

constraints

i_x & 64 == i_y & 64
```

write-write conflict checking

```
for (i = 0; i < 128; i++) {
    a[i%64]= a[i+64]*2;
}</pre>
```

	two integers: $i_x != i_y$
	i _x >= 0
	$i_x < 128$
	i _y >= 0
push bounds	i _y < 128
constraints	

```
for (i = 0; i < 128; i++) {
    a[i%64]= a[i+64]*2;
}</pre>
```

```
two integers: i_x != i_y

i_x \ge 0

i_x < 128

i_y \ge 0

push bounds

constraints

i_x & 64 == i_y + 64
```

read-write conflict checking

Moving onto loop structures

• Locality is key for good parallel performance:

- Locality is key for good parallel performance:
- Two types of locality:
 - Temporal locality
 - Spatial locality



- Locality is key for good parallel performance:
- Two types of locality:
 Temporal locality
 Spatial locality
 r1 = a[2];
 r2 = a[3];

how far apart can memory locations be?

• Locality is key for good parallel performance:

good data locality: cores will spend most of their time accessing private caches



• Locality is key for good parallel performance:

Bad data locality: cores will pressure and thrash shared memory resources





Row major



Row major



Row major



Column major? Fortran Matlab R



Column major? Fortran Matlab R



x1 = a[0,0]; x2 = a[0,1];

good pattern for row major bad pattern for column major

-			

unrolled row major: still has locality

x1 = a[x,y]; x2 = a[x, y+1];

good pattern for row major bad pattern for column major

x1 = a[x,y]; x2 = a[x, y+1];

good pattern for row major bad pattern for column major unrolled column major: Bad locality



x1 = a[0,0]; x2 = a[1, 0];

good pattern for column major bad pattern for row major

row major unrolled: bad spatial locality

unrolled

column

major:

x1 = a[x, y]; $x^{2} = a[x+1, y];$

good pattern for column major bad pattern for row major



How much does this matter?

which will be faster? by how much?

Demo

How to reorder loop nestings?

- For a DOALL loop, if loop bounds are independent, they can simply be re-ordered.
- If they are dependent...

bad nesting order for row-major!

bad nesting order for row-major!

but iteration variables are dependent

bad nesting order for row-major!

but iteration variables are dependent

```
loop constraints
y >= 0
y <= 5
x >= y
x <= 7</pre>
```

x = y

loop constraints
y >= 0
y <= 5
x >= y
x <= 7</pre>

У

System with N variables can be viewed as an N dimensional polyhedron



Fourier-Motzkin elimination:

- Given a system of inequalities with N variables, reduce it to a system with N 1 variables.
- A system of inequalities describes an N-dimensional polyhedron. Produce a system of equations that projects the polyhedron onto an N-1 dimensional space

x = y

loop constraints
y >= 0
y <= 5
x >= y
x <= 7</pre>

System with N variables can be viewed as an N dimensional polyhedron



Fourier-Motzkin elimination:

• To eliminate variable x_i :

For every pair of lower bound L_i and upper bound U_i on x_i , create:

 $L_i \leq x_i \leq U_i$ Then simply remove x_i :

 $L_i \leq U_i$

```
for (y = 0; y <= 5; y++) {
  for (x = y; x <= 7; x++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

All pairs of upper/lower bounds on y:

loop constraints

- y >= 0
- y <= 5
- x >= y
- x <= 7

0 <= y <= 5

0 <= y <= x

```
for (y = 0; y <= 5; y++) {
   for (x = y; x <= 7; x++) {
      a[x,y] = b[x,y] + c[x,y];
   }
}</pre>
```

All pairs of upper/lower bounds on y:

loop constraints

y >= 0 y <= 5 x >= y

```
for (y = 0; y <= 5; y++) {
  for (x = y; x <= 7; x++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

All pairs of upper/lower bounds on y:

loop constraints
y >= 0
y <= 5
x >= y
x <= 7</pre>

 $0 \le y \le x$

0 <= y <= 5

Then eliminate y:

0 <= 5 0 <= x

```
for (y = 0; y <= 5; y++) {
  for (x = y; x <= 7; x++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

All pairs of upper/lower bounds on y:

loop constraints

y >= 0 y <= 5 x >= y x <= 7 $0 \le y \le 5$ $0 \le y \le x$

Then eliminate y:

<mark>0 <= 5</mark> 0 <= x

```
for (y = 0; y <= 5; y++) {
  for (x = y; x <= 7; x++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

All pairs of upper/lower bounds on y:

loop constraints
y >= 0
y <= 5</pre>

x >= y x <= 7 0 <= y <= x

0 <= y <= 5

Then eliminate y:

0 <= x

```
for (y = 0; y <= 5; y++) {
  for (x = y; x <= 7; x++) {
    a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

All pairs of upper/lower bounds on y:

loop constraints $y \ge 0$

- y <= 5
- x >= y

x <= 7

 $0 \le y \le x$ Then eliminate y:

0 <= y <= 5

loop constraints without y:

x >= 0 x <= 7

0 <= x

x = y

loop constraints
y >= 0
y <= 5
x >= y
x <= 7</pre>

System with N variables can be viewed as an N dimensional polyhedron



Reording Loop bounds:

- Given a new order: $[x_0, x_1, x_2, ..., x_n]$
- For each variable x_i : perform Fourier-Motzkin elimination to eliminate any variables that come after x_i in the new order.
- Instantiate loop conditions for x_i, potentially using max/min operators

loop constraints

y >= 0 y <= 5

x >= y

```
for (y = 0; y <= 5; y++) {
  for (x = y; x <= 7; x++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

new order: [x,y]

for x: eliminate y using FM elimination:

loop constraints

- y >= 0 y <= 5
- $x \ge y$

new order: [x,y]

for x: eliminate y using FM elimination:

x loop constraints without y:

x >= 0 x <= 7

loop constraints

y >= 0

y <= 5

x >= y

loop constraints

y >= 0

y <= 5

x >= y

x <= 7

new order: [x,y]

for x: eliminate y using FM elimination:

x loop constraints without y:

x >= 0 x <= 7

- y loop constraints:
- y >= 0
- y <= 5
- y <= x

new order: [x,y]

for x: eliminate y using FM elimination:

x loop constraints without y:

x >= 0 x <= 7

- y loop constraints:
- y >= 0
- <mark>y <= 5</mark>
- <mark>y <= x</mark>

- loop constraints
- y >= 0
- y <= 5
- x >= y
- x <= 7

new order: [x,y]

for x: eliminate y using FM elimination:

x loop constraints without y:

x >= 0 x <= 7

y loop constraints: y >= 0 y <= min(x,5)</pre>

loop constraints

- y >= 0
- y <= 5
- x >= y
- x <= 7

У

x loop constraints without y:

x >= 0 x <= 7

y loop constraints: y >= 0 y <= min(x,5)</pre>



Х

Reordering loop bounds

- only works if loop increments by 1; assumes a closed polyhedron
- best performance when array indexes are simple:
 - e.g.: a[x,y]
 - harder with, e.g.: a[x*5+127, y+x*37]
 - There exists schemes to automatically detect locality. Reach chapter 10 of the Dragon book
- compiler implementation allows exploration and auto-tuning

 In some cases, there might not be a good nesting order for all accesses:

 $A = B + C^T$

 In some cases, there might not be a good nesting order for all accesses:

$$A = B + C^T$$

В







cold miss for all of them

Α

 In some cases, there might not be a good nesting order for all accesses:

$$A = B + C^T$$

В



Hit on A and B. Miss on C

Α





С

 In some cases, there might not be a good nesting order for all accesses:

$$A = B + C^T$$

В



Α





С

Hit on A and B. Miss on C

• Blocking operates on smaller chunks to exploit locality in column increment accesses. Example 2x2

A

$$A = B + C^T$$







 Blocking operates on smaller chunks to exploit locality in column increment accesses. Example 2x2

$$A = B + C^T$$







• Blocking operates on smaller chunks to exploit locality in column increment accesses. Example 2x2

 $A = B + C^T$







cold miss for all of them

• Blocking operates on smaller chunks to exploit locality in column increment accesses. Example 2x2

A

 $A = B + C^T$



С



Miss on C

• Blocking operates on smaller chunks to exploit locality in column increment accesses. Example 2x2

$$A = B + C^T$$



A





Miss on A,B, hit on C

• Blocking operates on smaller chunks to exploit locality in column increment accesses. Example 2x2

$$A = B + C^T$$



Α





Hit on all!

• Add two outer loops for both x and y

```
for (int x = 0; x < SIZE; x++) {
    for (int y = 0; y < SIZE; y++) {
        a[x*SIZE + y] = b[x*SIZE + y] + c[y*SIZE + x];
    }
}</pre>
```

Add two outer loops for both x and y

```
for (int xx = 0; xx < SIZE; xx += B) {
  for (int yy = 0; yy < SIZE; yy += B) {
    for (int x = xx; x < xx+B; x++) {
      for (int y = yy; y < yy+B; y++) {
         a[x*SIZE + y] = b[x*SIZE + y] + c[y*SIZE + x];
      }
    }
}</pre>
```

Add two outer loops for both x and y

```
for (int xx = 0; xx < SIZE; xx += B) {
  for (int yy = 0; yy < SIZE; yy += B) {
    for (int x = xx; x < xx+B; x++) {
      for (int y = yy; y < yy+B; y++) {
        a[x*SIZE + y] = b[x*SIZE + y] + c[y*SIZE + x];
      }
    }
}</pre>
```

Add two outer loops for both x and y

```
for (int xx = 0; xx < SIZE; xx += B) {
  for (int yy = 0; yy < SIZE; yy += B) {
    for (int x = xx; x < xx+B; x++) {
      for (int y = yy; y < yy+B; y++) {
         a[x*SIZE + y] = b[x*SIZE + y] + c[y*SIZE + x];
      }
    }
}</pre>
```

Add two outer loops for both x and y

```
for (int xx = 0; xx < SIZE; xx += B) {
  for (int yy = 0; yy < SIZE; yy += B) {
    for (int x = xx; x < xx+B; x++) {
      for (int y = yy; y < yy+B; y++) {
        a[x*SIZE + y] = b[x*SIZE + y] + c[y*SIZE + x];
      }
    }
}</pre>
```

Demo

Next class

• Topics:

• Implementing parallelism for DOALL loops

• Enjoy your weekend