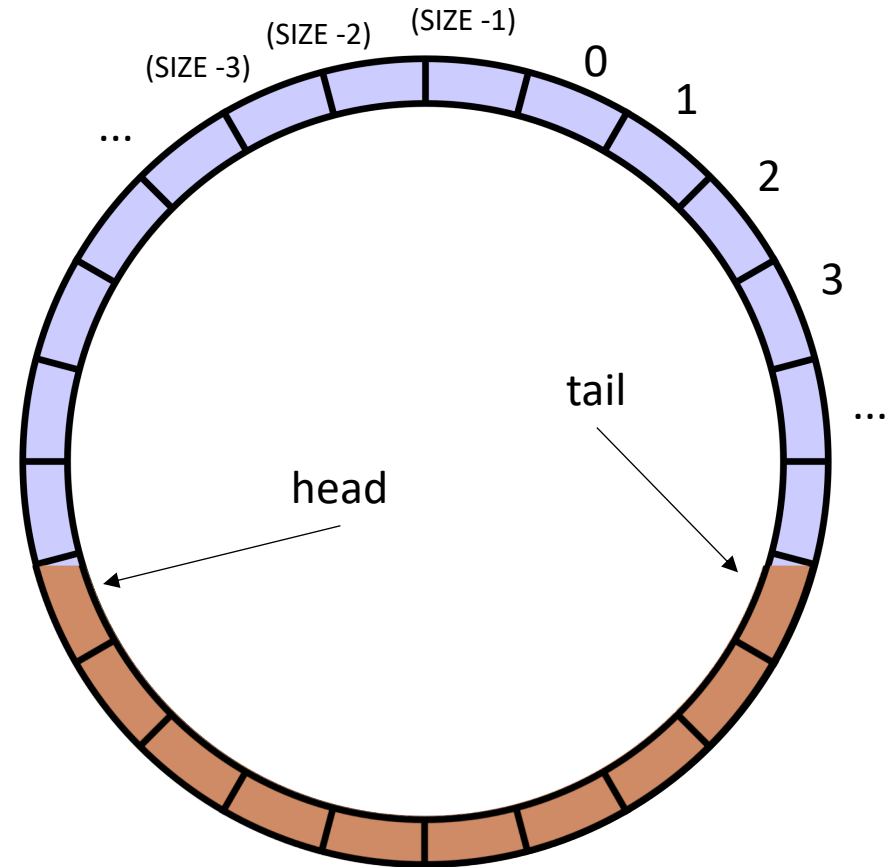


CSE113: Parallel Programming

Feb. 13, 2023

- **Topics:**

- Input/output queues
- Producer consumer queues
 - Synchronous
 - Circular buffer



Announcements

- HW1 grades will be out by the end of the day
 - Let us know ASAP if there are issues
- Homework 2 has a last due date today
 - We will keep an eye on Piazza and try to ask questions asked before 5 pm
- Homework 3 will be released today by midnight
 - Due in 10 days + 4 free late days

Announcements

- Midterm out!
 - asynchronous, 1 week (no time limit)
 - Open note, open internet (to a reasonable extent: no googling exact questions or asking questions on forums)
 - do not discuss with classmates AT ALL while the test is active
 - **No late tests will be accepted.**

- **Prioritize midterm next week!**

Previous quiz

What is the relationship between linearizable (L) and sequentially consistent (SC)?

-
- Objects can be one or the other, but not both

 - Objects that are L are also SC, but not the other way around

 - Objects that are SC are also L, but not the other way around

 - SC and L are the different definitions for the same concept

Previous quiz

Lock-free data structures are technically undefined because they contain data conflicts

True

False

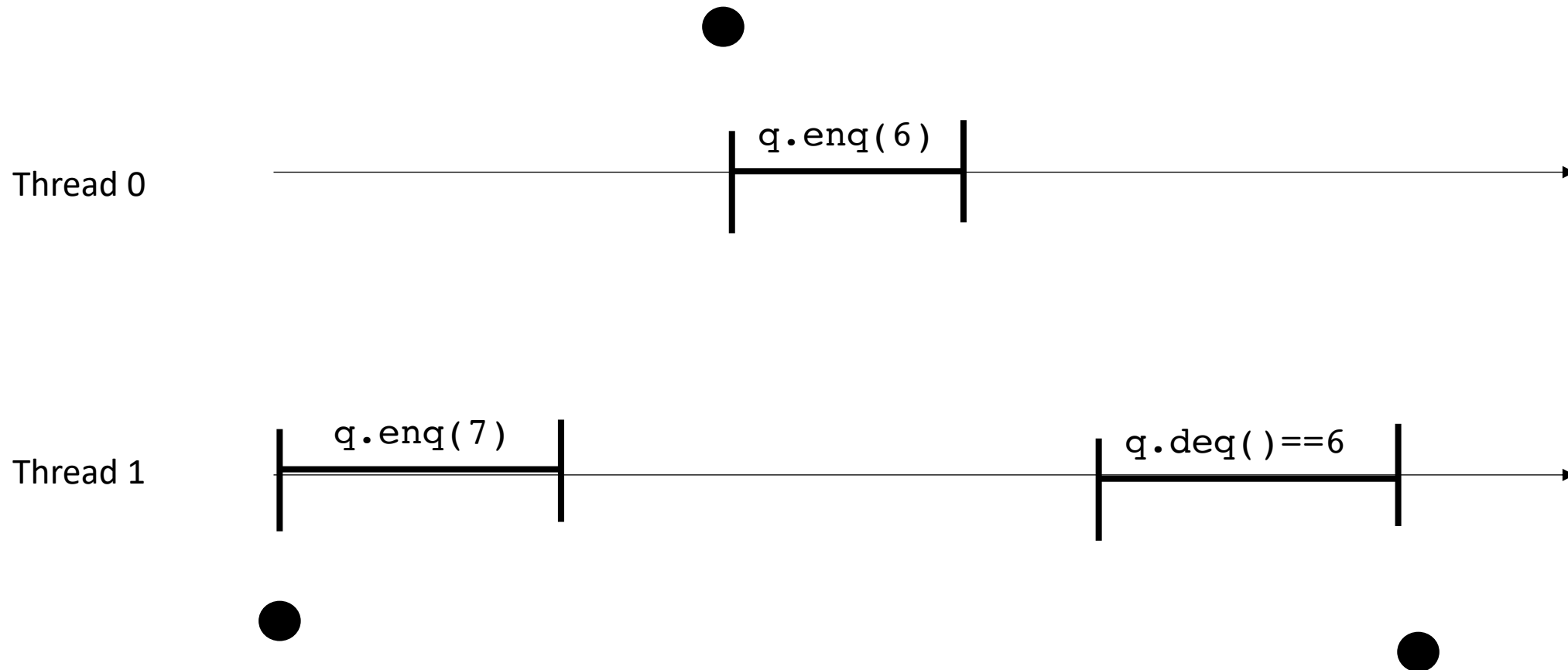
Review

Linearizability

Linearizability

each command gets a linearization point.

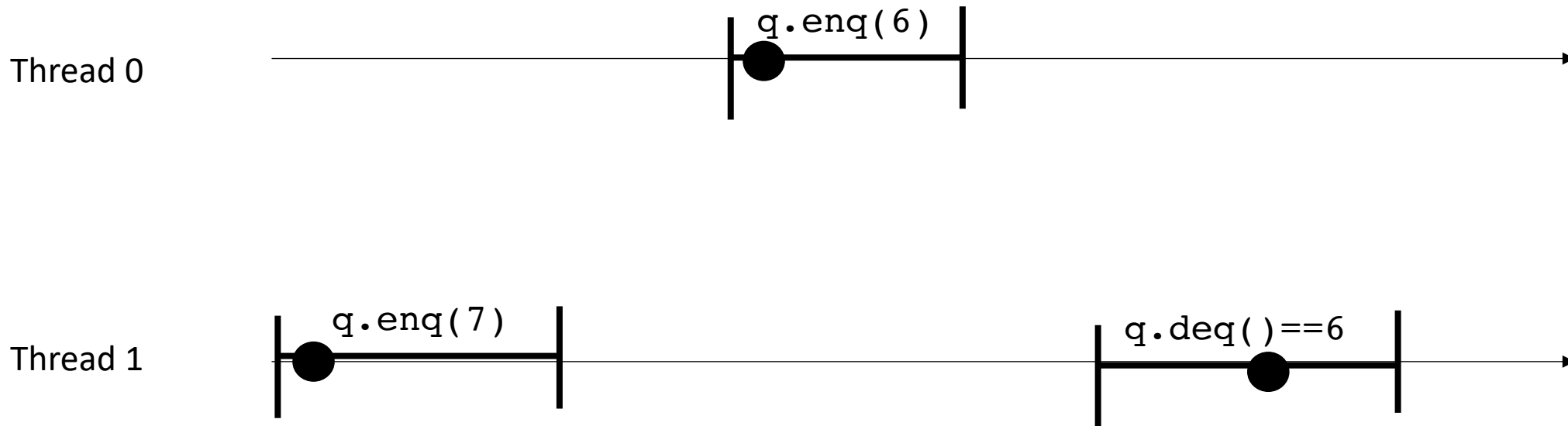
You can place the point anywhere between its innovation and response!



Linearizability

each command gets a linearization point.

You can place the point anywhere between its innovation and response!

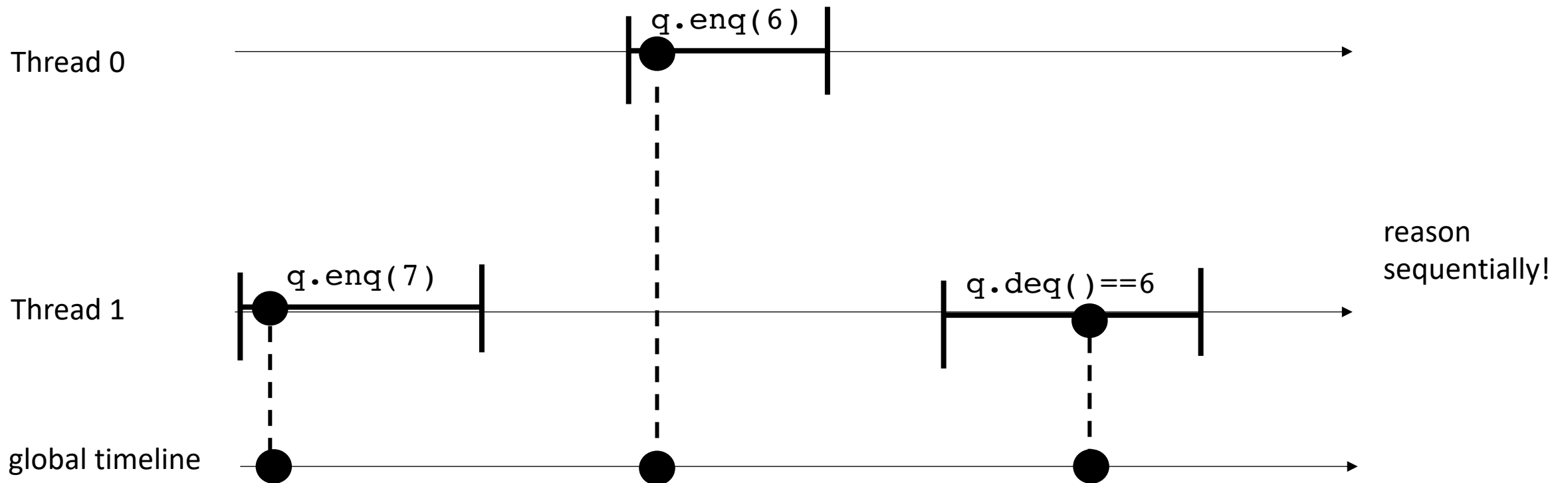


Linearizability

each command gets a linearization point.

You can place the point anywhere between its innovation and response!

Project the linearization points to a global timeline

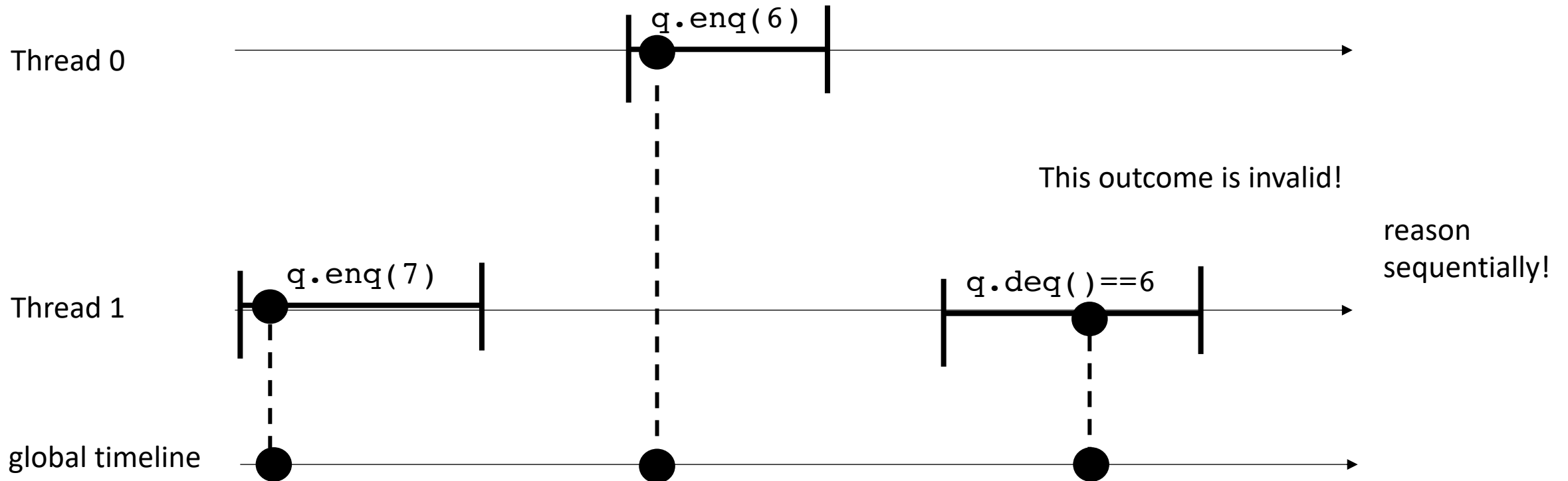


Linearizability

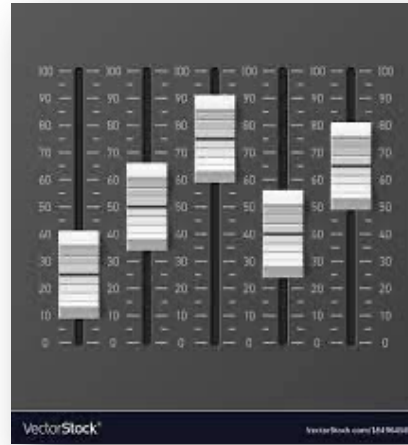
each command gets a linearization point.

You can place the point anywhere between its innovation and response!

Project the linearization points to a global timeline



Linearizability

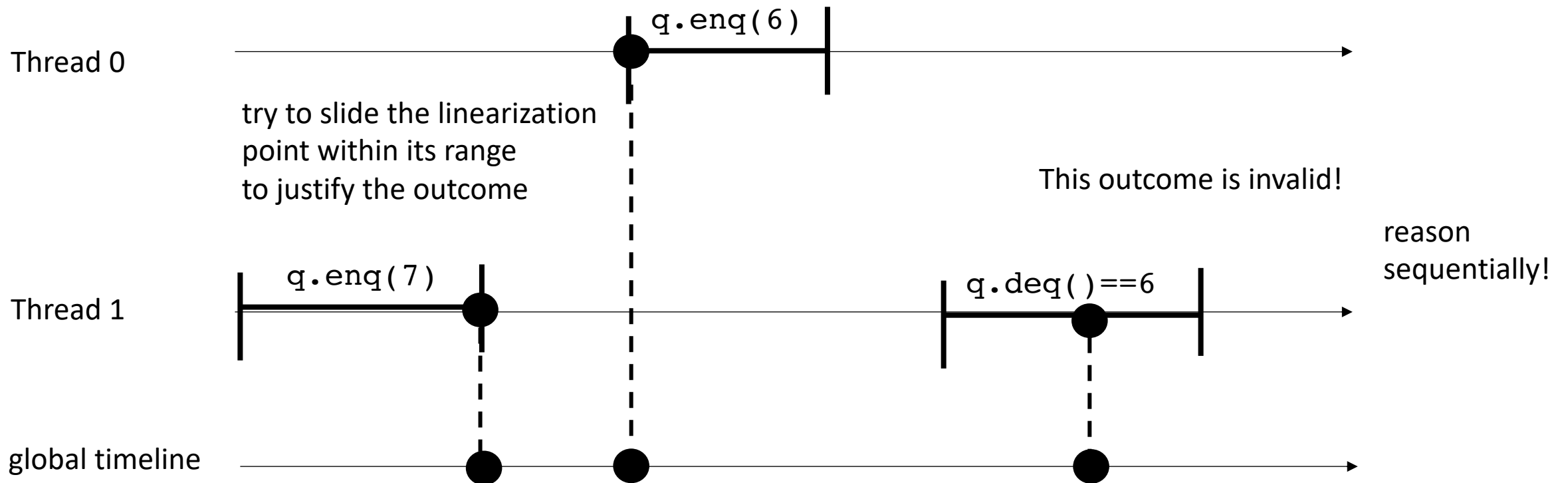


each command gets a linearization point.

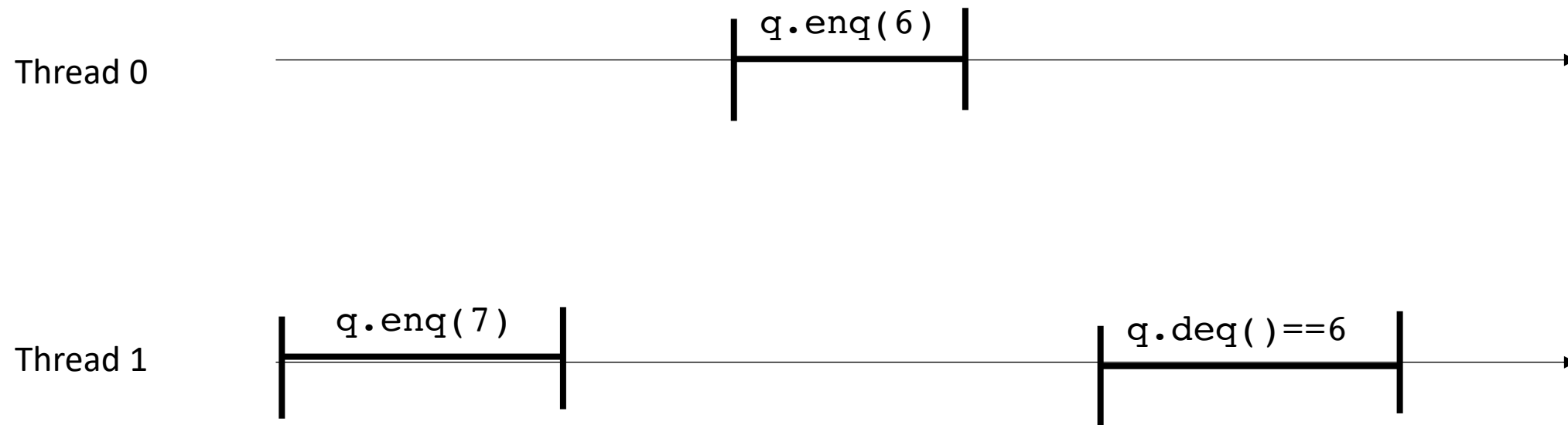
You can place the point anywhere between its innovation and response!

Project the linearization points to a global timeline

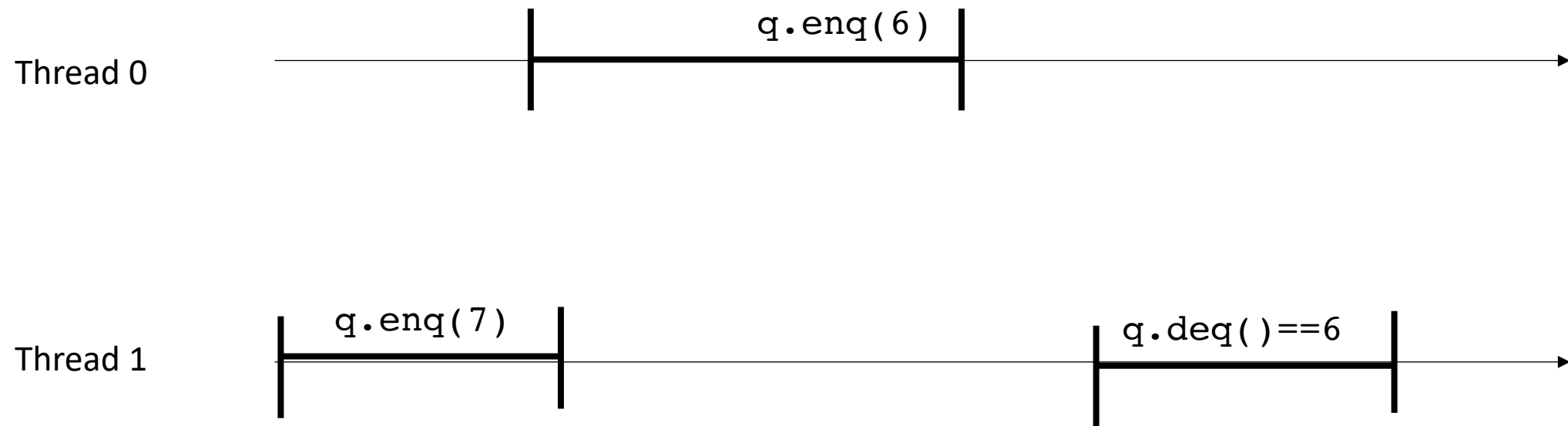
slider game!



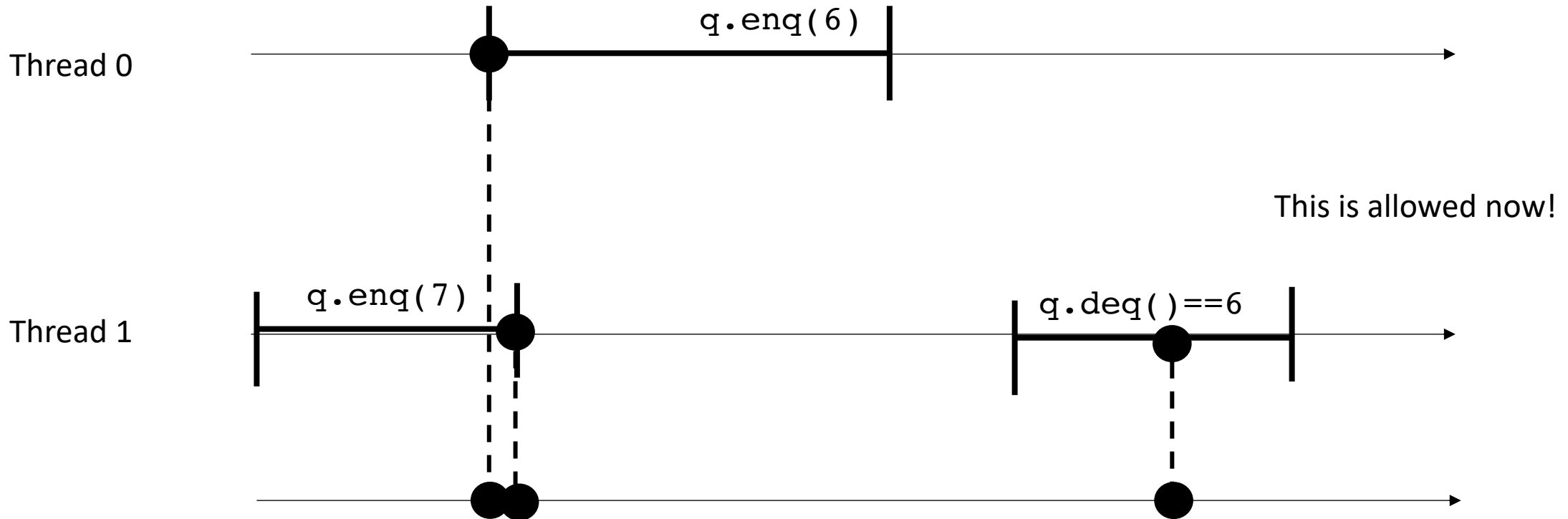
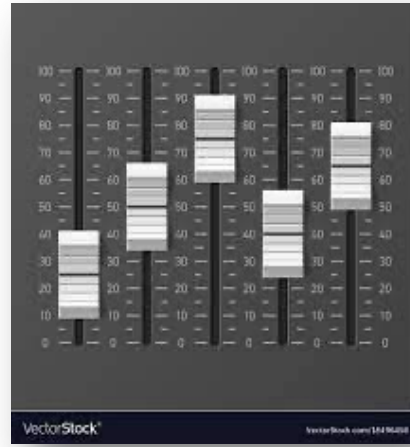
Linearizability



Linearizability

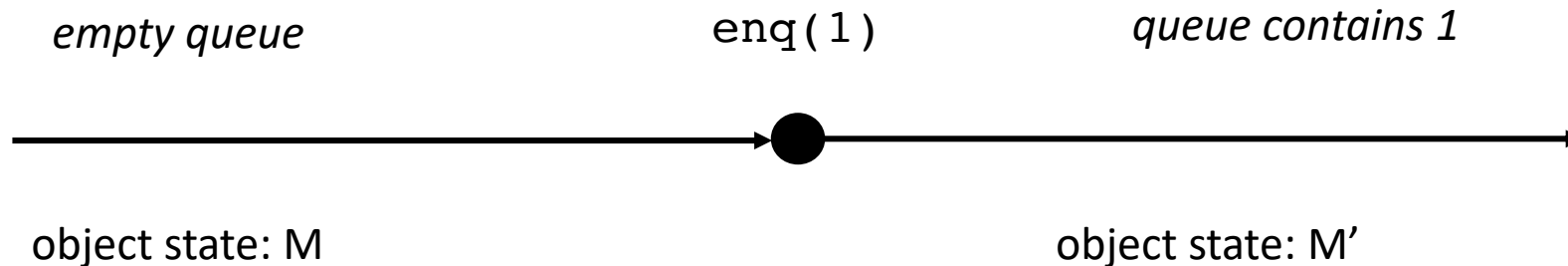


Linearizability



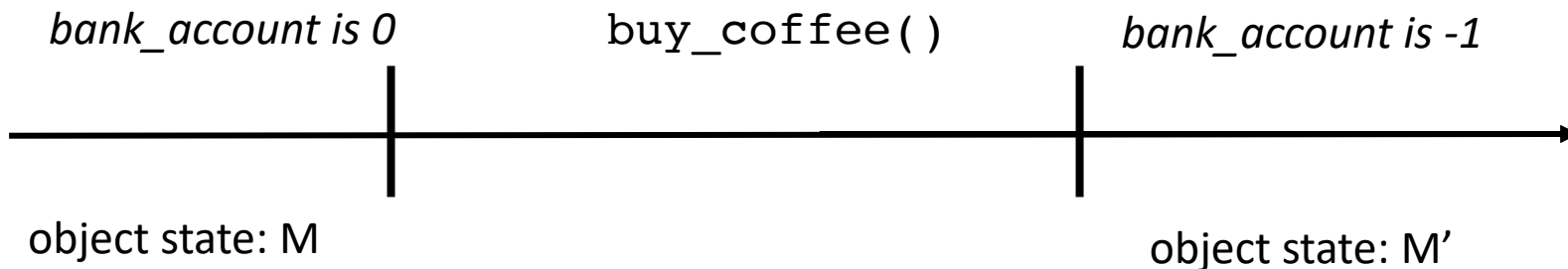
Linearizability

- How do we write our programs to be linearizable?
 - Identify the linearizability point
 - One indivisible region (e.g. an atomic store, atomic load, atomic RMW, or critical section) where the method call takes effect. Modeled as a point.



Linearizability

- Locked data structures are linearizable.

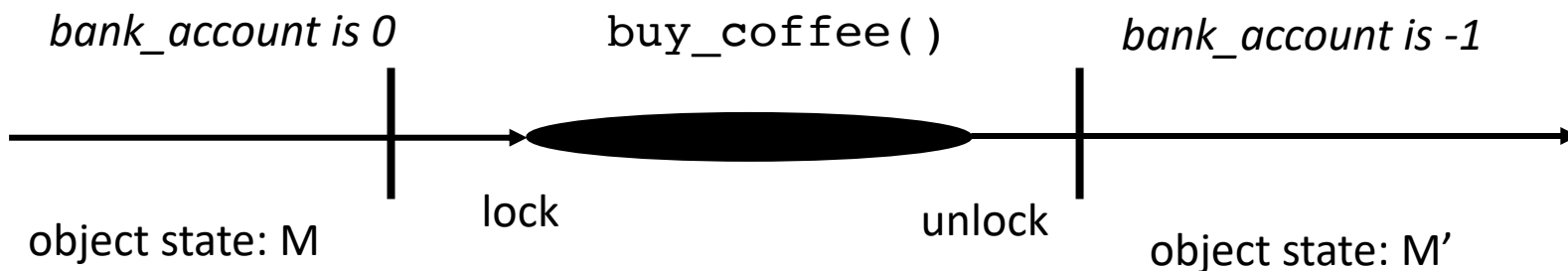


```
class bank_account {  
    public:  
    bank_account() {  
        balance = 0;  
    }  
  
    void buy_coffee() {  
        m.lock();  
        balance -= 1;  
        m.unlock();  
    }  
  
    void get_paid() {  
        m.lock();  
        balance += 1;  
        m.unlock();  
    }  
  
    private:  
    int balance;  
    mutex m;  
};
```

Linearizability

- Locked data structures are linearizable.

typically modeled as the point the lock is acquired or released

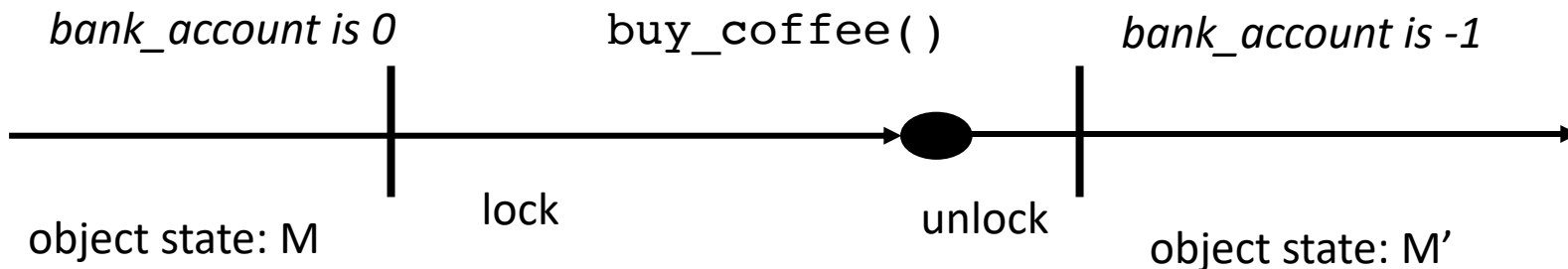


```
class bank_account {  
    public:  
        bank_account() {  
            balance = 0;  
        }  
  
        void buy_coffee() {  
            m.lock();  
            balance -= 1;  
            m.unlock();  
        }  
  
        void get_paid() {  
            m.lock();  
            balance += 1;  
            m.unlock();  
        }  
  
    private:  
        int balance;  
        mutex m;  
};
```

Linearizability

- Locked data structures are linearizable.

*typically modeled as the point the lock is acquired or released
lets say released.*

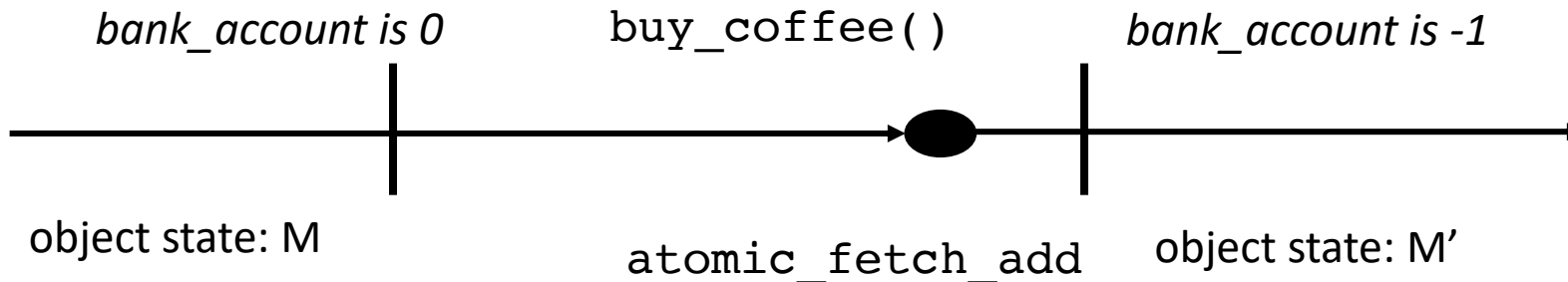


```
class bank_account {  
  public:  
    bank_account() {  
      balance = 0;  
    }  
  
    void buy_coffee() {  
      m.lock();  
      balance -= 1;  
      m.unlock();  
    }  
  
    void get_paid() {  
      m.lock();  
      balance += 1;  
      m.unlock();  
    }  
  
  private:  
    int balance;  
    mutex m;  
};
```

Linearizability

- Our lock-free bank account is linearizable:
 - The atomic operation is the linearizable point

```
class bank_account {  
    public:  
    bank_account() {  
        balance = 0;  
    }  
  
    void buy_coffee() {  
        atomic_fetch_add(&balance, -1);  
    }  
  
    void get_paid() {  
        atomic_fetch_add(&balance, 1);  
    }  
  
    private:  
    atomic_int balance;  
};
```



Input/Output Queues

Concurrent Queues

- List of items, accessed in a first-in first-out (FIFO) way
- *duplicates allowed*
- Methods
 - **enq(x)** put **x** in the list at the end
 - **deq()** remove the item at the front of the queue and return it.
 - **size()** returns how many items are in the queue

Concurrent Queues

- General implementation given in Chapter 10 of the book.
- Similar types of reasoning as the linked list
 - Lots of reasoning about node insertion, node deletion
 - Using atomic RMWs (CAS) in clever ways
- We will think about specialized queues
 - Implementations can be simplified!

Input/Output Queues

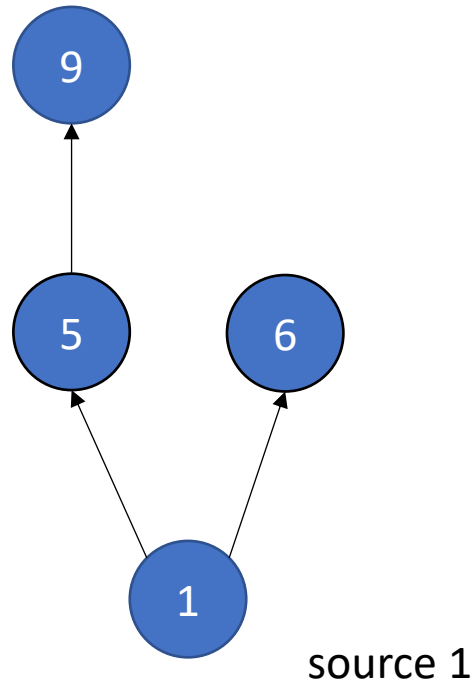
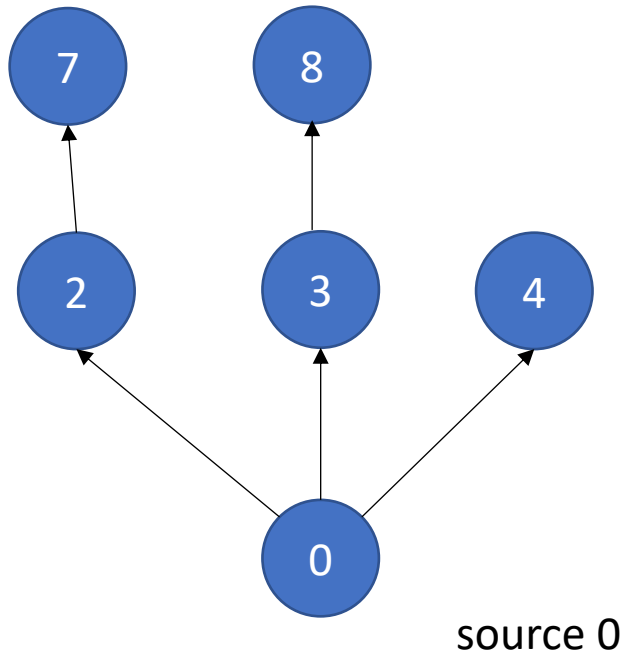
- Queue in which multiple threads read (deq), or write (enq), but not both.
- Why would we want a thing?
- Computation done in phases:
 - First phase prepares the queue (by writing into it)
 - All threads join
 - Second phase reads values from the queue.

Input/Output Queues

- Example: Information flow in graph applications:

Input/Output Queues

- Example: Information flow in graph applications:



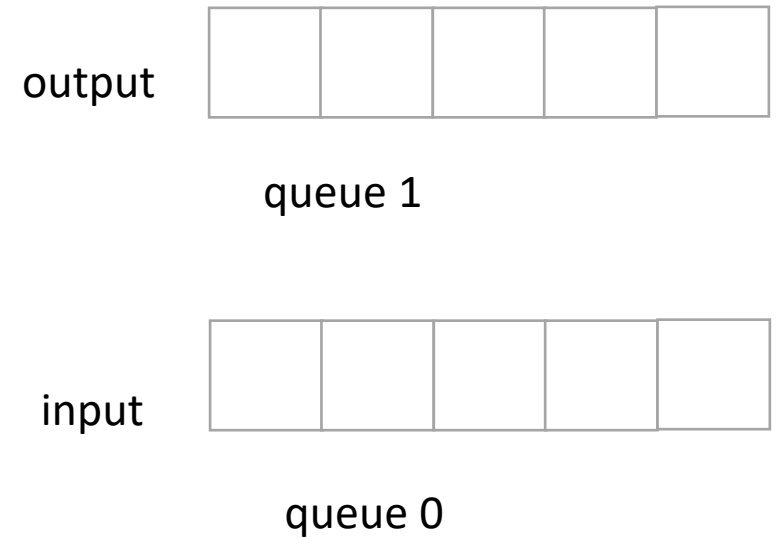
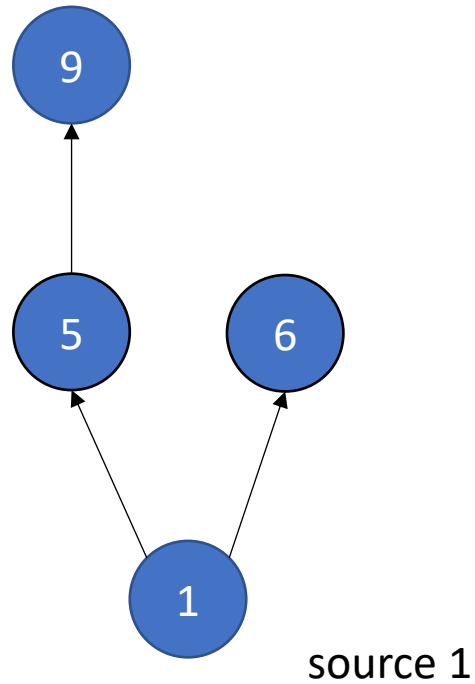
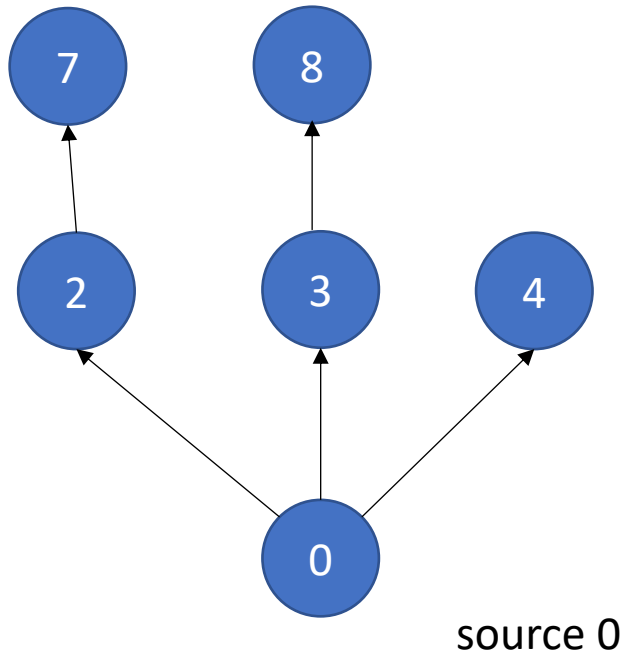
queue 1



queue 0

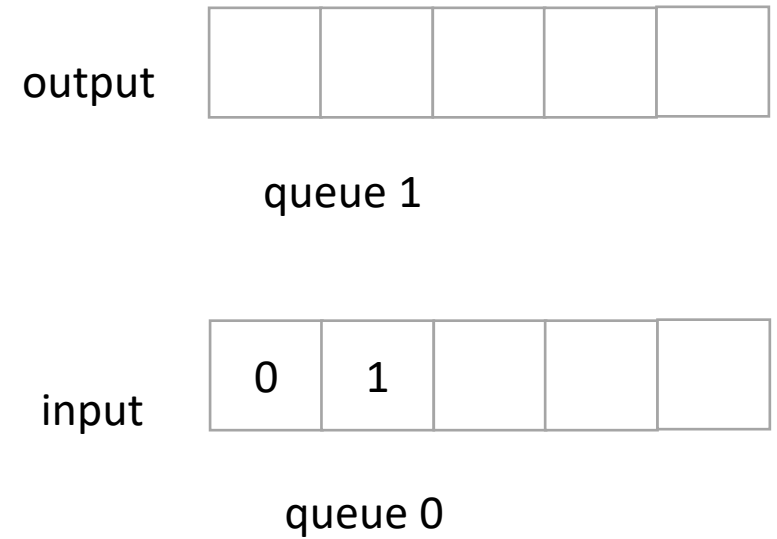
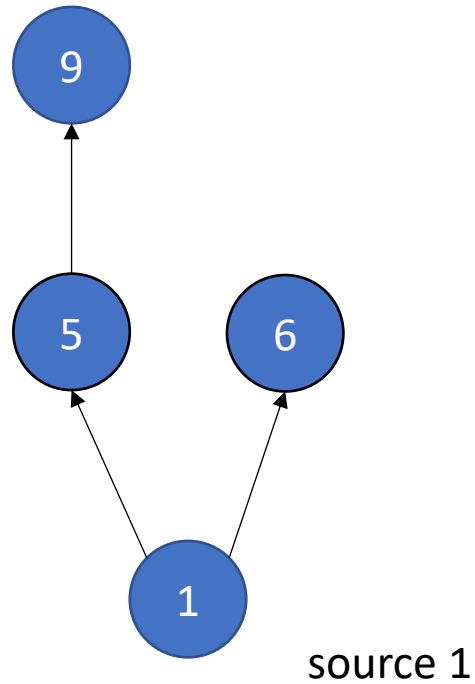
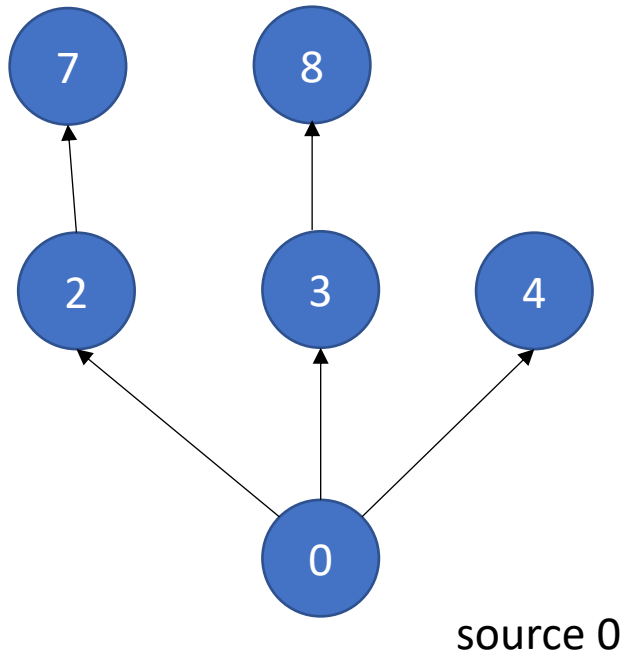
Input/Output Queues

- Example: Information flow in graph applications:



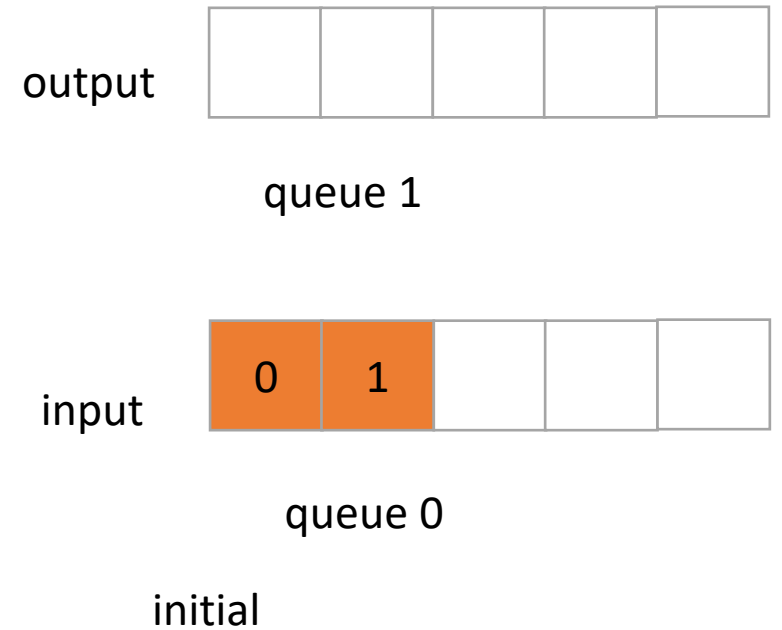
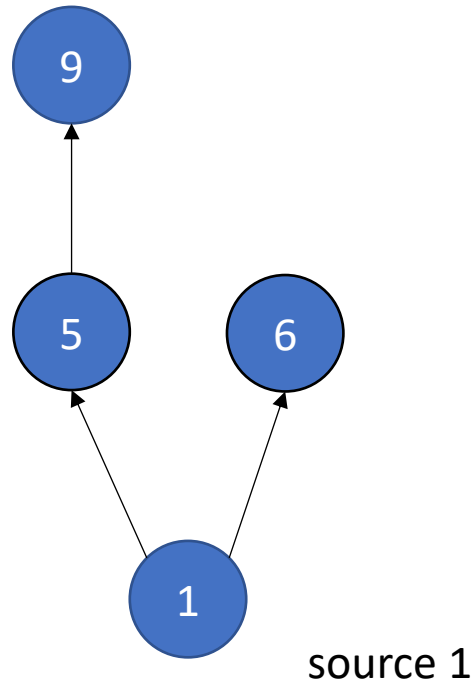
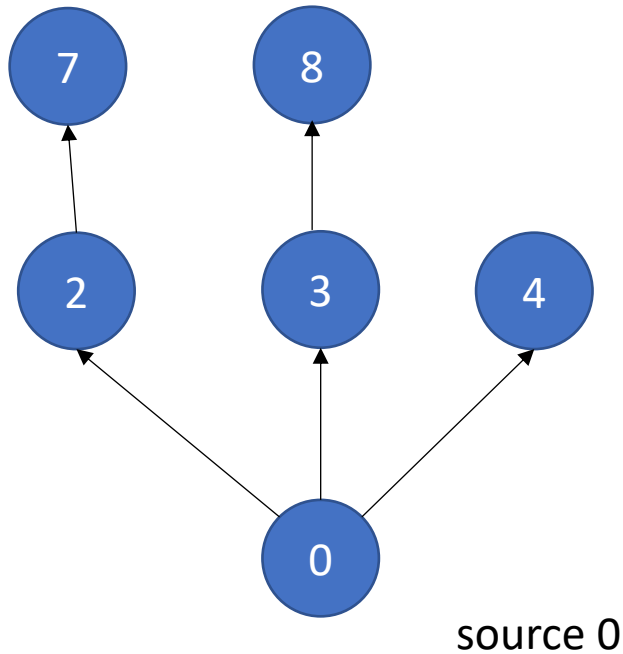
Input/Output Queues

- Example: Information flow in graph applications:



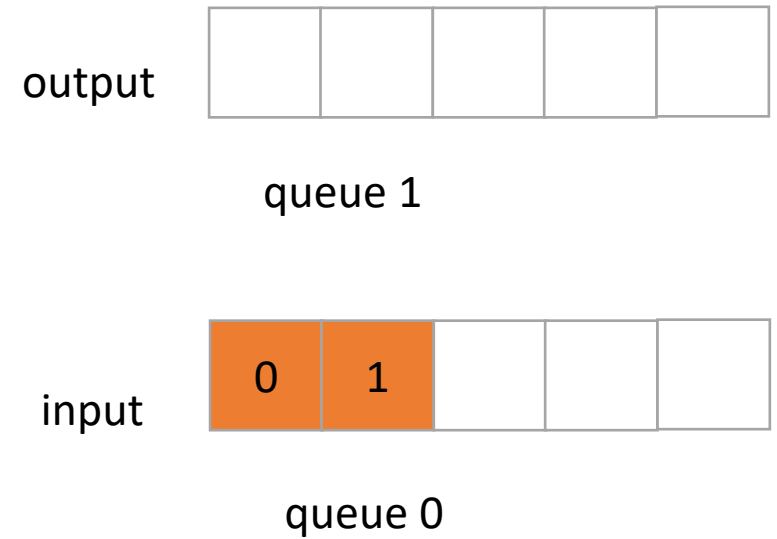
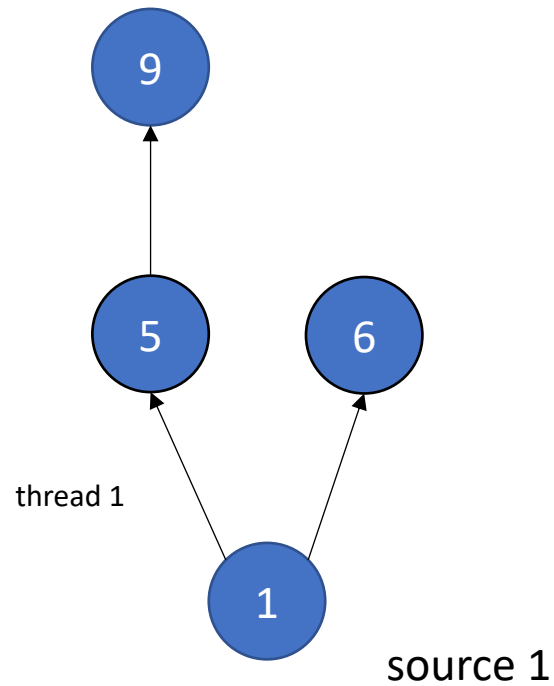
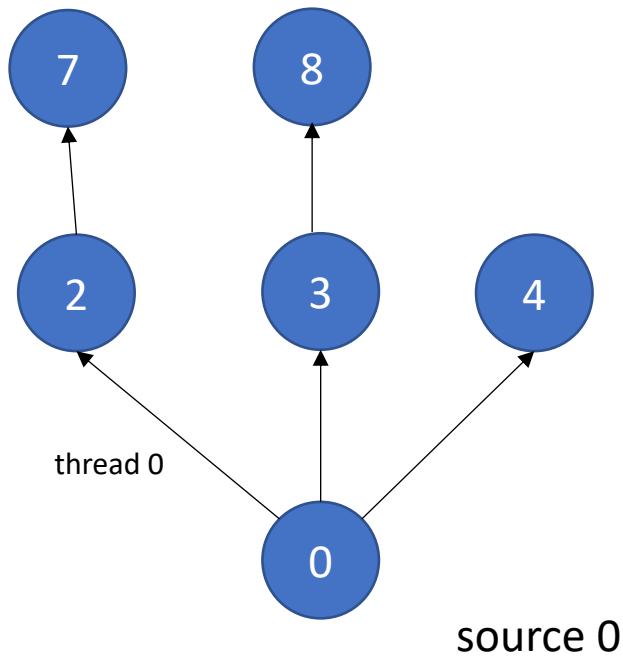
Input/Output Queues

- Example: Information flow in graph applications:



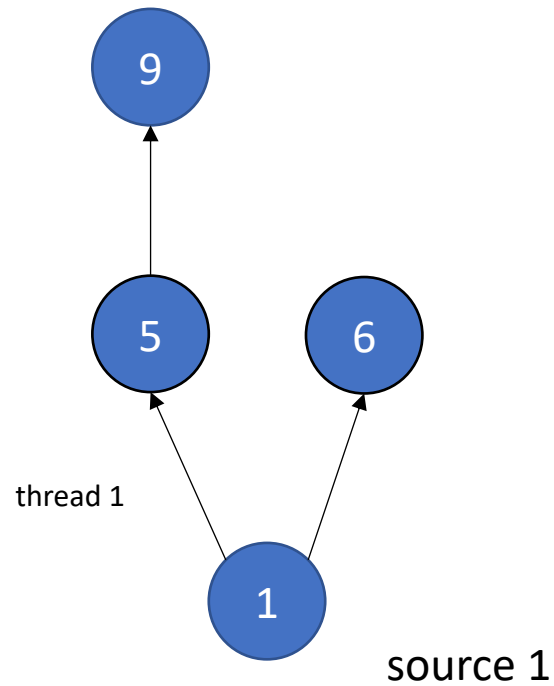
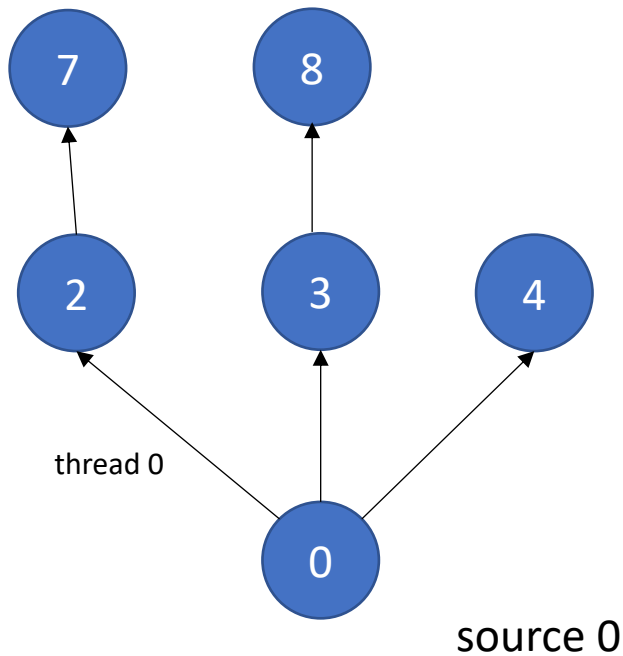
Input/Output Queues

- Example: Information flow in graph applications:

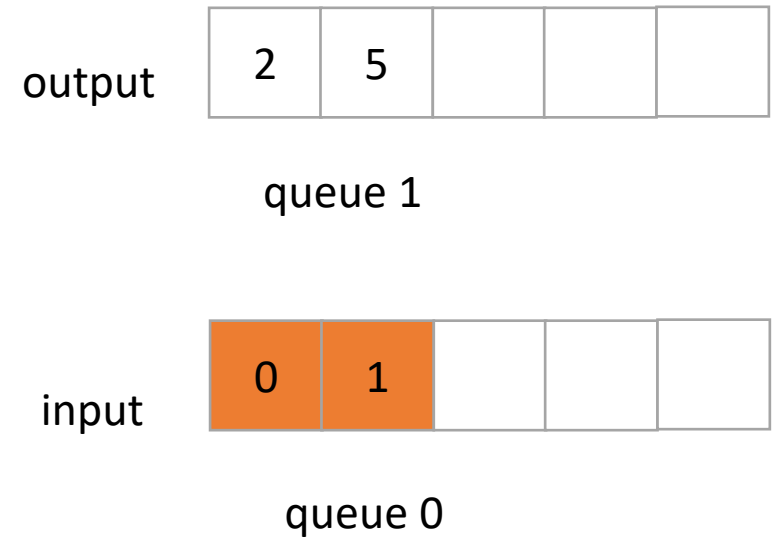


Input/Output Queues

- Example: Information flow in graph applications:

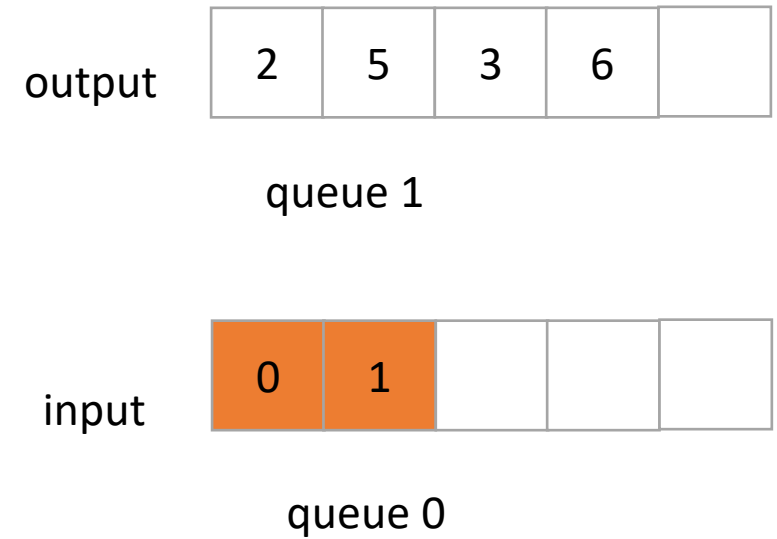
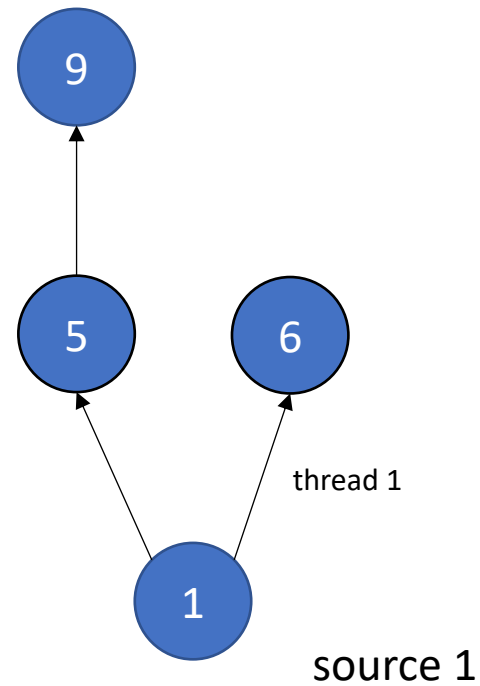
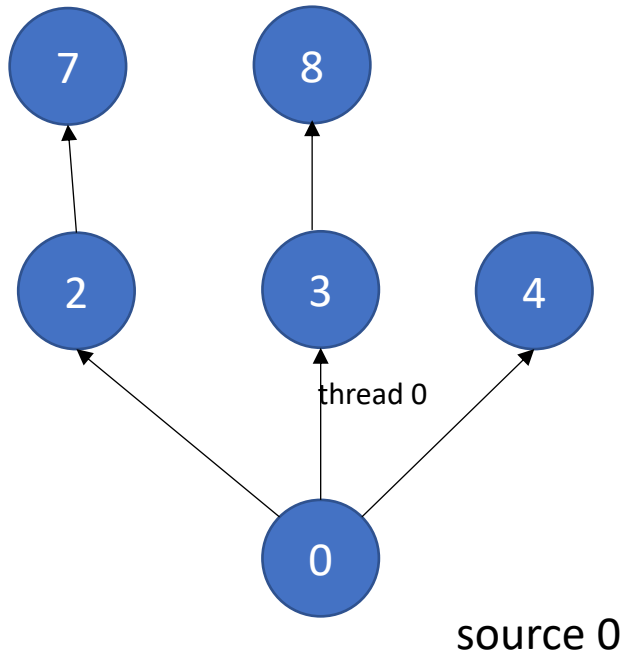


concurrent enqueues!



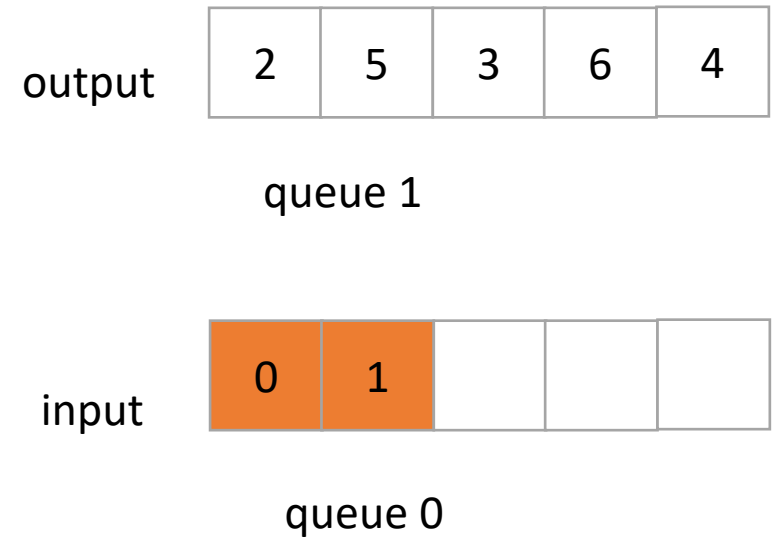
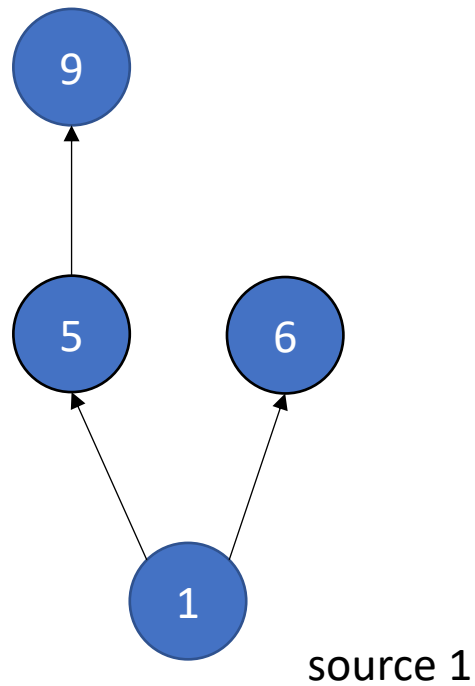
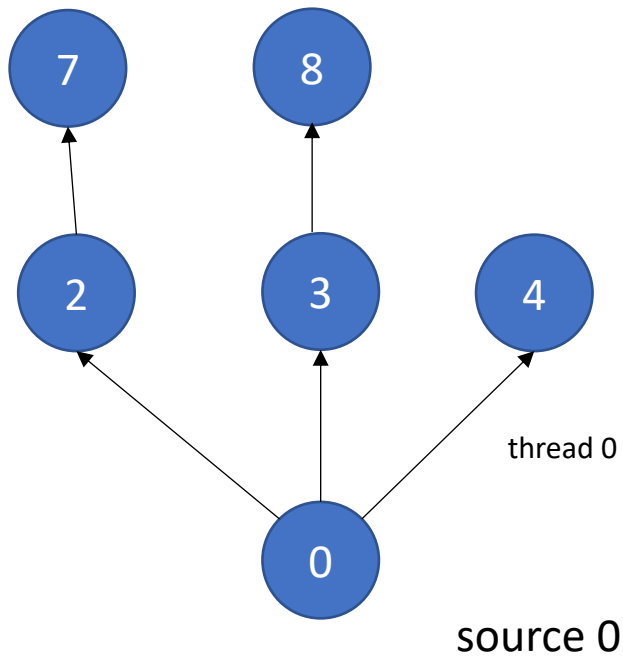
Input/Output Queues

- Example: Information flow in graph applications:



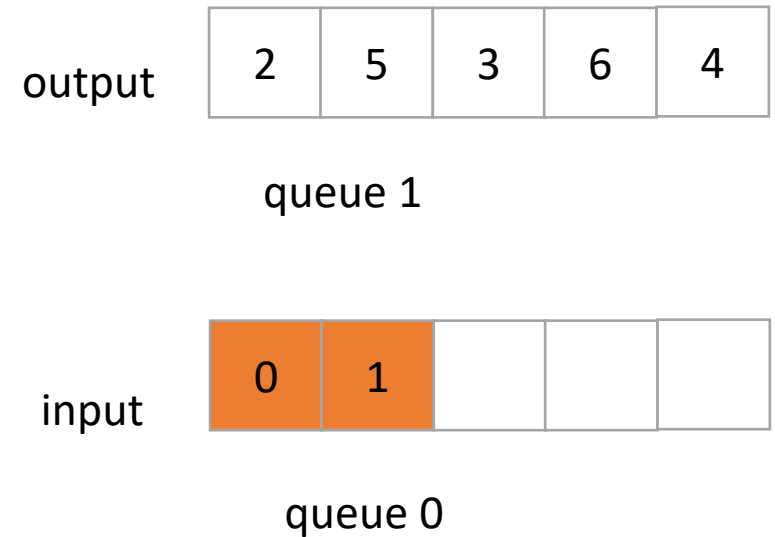
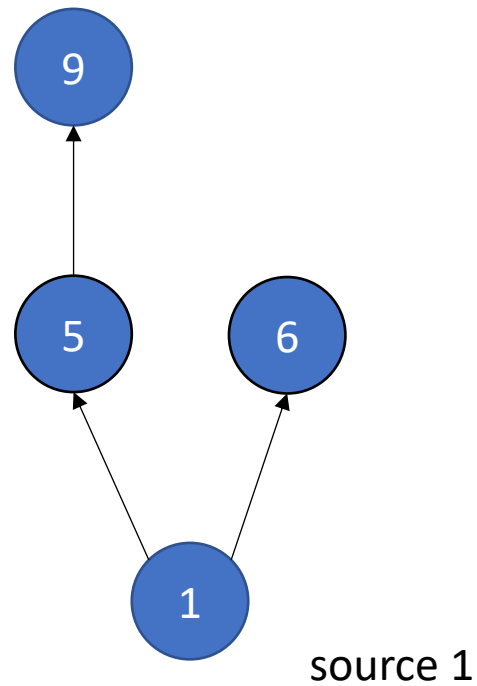
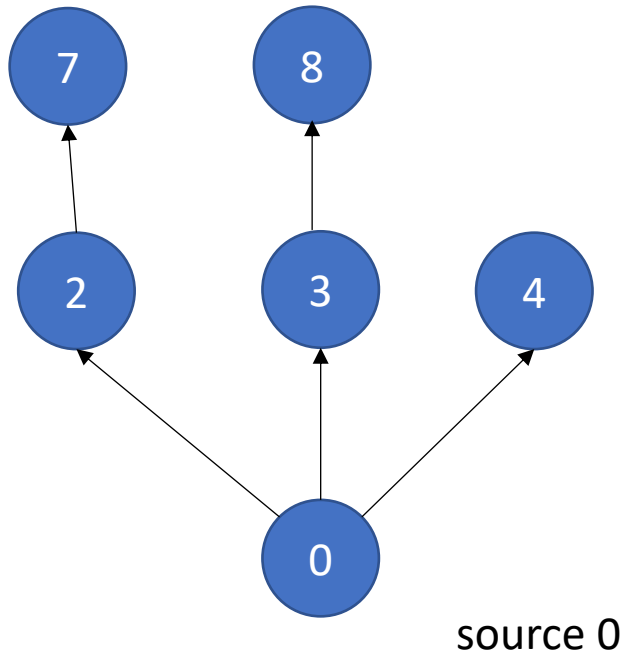
Input/Output Queues

- Example: Information flow in graph applications:



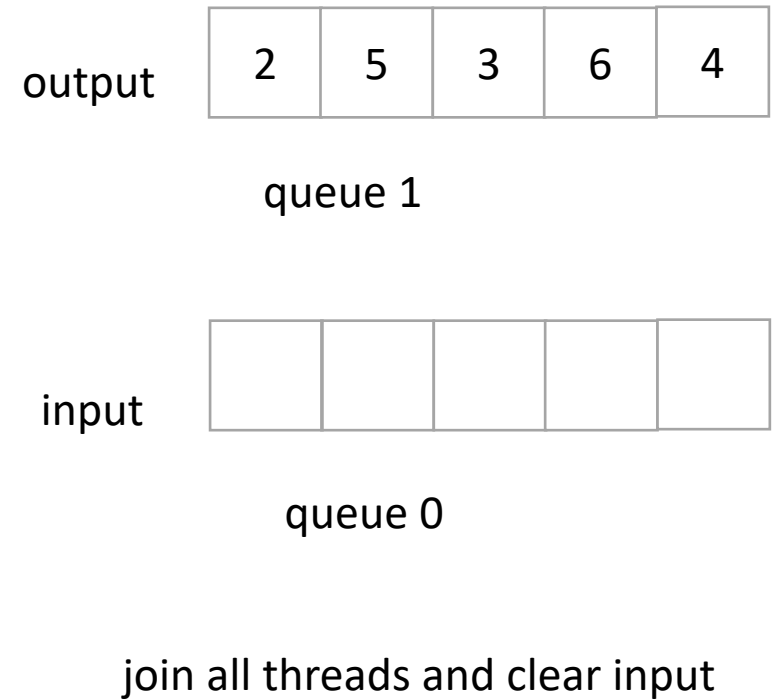
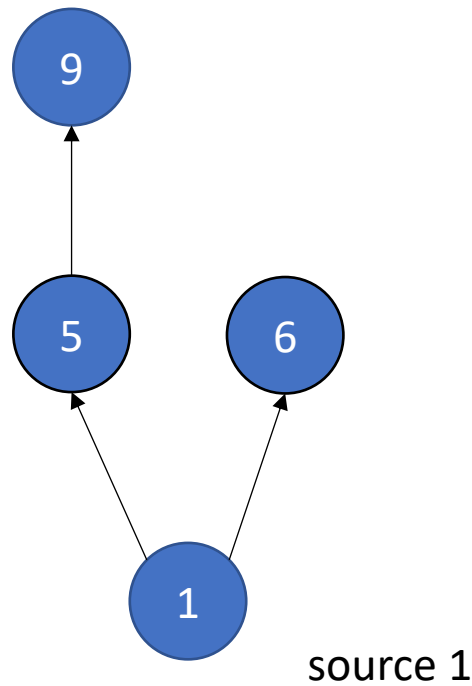
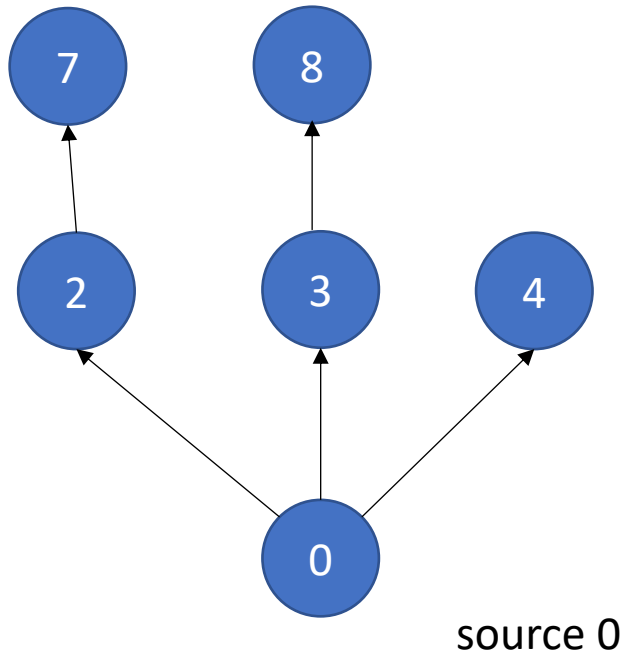
Input/Output Queues

- Example: Information flow in graph applications:



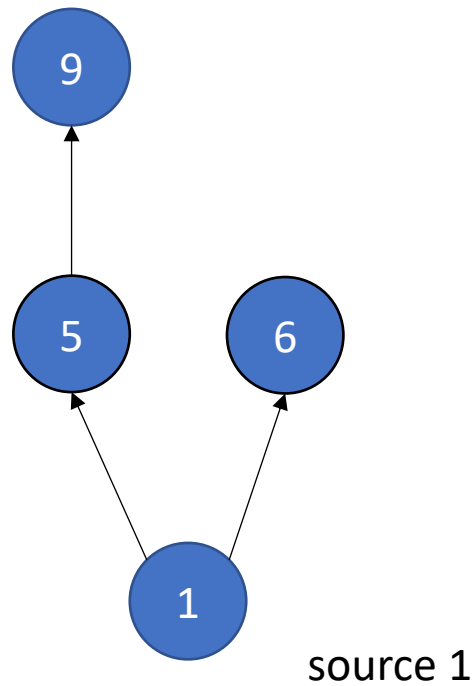
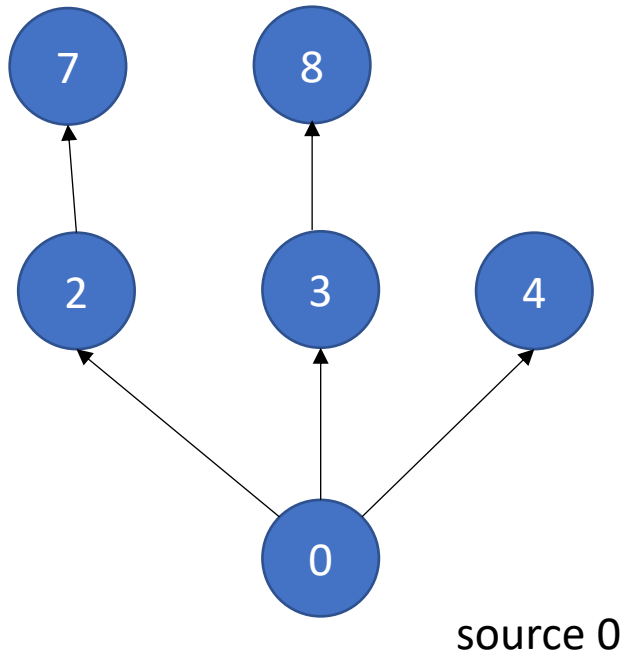
Input/Output Queues

- Example: Information flow in graph applications:



Input/Output Queues

- Example: Information flow in graph applications:



input



swap!

queue 1

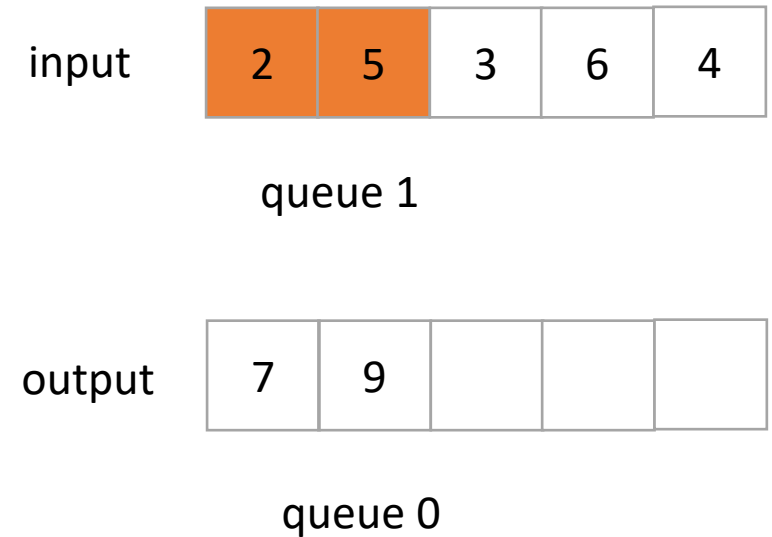
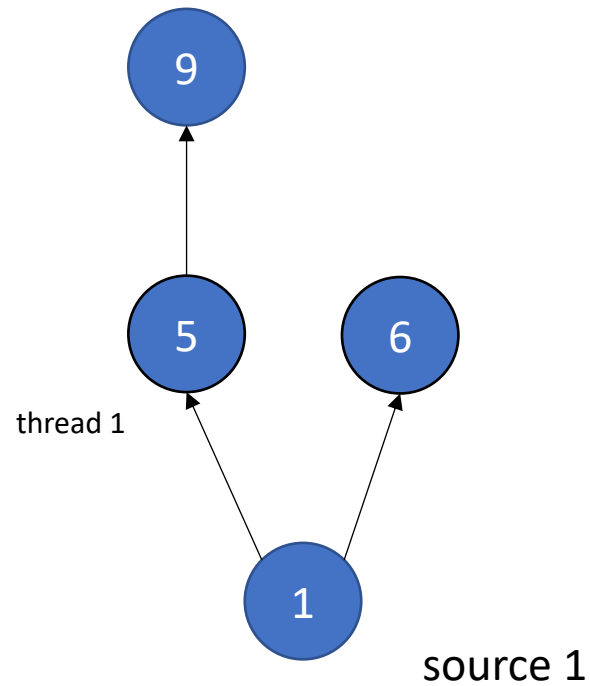
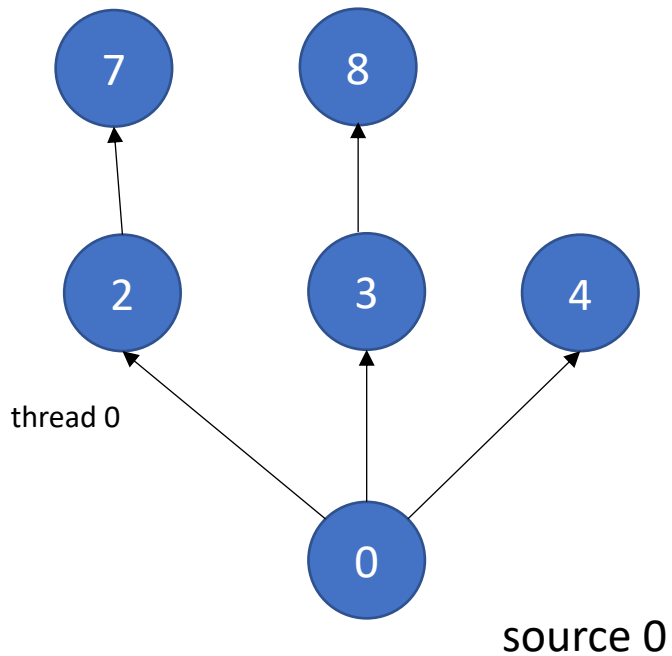
output



queue 0

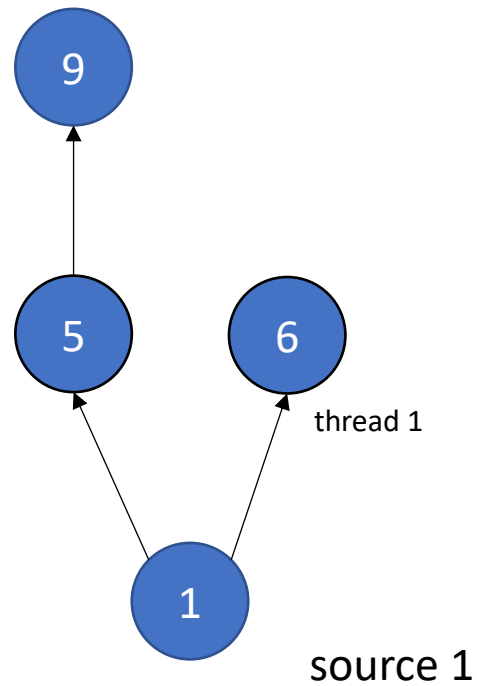
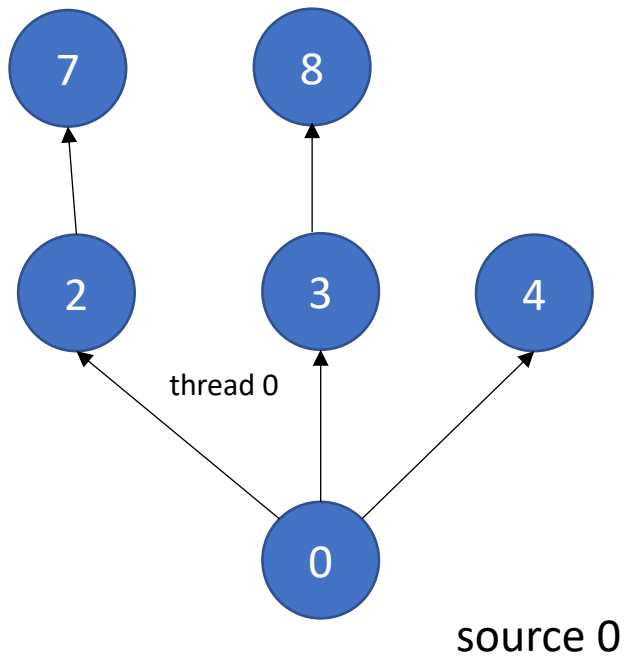
Input/Output Queues

- Example: Information flow in graph applications:



Input/Output Queues

- Example: Information flow in graph applications:



input



queue 1

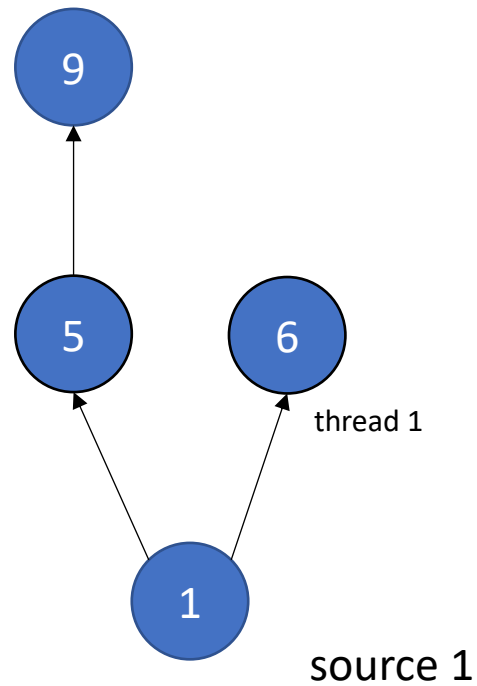
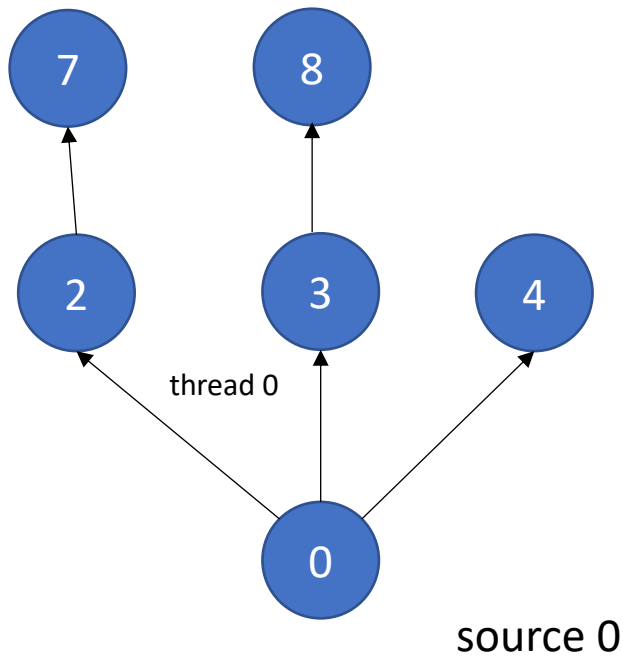
output



queue 0

Input/Output Queues

- Example: Information flow in graph applications:

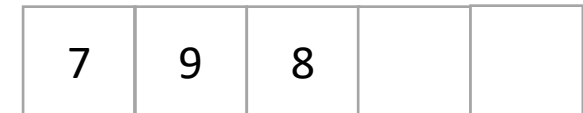


input



queue 1

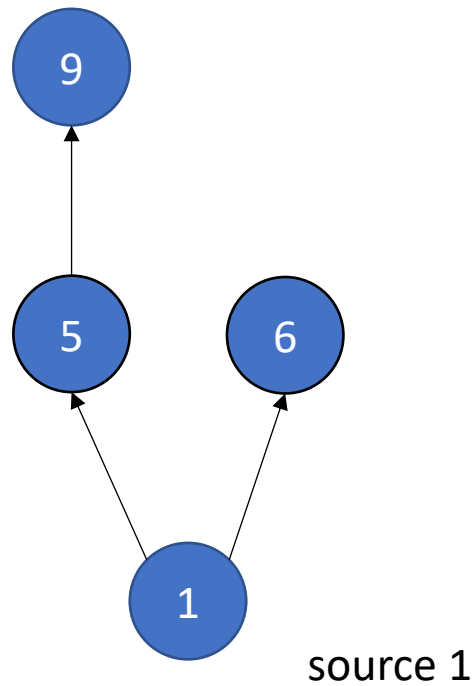
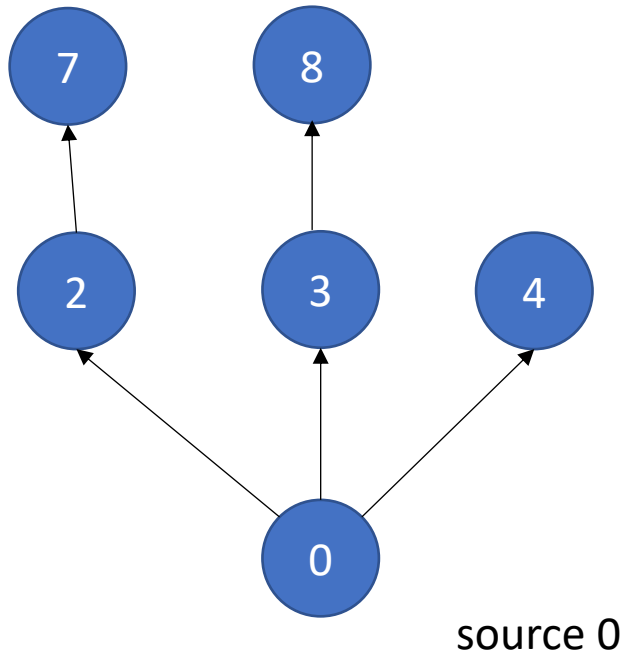
output



queue 0

Input/Output Queues

- Example: Information flow in graph applications:



input

2	5	3	6	4
---	---	---	---	---

queue 1

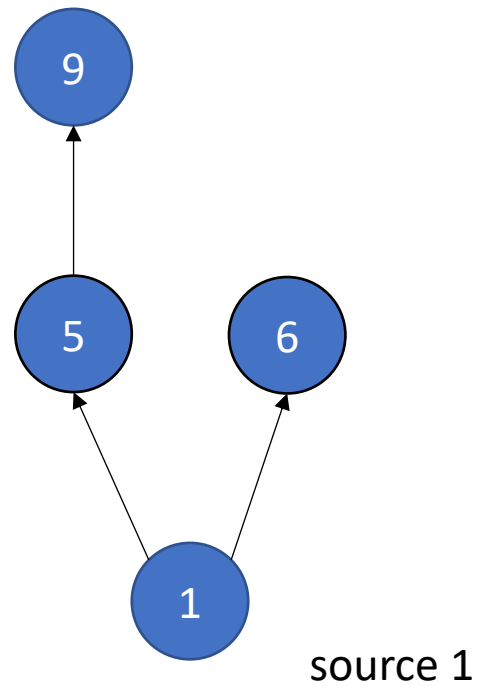
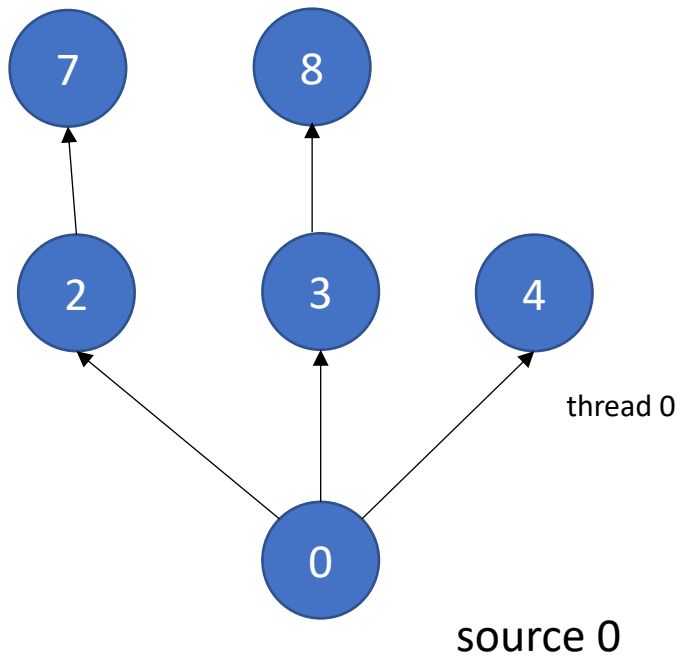
output

7	9	8		
---	---	---	--	--

queue 0

Input/Output Queues

- Example: Information flow in graph applications:

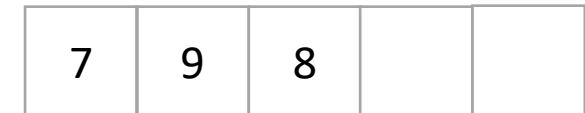


input



queue 1

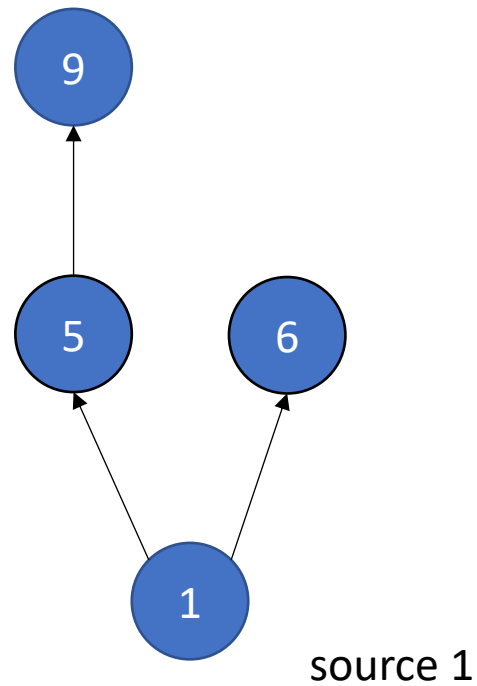
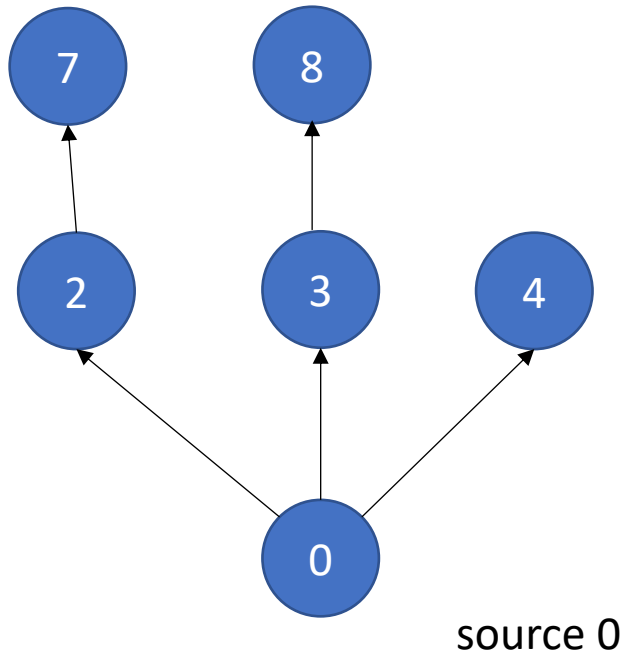
output



queue 0

Input/Output Queues

- Example: Information flow in graph applications:



input

2	5	3	6	4
---	---	---	---	---

queue 1

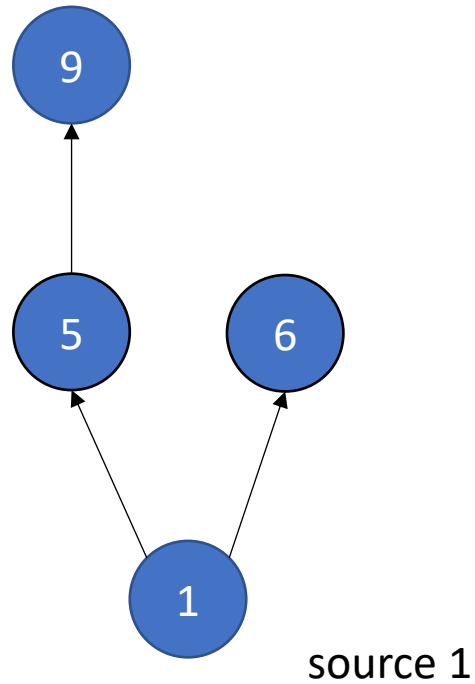
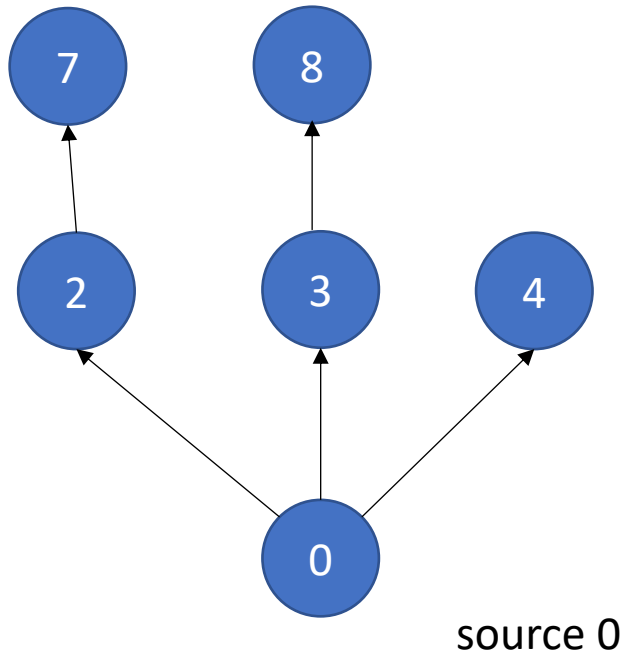
output

7	9	8		
---	---	---	--	--

queue 0

Input/Output Queues

- Example: Information flow in graph applications:



and so on...

output



queue 1

input



queue 0

Implementation

Implementation

Allocate a contiguous array



Pros:

?

Cons:

?

Implementation

Allocate a contiguous array



Pros:

+ fast!

+ we can use indexes instead of addresses

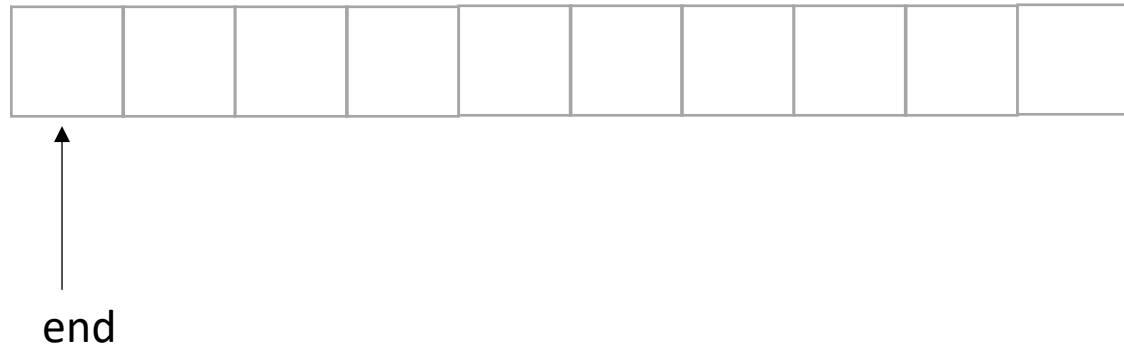
Cons:

- need to reason about overflow!

Note on terminology

- Head/tail - often used in queue implementations, but switches when we start doing circular buffers.
- Front/end - To avoid confusion, we will use front/end for input/output queues.

Implementation



Implementation



↑
end

What happens if a thread wants
to add an element?

Implementation



↑
end

What happens if a thread wants to add an element?

Think sequentially:

Implementation



↑
end

What happens if a thread wants to add an element?

Think sequentially:

*reserve a space - increment end

Implementation

reserved!



↑
end

What happens if a thread wants to add an element?

Think sequentially:

*reserve a space - increment end

Implementation

reserved!



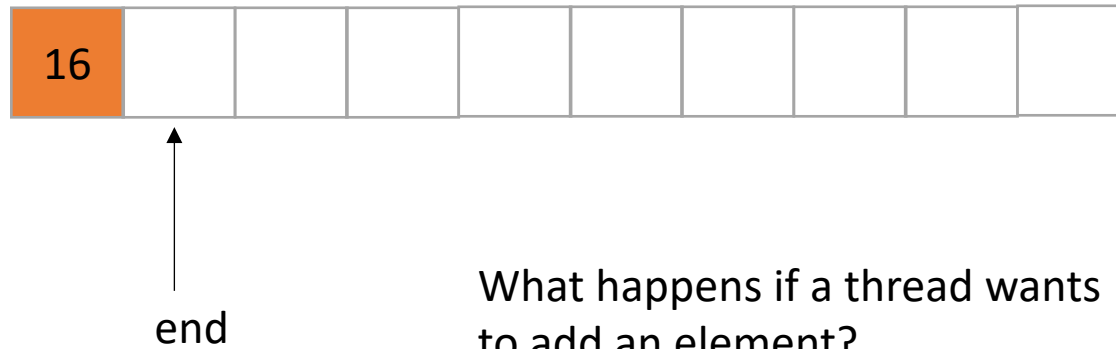
end

What happens if a thread wants to add an element?

Think sequentially:

- * reserve a space - increment end
- * add the element

Implementation

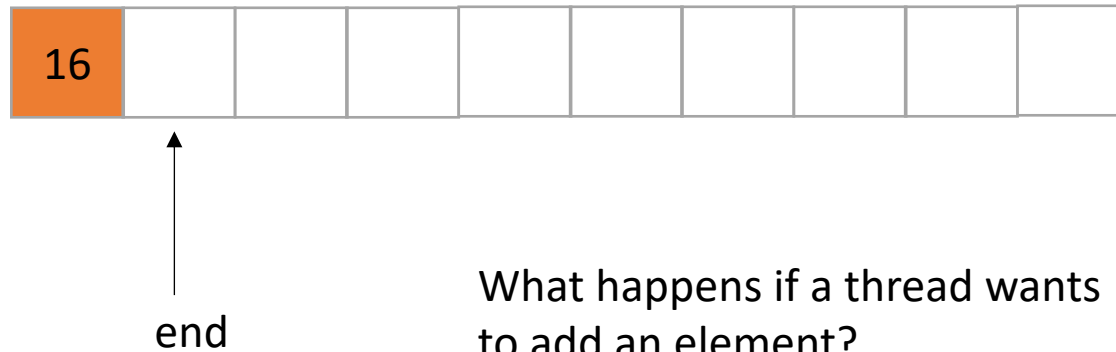


What happens if a thread wants to add an element?

Think sequentially:

- * reserve a space - increment end
- * add the element

Implementation



What happens if a thread wants to add an element?

Think sequentially:

- * reserve a space - increment end
- * add the element

done!

Implementation

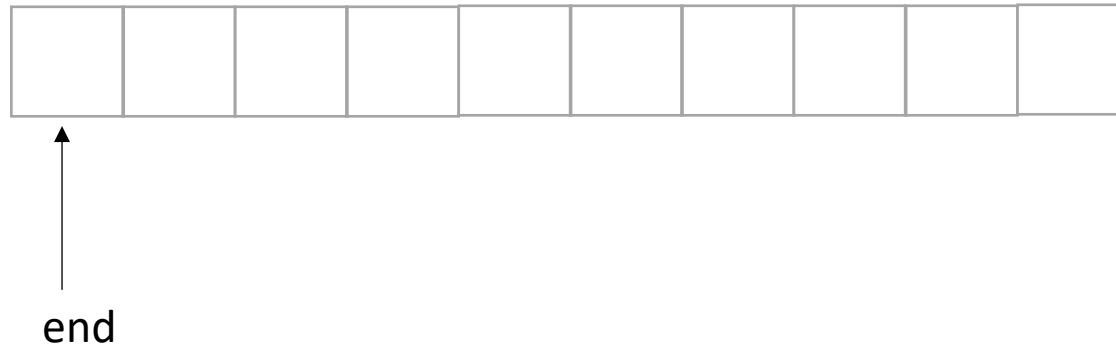


What happens if a thread wants to add an element?

Think concurrently:

*Two threads cannot reserve the same space!
We've seen this before*

Implementation

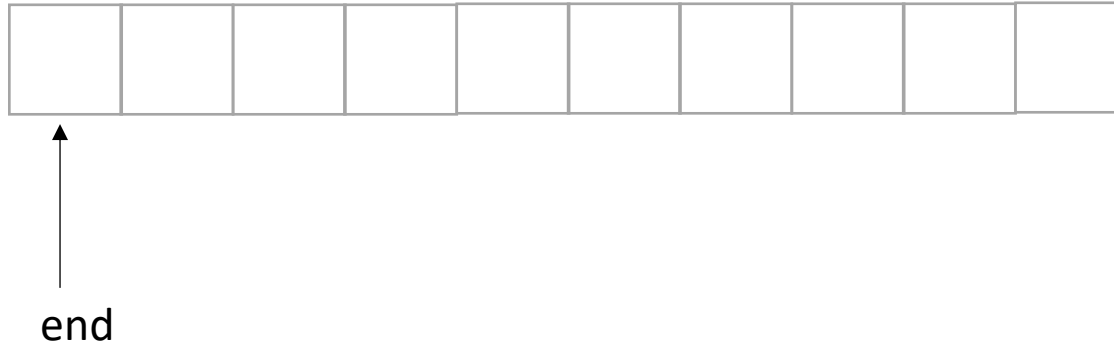


What happens if a thread wants to add an element?

Think concurrently:

```
reserved_index = atomic_fetch_add(&end, 1);
```

Implementation



Thread 0:
enq(6);

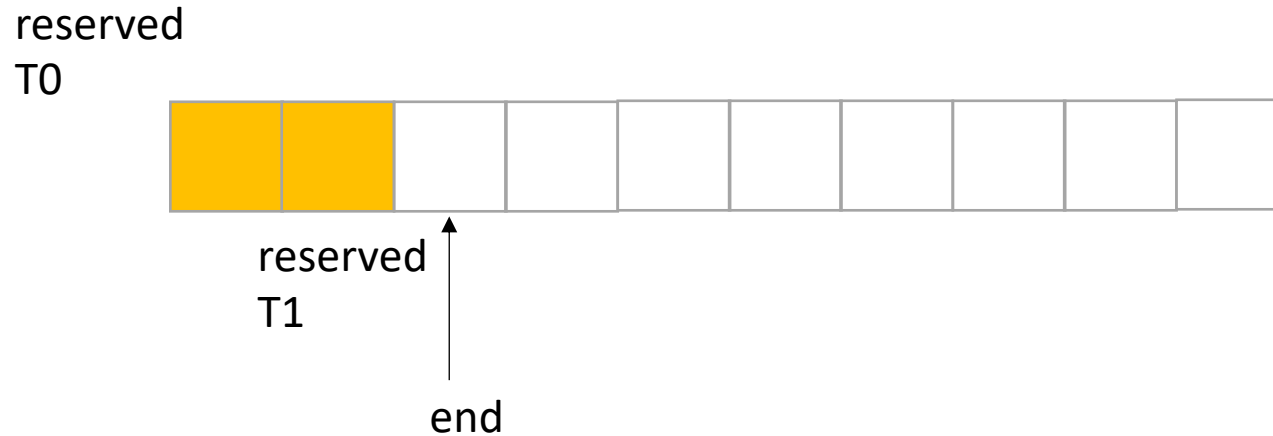
Thread 1:
enq(7);

What happens if a thread wants to add an element?

Think concurrently:

```
reserved_index = atomic_fetch_add(&end, 1);
```

Implementation



Thread 0:
enq(6);

Thread 1:
enq(7);

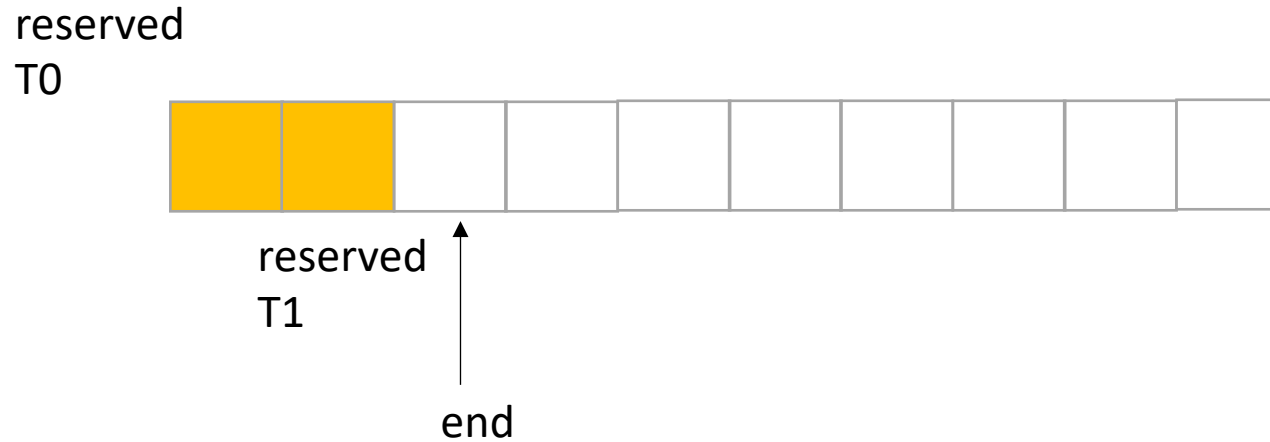
What happens if a thread wants to add an element?

Think concurrently:

```
reserved_index = atomic_fetch_add(&end, 1);
```

Implementation

*does it matter which order
threads add their data?*



Thread 0:
`enq(6);`

Thread 1:
`enq(7);`

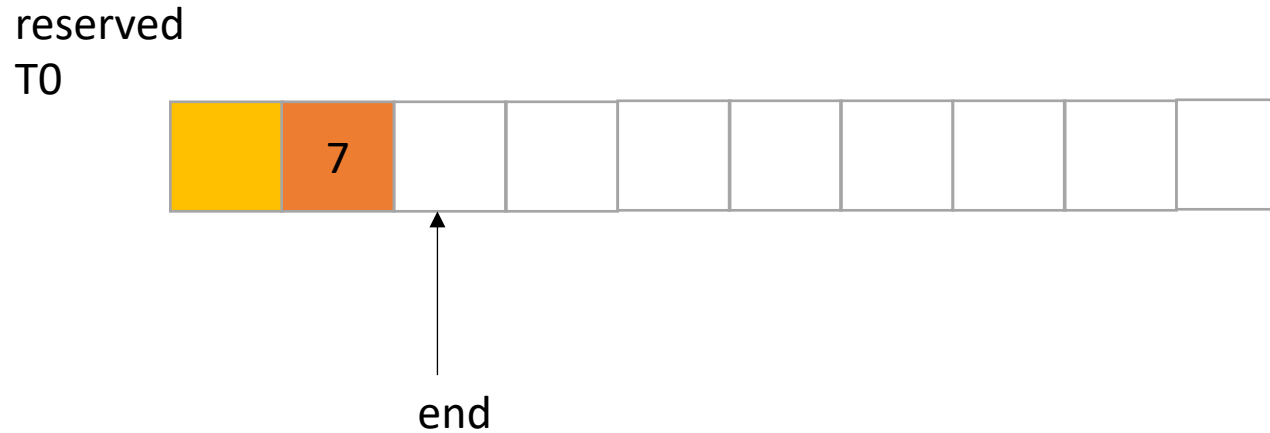
What happens if a thread wants
to add an element?

Think concurrently:

```
reserved_index = atomic_fetch_add(&end, 1);
```

Implementation

*does it matter which order
threads add their data?*



Thread 0:
`enq(6);`

Thread 1:
`enq(7);`

What happens if a thread wants
to add an element?

Think concurrently:

```
reserved_index = atomic_fetch_add(&end, 1);
```

Implementation

*does it matter which order
threads add their data? No!
Because there are no deqs!*

reserved
T0



end

Thread 0:
`enq(6);`

Thread 1:
`enq(7);`

What happens if a thread wants
to add an element?

Think concurrently:

```
reserved_index = atomic_fetch_add(&end, 1);
```

```
class InputOutputQueue {
    private:
        atomic_int end;
        int list[SIZE];

    public:
        InputOutputQueue() {
            end = 0;
        }

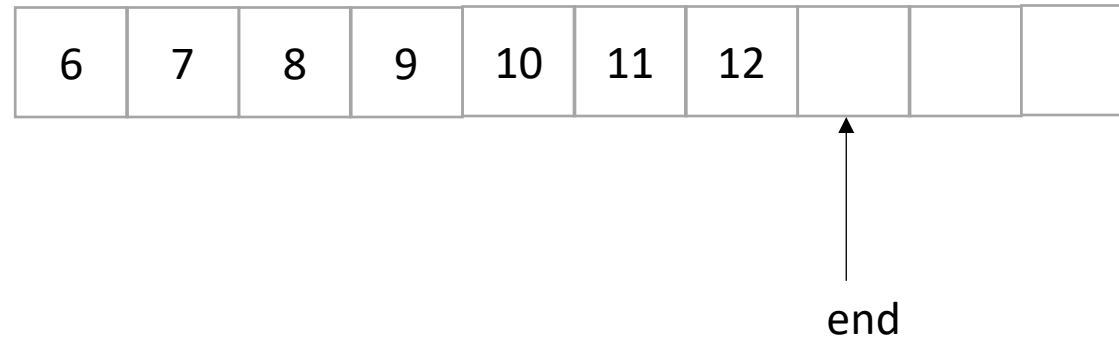
        void enq(int x) {
            int reserved_index = atomic_fetch_add(&end, 1);
            list[reserved_index] = x;
        }

        int size() {
            return end.load();
        }
}
```

How to protect against overflows?

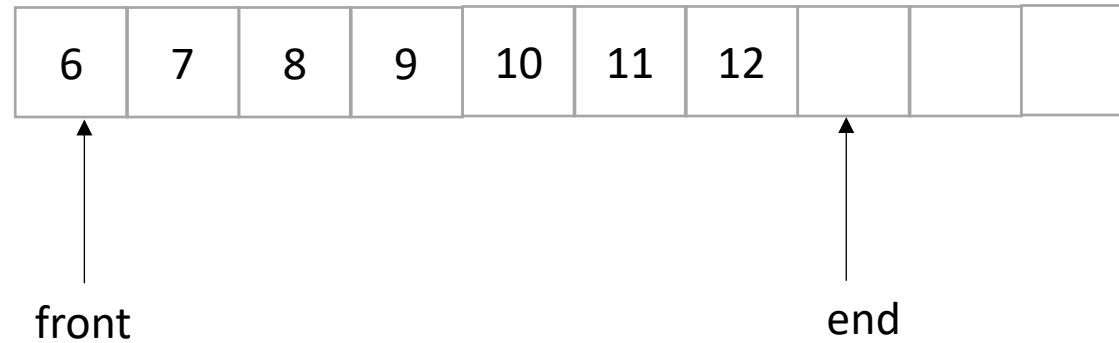
What about Input?

- Now we only do deqs



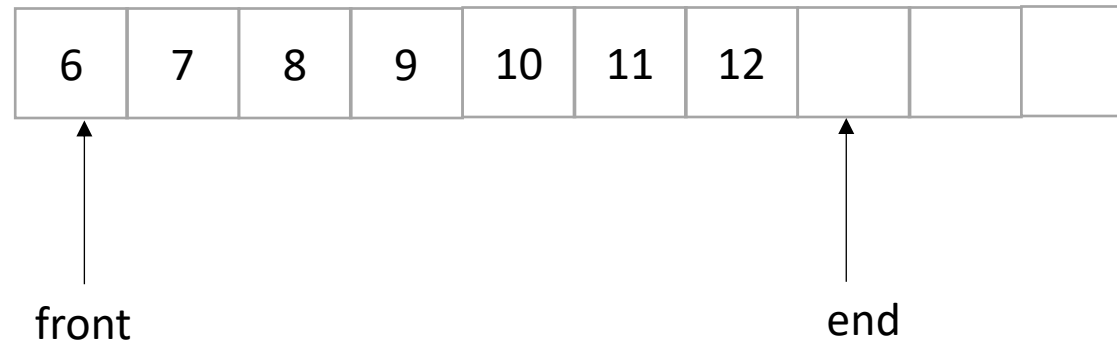
What about Input?

- Now we only do deqs



What about Input?

- Now we only do deqs



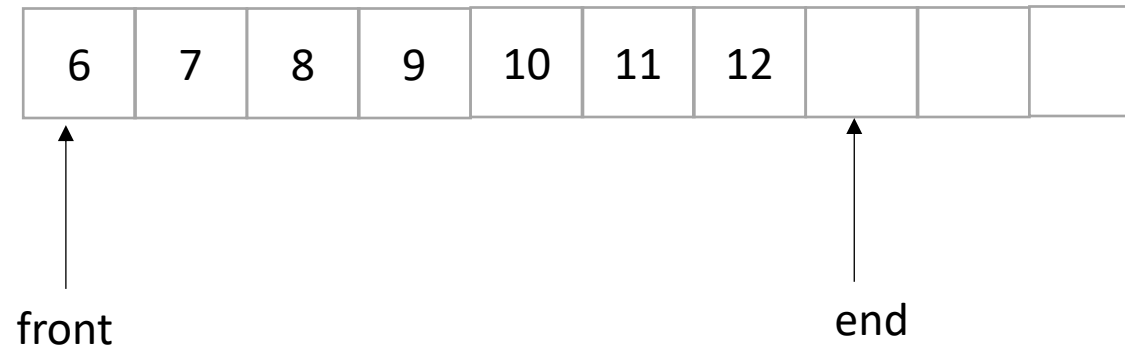
What happens if a thread wants to add an element?

Think concurrently:

```
data_index = atomic_fetch_add(&front, 1);
```

What about Input?

- Now we only do deqs



Thread 0:
deq();

Thread 1:
deq();

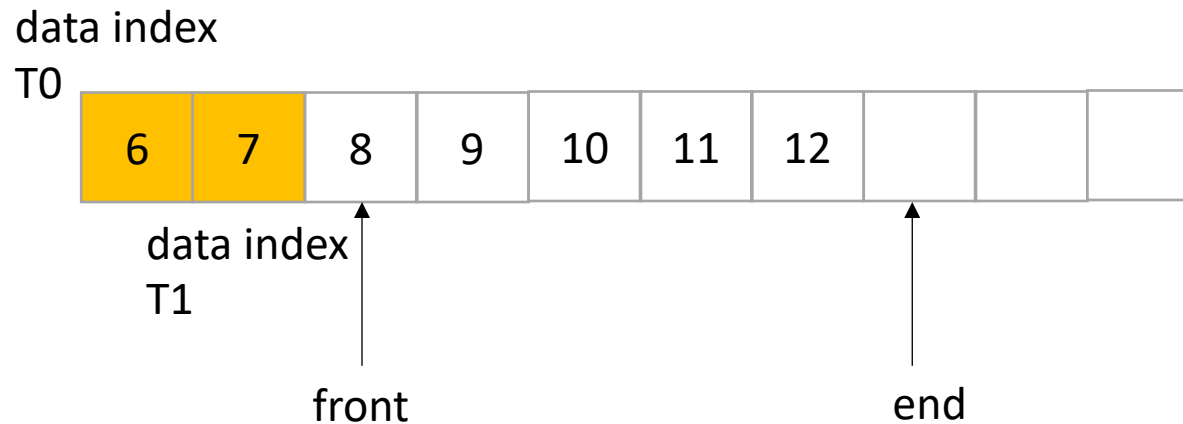
What happens if a thread wants to add an element?

Think concurrently:

```
data_index = atomic_fetch_add(&front, 1);
```

What about Input?

- Now we only do deqs



Thread 0:
deq();

Thread 1:
deq();

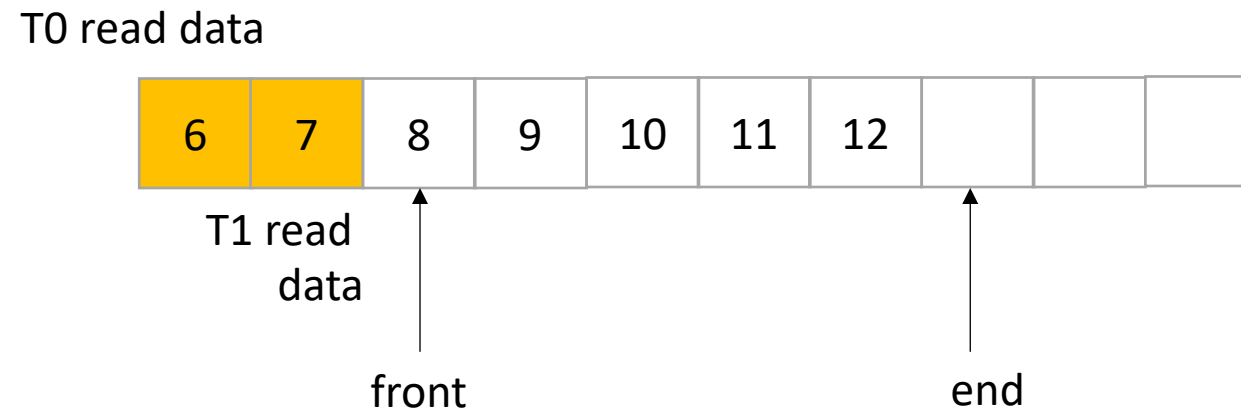
What happens if a thread wants to add an element?

Think concurrently:

```
data_index = atomic_fetch_add(&front, 1);
```

What about Input?

- Now we only do deqs



Thread 0:
`deq(); // reads 6`

Thread 1:
`deq(); // reads 7`

What happens if a thread wants to add an element?

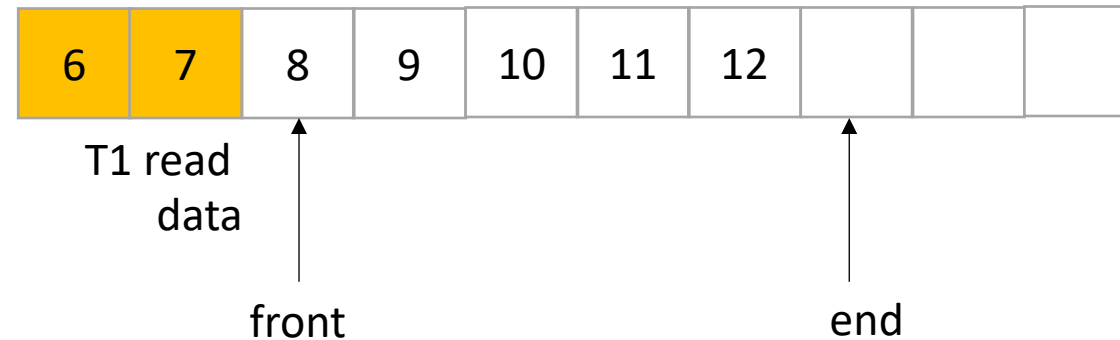
Think concurrently:

```
data_index = atomic_fetch_add(&front, 1);
```

What about Input?

- Now we only do deqs

T0 read data



T1 read data

Thread 0:
`deq(); // reads 6`

Thread 1:
`deq(); // reads 7`

What happens if a thread wants to add an element?

Think concurrently:

```
data_index = atomic_fetch_add(&front, 1);
```

How to implement a stack?

```
class InputOutputQueue {  
  private:  
    atomic_int front;  
    atomic_int end;  
    int list[SIZE];  
  
  public:  
    InputOutputQueue() {  
      front = end = 0;  
    }  
  
    void enq(int x) {  
      int reserved_index = atomic_fetch_add(&end, 1);  
      list[reserved_index] = x;  
    }  
  
    void deq() {  
      int reserved_index = atomic_fetch_add(&front, 1);  
      return list[reserved_index];  
    }  
  
    int size() {  
      return ??;  
    }  
}
```

```
class InputOutputQueue {
    private:
        atomic_int front;
        atomic_int end;
        int list[SIZE];

    public:
        InputOutputQueue() {
            front = end = 0;
        }

        void enq(int x) {
            int reserved_index = atomic_fetch_add(&end, 1);
            list[reserved_index] = x;
        }

        void deq() {
            int reserved_index = atomic_fetch_add(&front, 1);
            return list[reserved_index];
        }

        int size() {
            return ??;
        }
}
```

How about size?


```
class InputOutputQueue {
  private:
    atomic_int front;
    atomic_int end;
    int list[SIZE];

  public:
    InputOutputQueue() {
      front = end = 0;
    }

    void enq(int x) {
      int reserved_index = atomic_fetch_add(&end, 1);
      list[reserved_index] = x;
    }

    void deq() {
      int reserved_index = atomic_fetch_add(&front, 1);
      return list[reserved_index];
    }

    int size() {
      return end.load() - front.load();
    }
}
```

how about size?

how do we reset?

```
class InputOutputQueue {
  private:
    atomic_int front;
    atomic_int end;
    int list[SIZE];

  public:
    InputOutputQueue() {
      front = end = 0;
    }

    void enq(int x) {
      int reserved_index = atomic_fetch_add(&end, 1);
      list[reserved_index] = x;
    }

    void deq() {
      int reserved_index = atomic_fetch_add(&front, 1);
      return list[reserved_index];
    }

    int size() {
      return end.load() - front.load();
    }
}
```

how about size?

how do we reset?
Reset front and end

```

class InputOutputQueue {
    private:
        atomic_int front;
        atomic_int end;
        int list[SIZE];

    public:
        InputOutputQueue() {
            front = end = 0;
        }

        void enq(int x) {
            int reserved_index = atomic_fetch_add(&end, 1);
            list[reserved_index] = x;
        }

        void deq() {
            int reserved_index = atomic_fetch_add(&front, 1);
            return list[reserved_index];
        }

        int size() {
            return end.load() - front.load();
        }
}

```

how about size?

how do we reset?
Reset front and end

does the list need
to be atomic?

Schedule

- Producer Consumer queues
 - **Synchronous**
 - Circular buffer

Producer Consumer Queues

- 1 enq, 1 deq
 - enq'er cannot deq
 - deq'er cannot enq
- Example: printf:
 - your program equeues values to print
 - the terminal process dequeues values and prints them

Synchronous Producer Consumer Queues

- First implementation:
 - Synchronous
 - Slow
 - Good for debugging

Synchronous Producer Consumer Queues

- First implementation:
 - Synchronous
 - Slow
 - Good for debugging
- enq does not return until value is deq'ed

Synchronous Producer Consumer Queues

Producer Thread
enq(7);



Consumer Thread
deq();

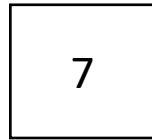
Synchronous Producer Consumer Queues

Producer Thread

`enq(7);`



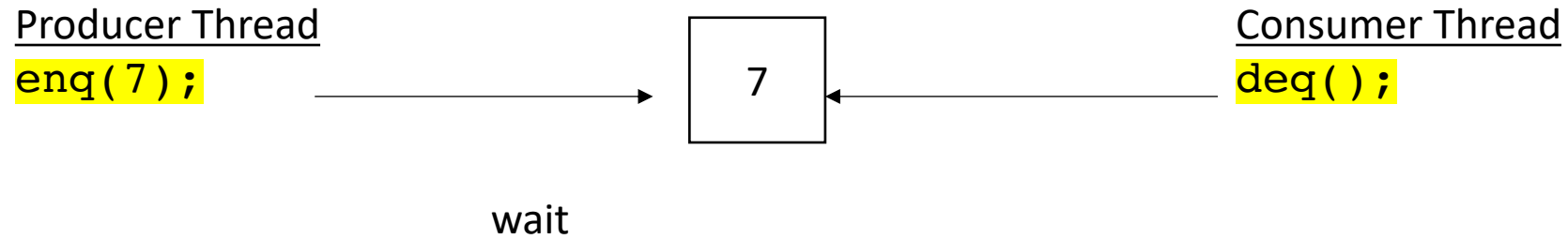
wait



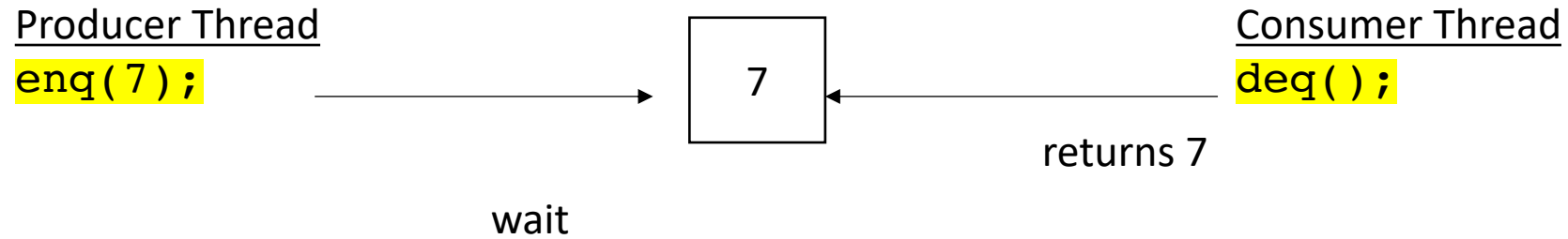
Consumer Thread

`deq();`

Synchronous Producer Consumer Queues



Synchronous Producer Consumer Queues



Synchronous Producer Consumer Queues

Producer Thread
enq(7);



Consumer Thread
deq();

both can continue

Synchronous Producer Consumer Queues

Producer Thread

```
sleep();  
enq(7);
```



Consumer Thread

```
deq();
```

Synchronous Producer Consumer Queues

Producer Thread

`sleep();`

`enq(7);`

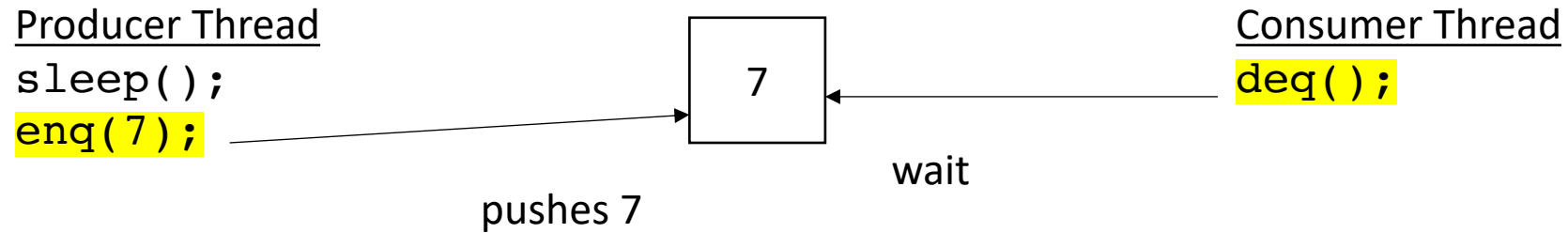


Consumer Thread

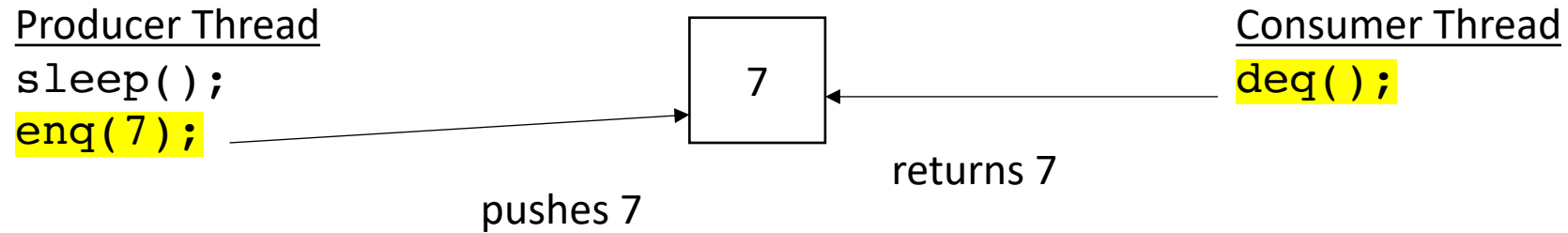
`deq();`

←
wait

Synchronous Producer Consumer Queues



Synchronous Producer Consumer Queues



They both can continue

Synchronous Producer Consumer Queues

Producer Thread
enq(7);



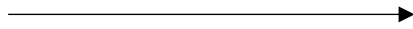
Consumer Thread
deq();

Synchronous Producer Consumer Queues

Producer Thread
enq(7);



Consumer Thread
deq();



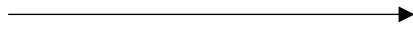
can the consumer just read?

Synchronous Producer Consumer Queues

Producer Thread
enq(7);



Consumer Thread
deq();



*can the consumer just read?
Needs to wait for a value to appear*

Synchronous Producer Consumer Queues

Producer Thread
enq(7);



flag

Consumer Thread
deq();

*can the consumer just read?
Needs to wait for a value to appear*

spin waiting for the flag to turn green

Synchronous Producer Consumer Queues

Producer Thread
enq(7);



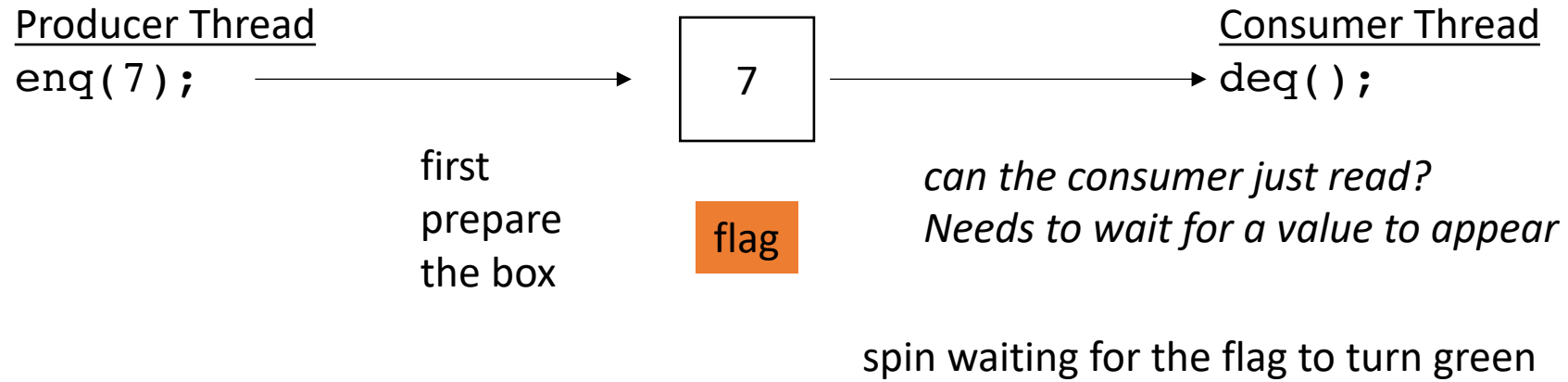
Consumer Thread
deq();

flag

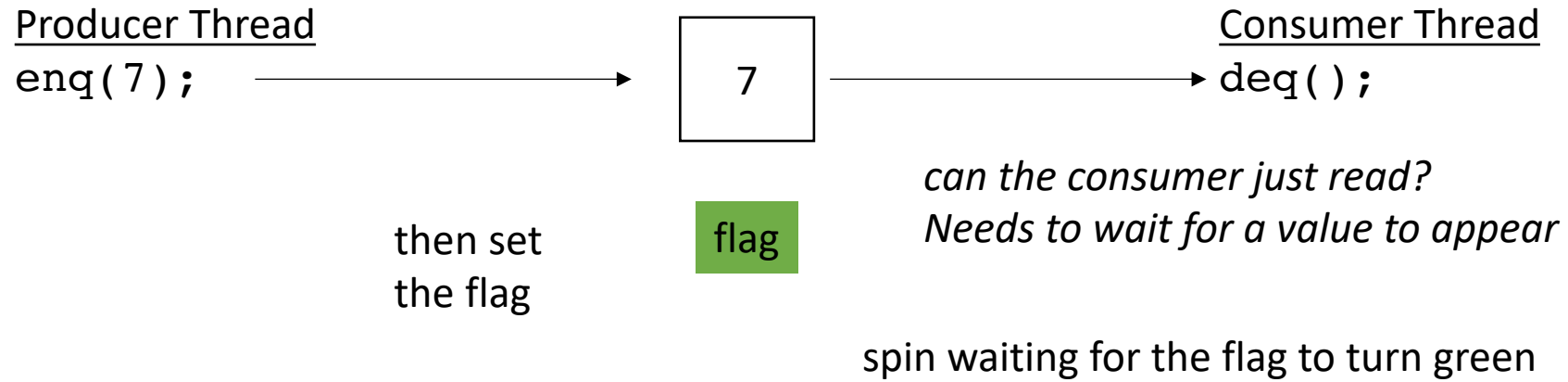
*can the consumer just read?
Needs to wait for a value to appear*

spin waiting for the flag to turn green

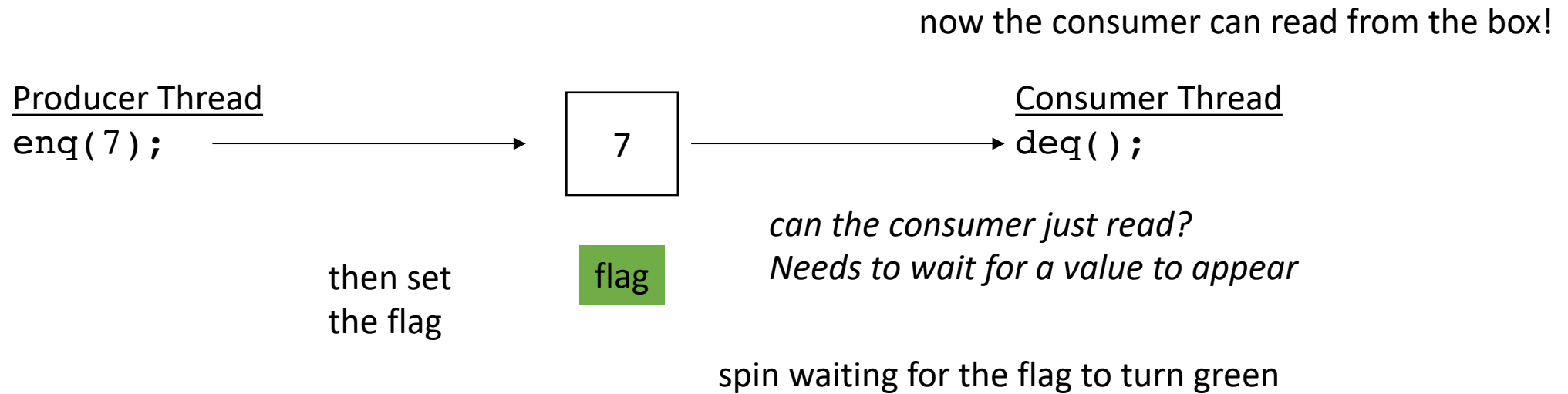
Synchronous Producer Consumer Queues



Synchronous Producer Consumer Queues



Synchronous Producer Consumer Queues



Synchronous Producer Consumer Queues

Producer Thread
enq(7);



flag

Consumer Thread
deq();

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
    }
}
```

Synchronous Producer Consumer Queues

Producer Thread
enq(7);



flag

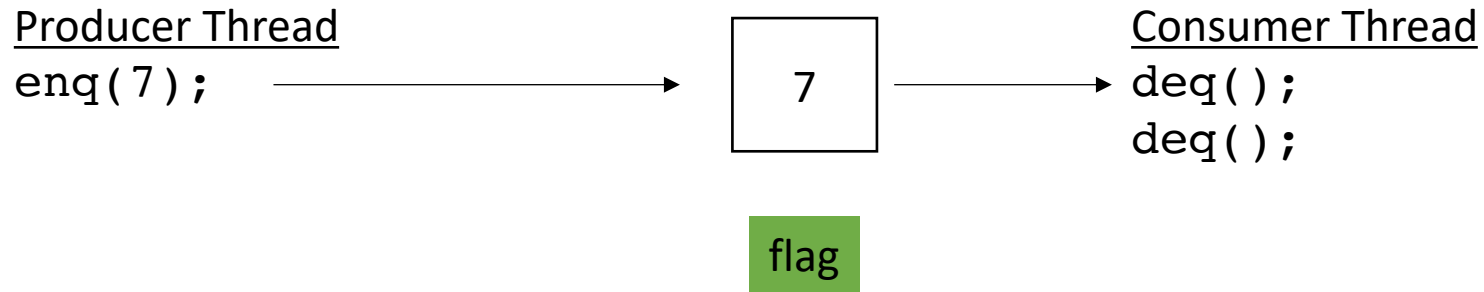
Consumer Thread
deq();
deq();

what happens
when there are
two deqs?

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

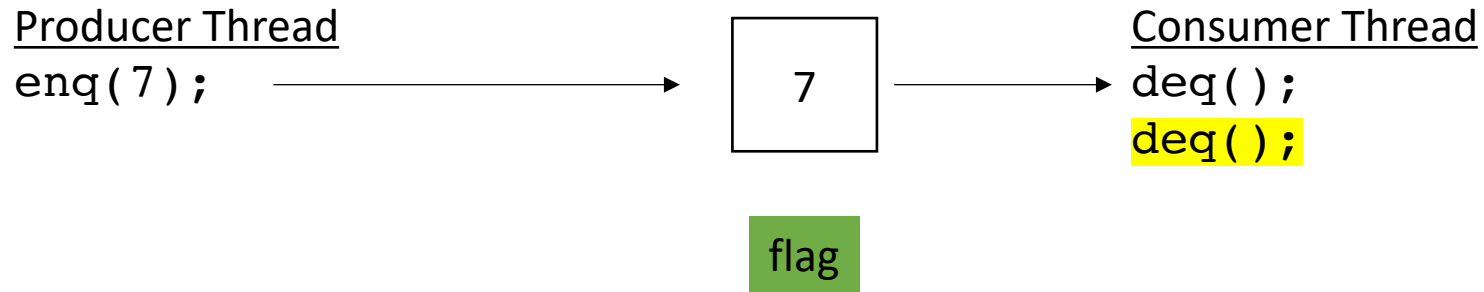
public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
    }
}
```

Synchronous Producer Consumer Queues



```
class SyncQueue {  
    private:  
        atomic_int box;  
        atomic_bool flag;  
  
    public:  
        void enq(int x) {  
            // put value in box  
            // set flag  
        }  
        void deq() {  
            // wait for flag to be set  
            // read from the box  
        }  
}
```

Synchronous Producer Consumer Queues

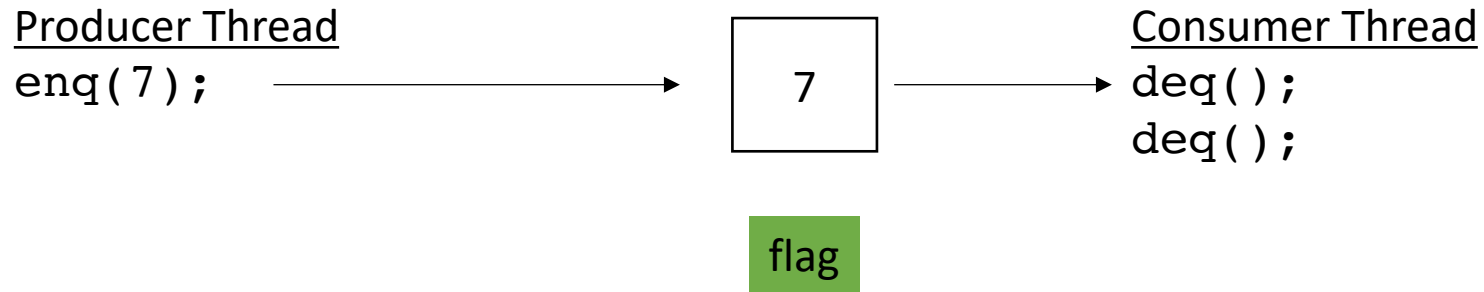


```
class SyncQueue {  
    private:  
        atomic_int box;  
        atomic_bool flag;  
  
    public:  
        void enq(int x) {  
            // put value in box  
            // set flag  
        }  
        void deq() {  
            // wait for flag to be set  
            // read from the box  
        }  
}
```

what happens in the
next deq?

How to fix?

Synchronous Producer Consumer Queues



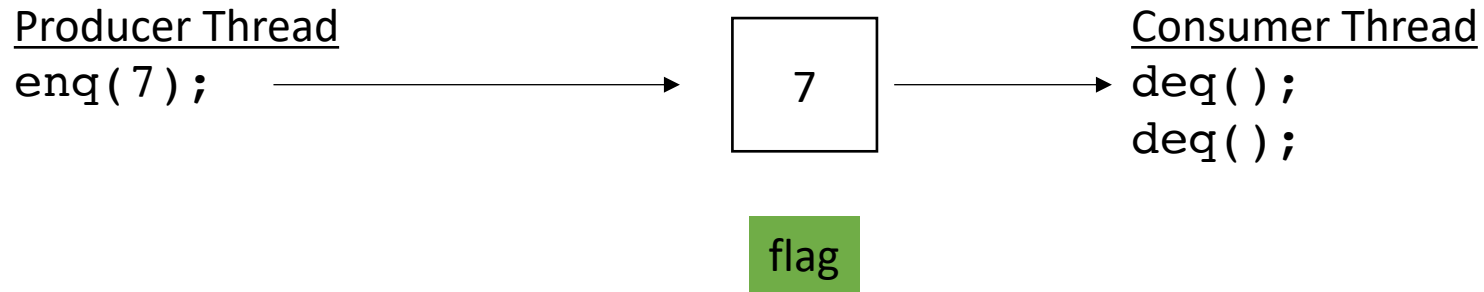
what happens in the
next deq?

How to fix?

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

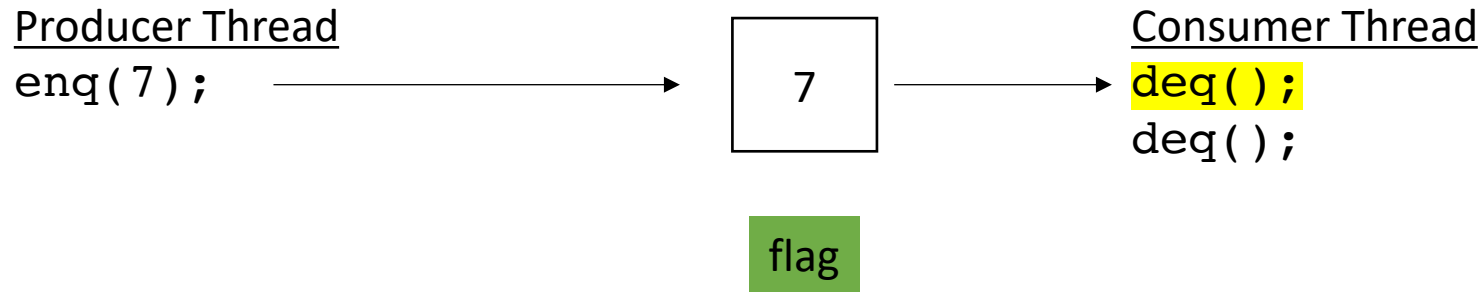
Synchronous Producer Consumer Queues



```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

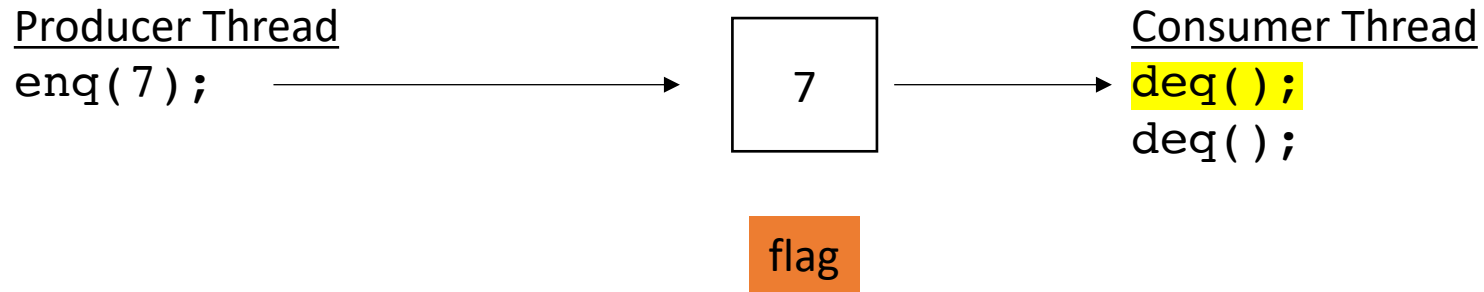
Synchronous Producer Consumer Queues



```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

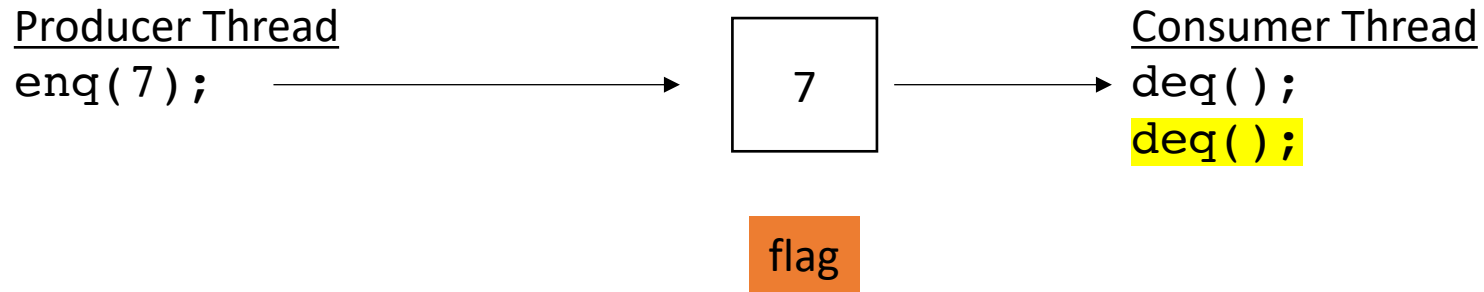
Synchronous Producer Consumer Queues



```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```


Synchronous Producer Consumer Queues



```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

waiting like we are
supposed to

Synchronous Producer Consumer Queues

Producer Thread

```
enq(7);  
enq(8);
```

extra enq

reset (now with extra enq)



flag

Consumer Thread

```
deq();  
deq();
```

```
class SyncQueue {  
    private:  
        atomic_int box;  
        atomic_bool flag;  
  
    public:  
        void enq(int x) {  
            // put value in box  
            // set flag  
        }  
        void deq() {  
            // wait for flag to be set  
            // read from the box  
            // reset flag  
        }  
}
```

Synchronous Producer Consumer Queues

Producer Thread

enq(7);

enq(8);

7

flag

Consumer Thread

deq();

deq();

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

Synchronous Producer Consumer Queues

Producer Thread

enq(7);

enq(8);

7

flag

Consumer Thread

deq();

deq();

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

Synchronous Producer Consumer Queues

Producer Thread

enq(7);

enq(8);

8

flag

Consumer Thread

deq();

deq();

7 was dropped!

how to fix?

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

Synchronous Producer Consumer Queues

Producer Thread

enq(7);

enq(8);

8

flag

Consumer Thread

deq();

deq();

7 was dropped!

how to fix?

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
        // wait for flag to be reset
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

Synchronous Producer Consumer Queues

reset

Producer Thread

```
enq(7);  
enq(8);
```



flag

Consumer Thread

```
deq();  
deq();
```

```
class SyncQueue {  
    private:  
        atomic_int box;  
        atomic_bool flag;  
  
    public:  
        void enq(int x) {  
            // put value in box  
            // set flag  
            // wait for flag to be reset  
        }  
        void deq() {  
            // wait for flag to be set  
            // read from the box  
            // reset flag  
        }  
}
```

Synchronous Producer Consumer Queues

Producer Thread

enq(7);

enq(8);

7

flag

Consumer Thread

deq();

deq();

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
        // wait for flag to be reset
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```


Synchronous Producer Consumer Queues

Producer Thread

`enq(7);`

`enq(8);`

7

flag

Consumer Thread

`deq();`

`deq();`

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
        // wait for flag to be reset
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

Synchronous Producer Consumer Queues

Producer Thread

`enq(7);`

`enq(8);`

7

flag

Consumer Thread

`deq();`

`deq();`

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

public:
    void enq(int x) {
        // put value in box
        // set flag
        // wait for flag to be reset
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

Synchronous Producer Consumer Queues

Producer Thread

enq(7);

enq(8);

7

flag

Consumer Thread

deq();

deq();

```
class SyncQueue {
private:
    atomic_int box;
    atomic_bool flag;

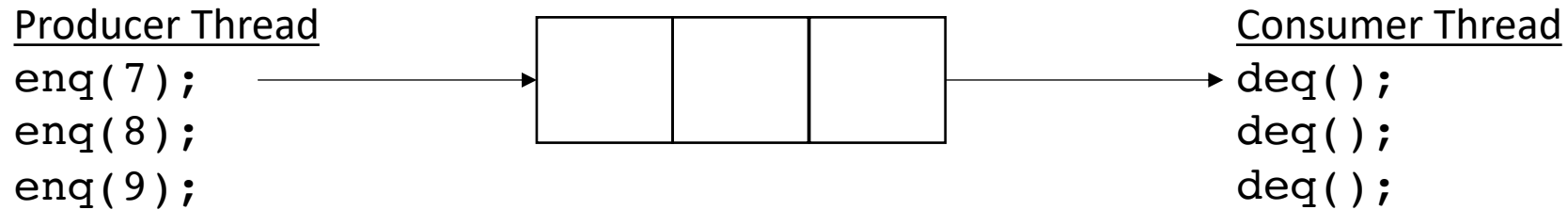
public:
    void enq(int x) {
        // put value in box
        // set flag
        // wait for flag to be reset
    }
    void deq() {
        // wait for flag to be set
        // read from the box
        // reset flag
    }
}
```

Schedule

- Producer Consumer Queues
 - Synchronous
 - Circular buffer

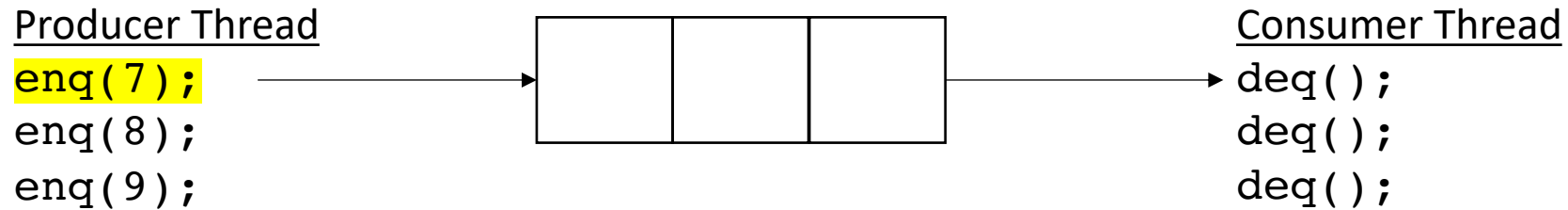
Producer Consumer Queues

- Asynchronous:



Producer Consumer Queues

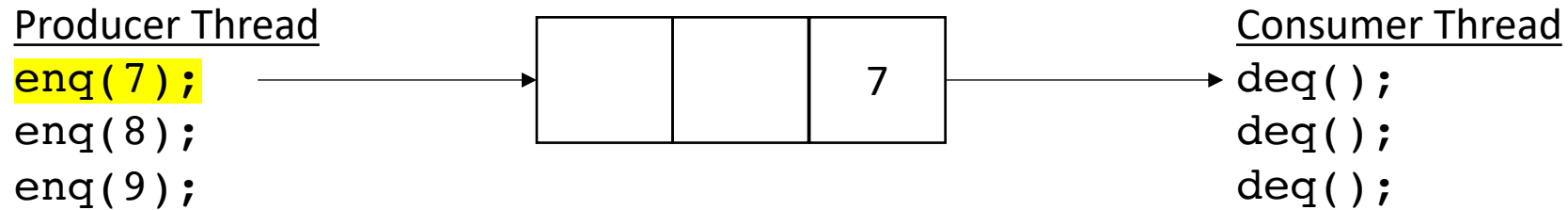
- Asynchronous:



no waiting for producer (while there is room)

Producer Consumer Queues

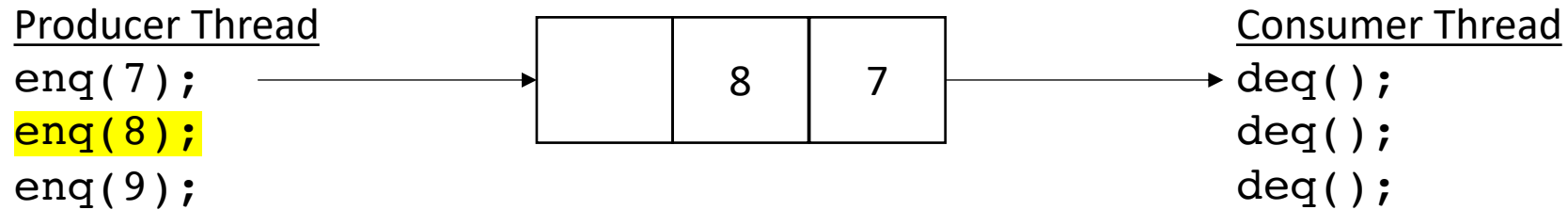
- Asynchronous:



no waiting for producer (while there is room)

Producer Consumer Queues

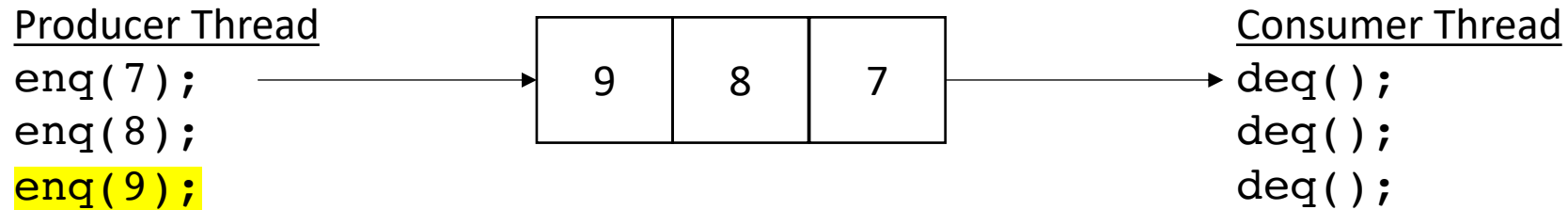
- Asynchronous:



no waiting for producer (while there is room)

Producer Consumer Queues

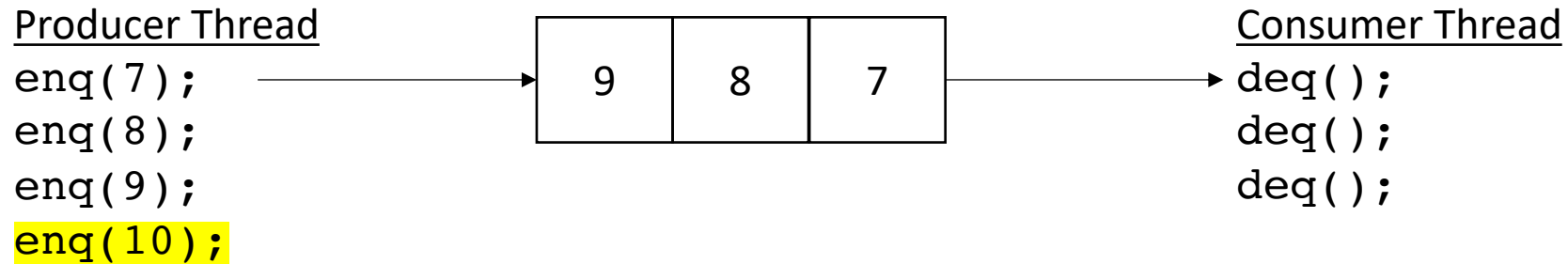
- Asynchronous:



no waiting for producer (while there is room)

Producer Consumer Queues

- Asynchronous:

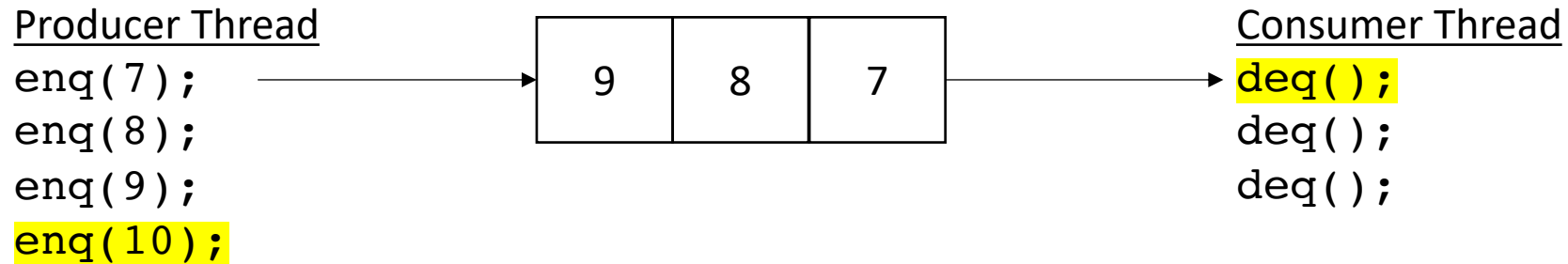


no waiting for producer (while there is room)

when there is no room, the queue will wait

Producer Consumer Queues

- Asynchronous:



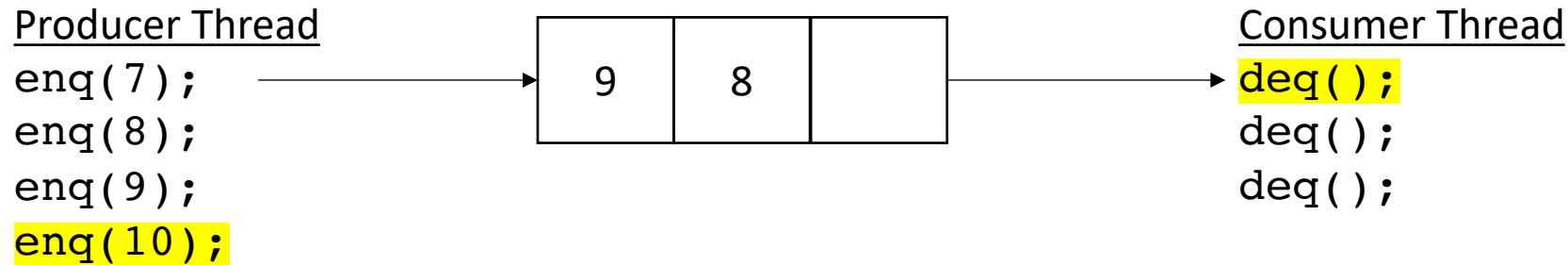
no waiting for producer (while there is room)

returns 7

when there is no room, the queue will wait

Producer Consumer Queues

- Asynchronous:



no waiting for producer (while there is room)

returns 7

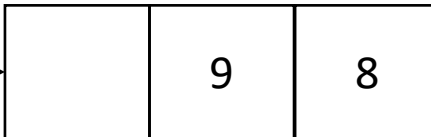
when there is no room, the queue will wait

Producer Consumer Queues

- Asynchronous:

Producer Thread

```
enq(7);  
enq(8);  
enq(9);  
enq(10);
```



Consumer Thread

```
deq();  
deq();  
deq();
```

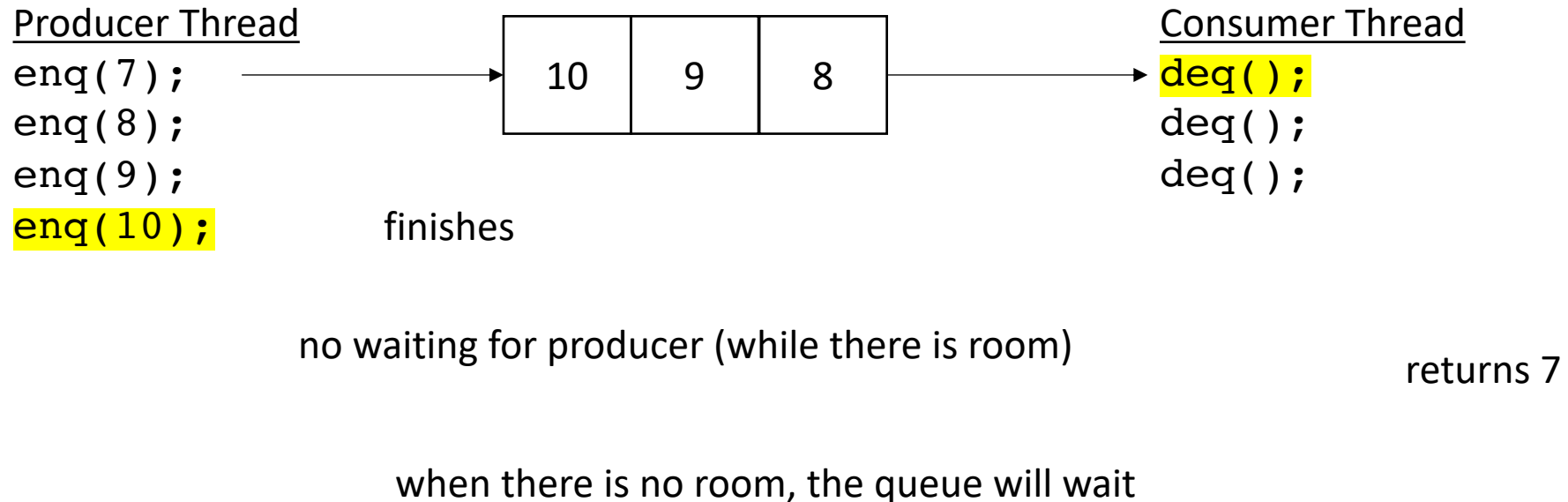
no waiting for producer (while there is room)

returns 7

when there is no room, the queue will wait

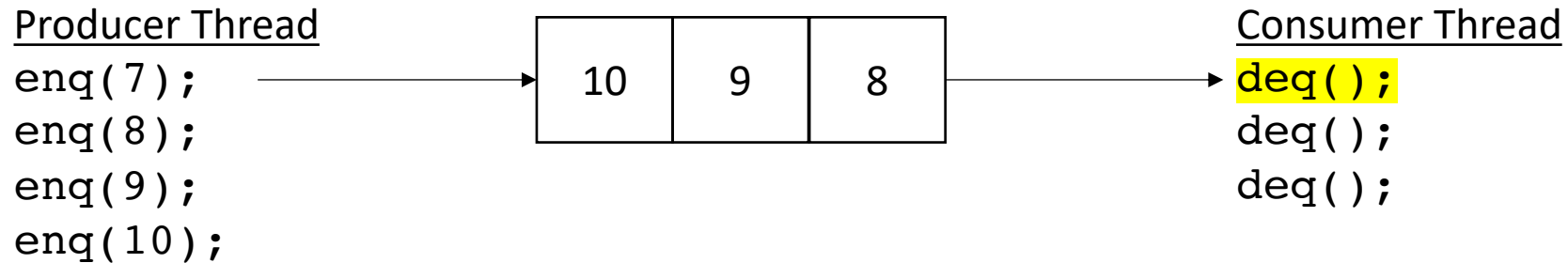
Producer Consumer Queues

- Asynchronous:



Producer Consumer Queues

- Asynchronous:



no waiting for producer (while there is room)

returns 7

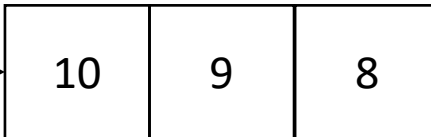
when there is no room, the queue will wait

Producer Consumer Queues

- Asynchronous:

Producer Thread

```
enq(7);  
enq(8);  
enq(9);  
enq(10);
```



Consumer Thread

```
deq();  
deq();  
deq();
```

no waiting for producer (while there is room)

returns 8

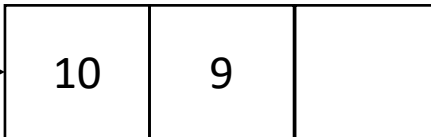
when there is no room, the queue will wait

Producer Consumer Queues

- Asynchronous:

Producer Thread

```
enq(7);  
enq(8);  
enq(9);  
enq(10);
```



Consumer Thread

```
deq();  
deq();  
deq();
```

no waiting for producer (while there is room)

returns 8

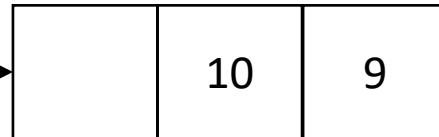
when there is no room, the queue will wait

Producer Consumer Queues

- Asynchronous:

Producer Thread

```
enq(7);  
enq(8);  
enq(9);  
enq(10);
```



Consumer Thread

```
deq();  
deq();  
deq();
```

no waiting for producer (while there is room)

returns 8

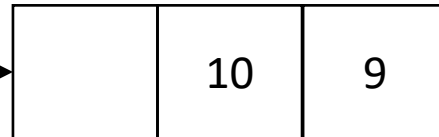
when there is no room, the queue will wait

Producer Consumer Queues

- Asynchronous:

Producer Thread

```
enq(7);  
enq(8);  
enq(9);  
enq(10);
```



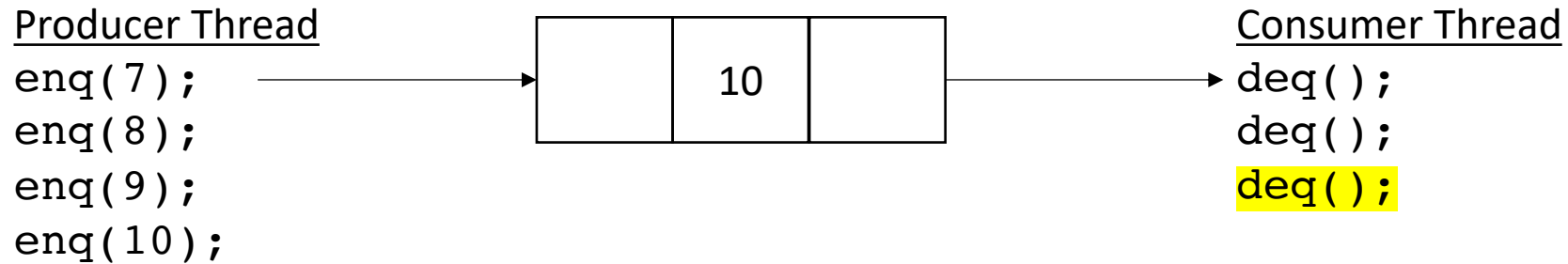
Consumer Thread

```
deq();  
deq();  
deq();
```

returns 9

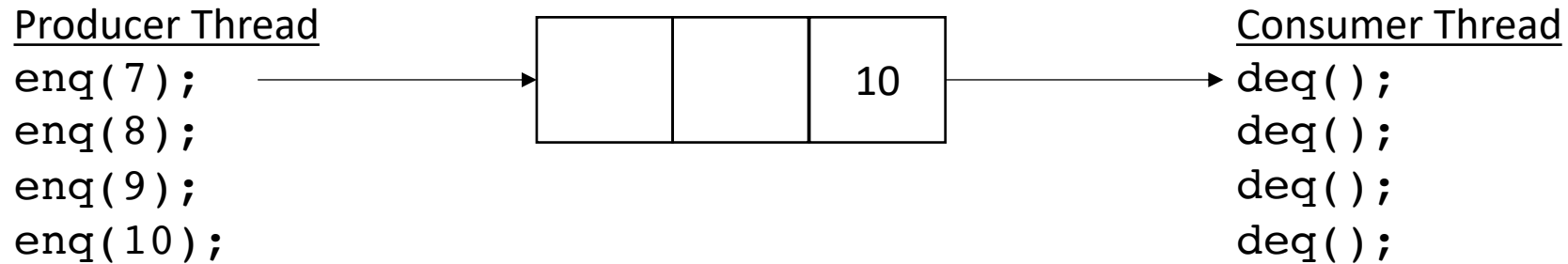
Producer Consumer Queues

- Asynchronous:



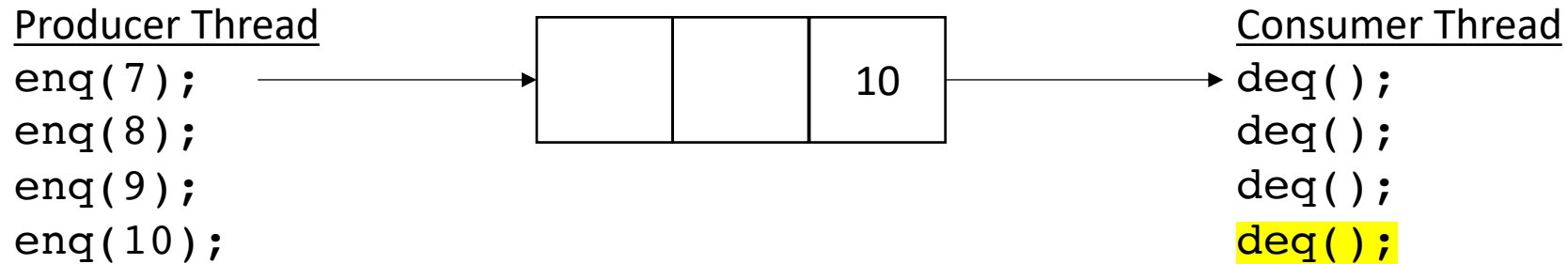
Producer Consumer Queues

- Asynchronous:



Producer Consumer Queues

- Asynchronous:



Producer Consumer Queues

- Asynchronous:

Producer Thread

```
enq(7);  
enq(8);  
enq(9);  
enq(10);
```



Consumer Thread

```
deq();  
deq();  
deq();  
deq();  
deq();
```

blocks when there is nothing in the queue

Producer Consumer Queues

- How do we implement it?

Producer Consumer Queues

- Start with a fixed size array



Producer Consumer Queues

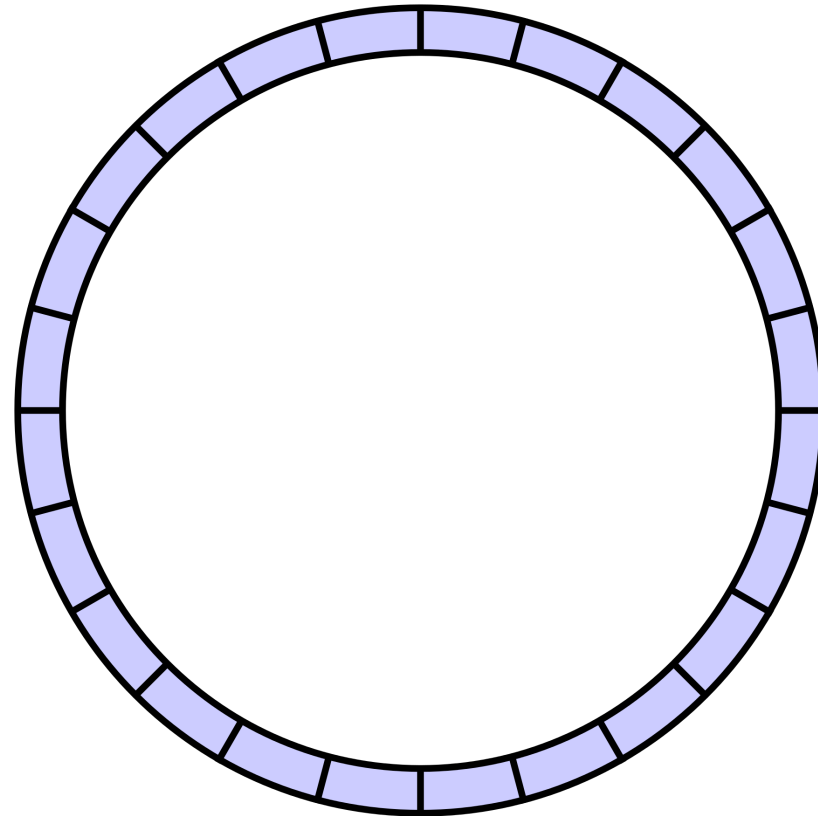
- Start with a fixed size array



We will use what is called a *circular buffer method*

Producer Consumer Queues

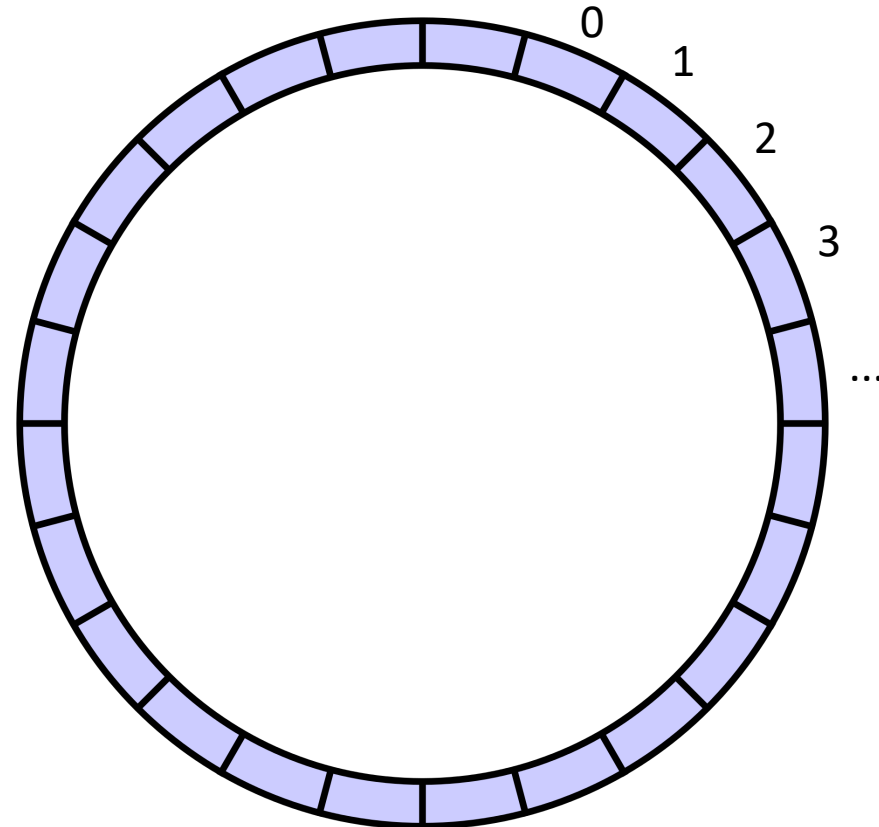
- Start with a fixed size array



conceptually it is a circle

Producer Consumer Queues

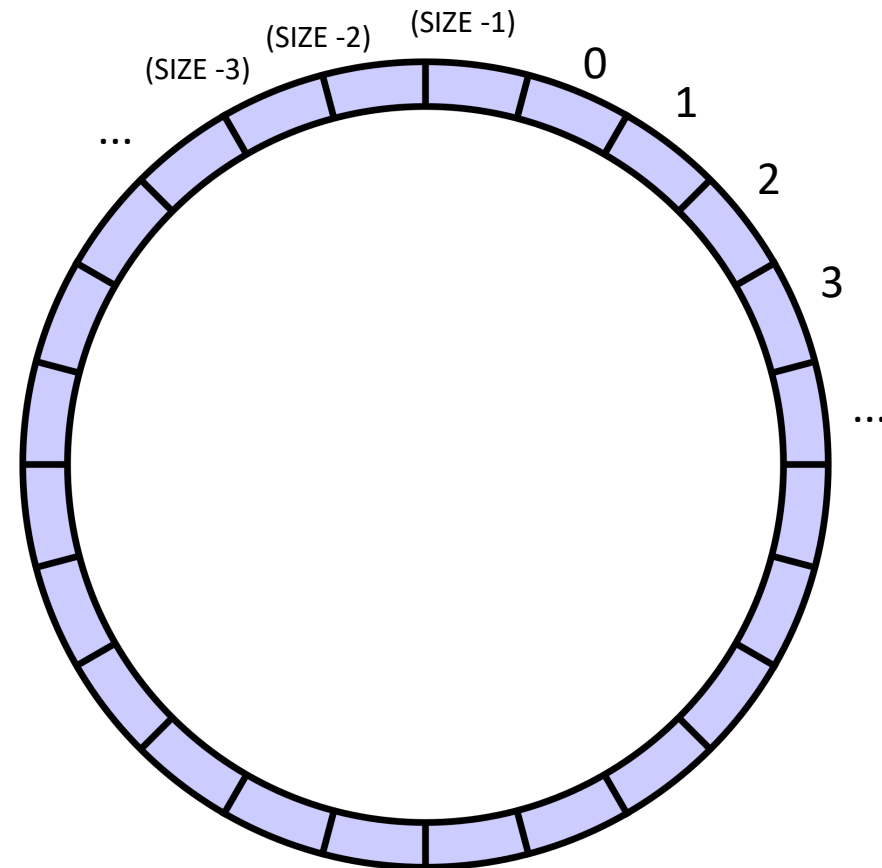
- Start with a fixed size array



conceptually it is a circle

Producer Consumer Queues

- Start with a fixed size array



indexes will circulate in order and wrap around

conceptually it is a circle

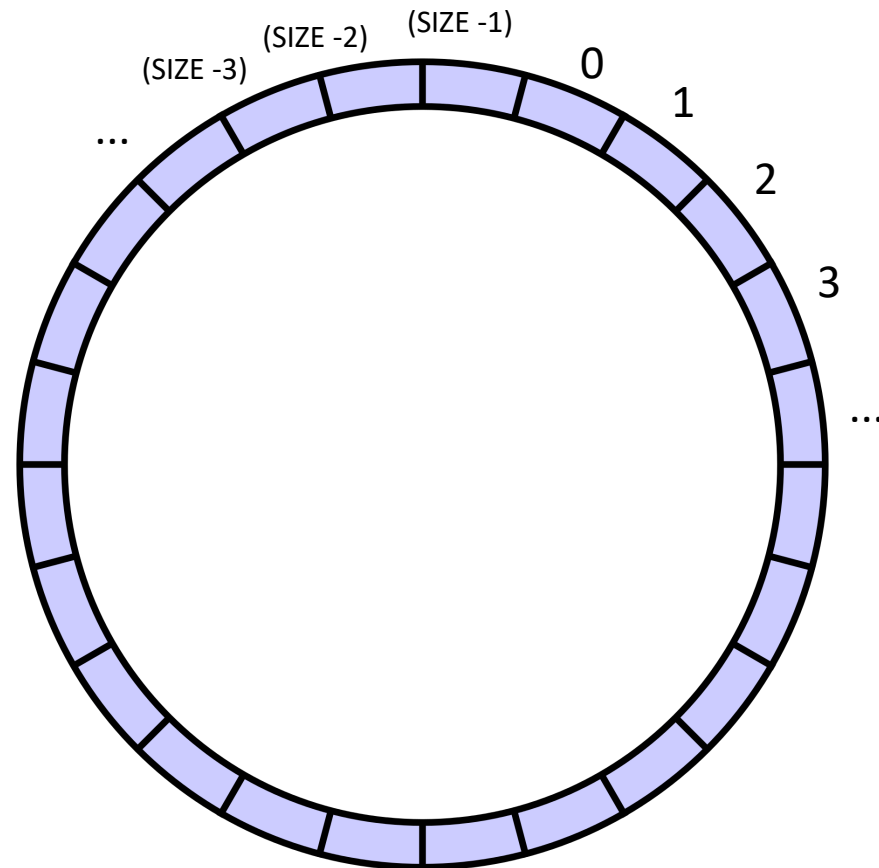
Producer Consumer Queues

- Start with a fixed size array

we will assume modular arithmetic:

if $x = (\text{SIZE} - 1)$ then
 $x + 1 == 0$;

conceptually it is a circle



indexes will circulate in order and wrap around

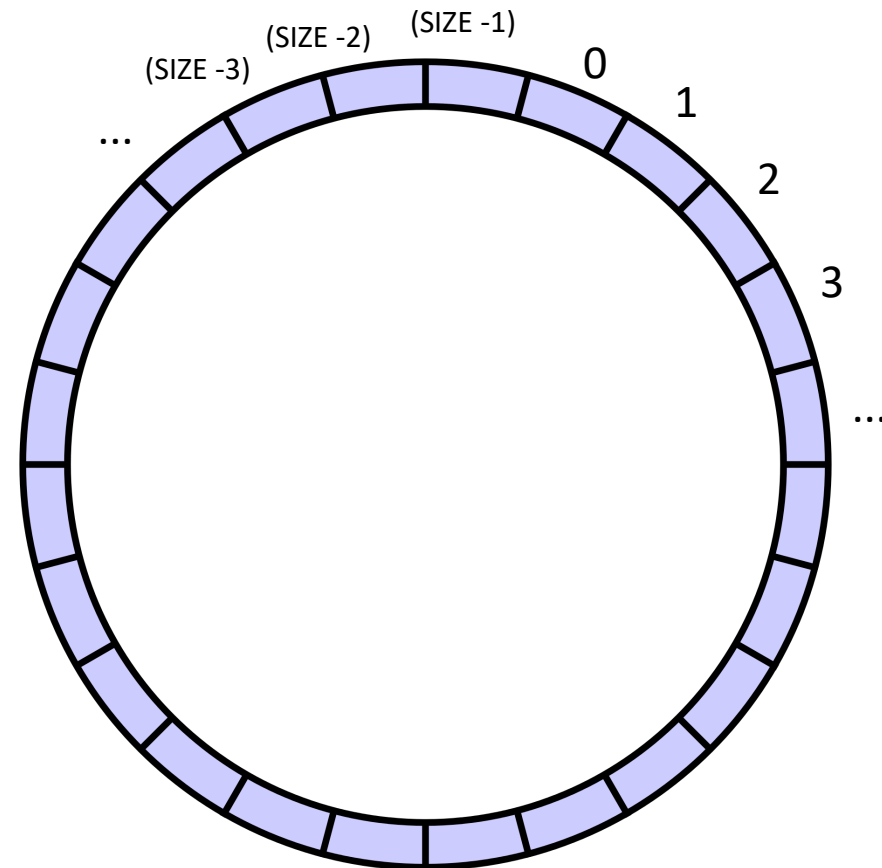
Producer Consumer Queues

- Start with a fixed size array

Two variables to keep track of where to deq and enq:

head and tail

conceptually it is a circle



indexes will circulate in order and wrap around

Producer Consumer Queues

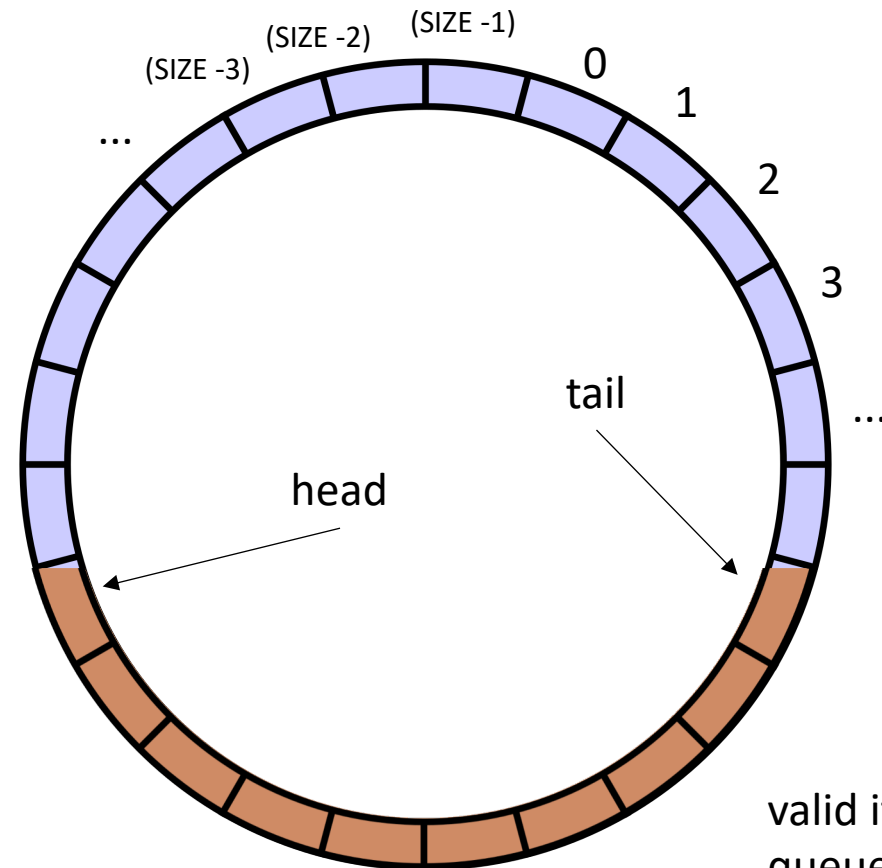
- Start with a fixed size array

Two variables to keep track of where to deq and enq:

head and tail:

enq to the head, deq from the tail

conceptually it is a circle



indexes will circulate in order and wrap around

valid items in the queue

Producer Consumer Queues

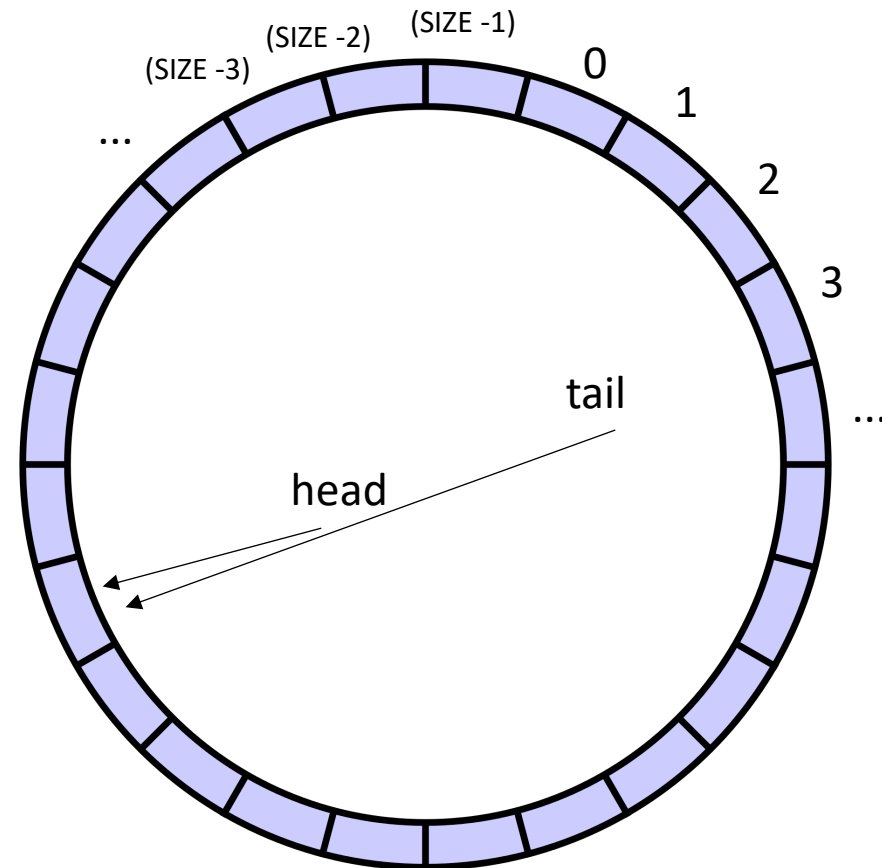
- Start with a fixed size array

Two variables to keep track of where to deq and enq:

head and tail

Empty queue is when $head == tail$

conceptually it is a circle



indexes will circulate in order and wrap around

Producer Consumer Queues

- Start with a fixed size array

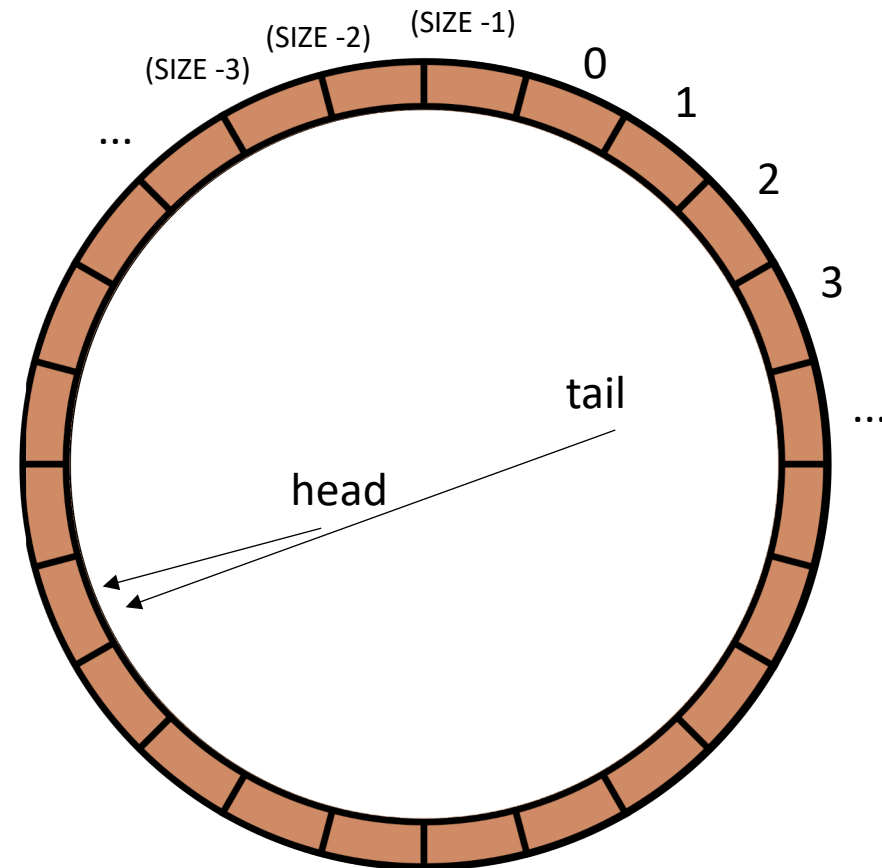
Two variables to keep track of where to deq and enq:

head and tail

Empty queue is when
 $head == tail$

Full queue is when
 $head == tail?$

conceptually it is a circle



indexes will circulate in order and wrap around

Producer Consumer Queues

- Start with a fixed size array

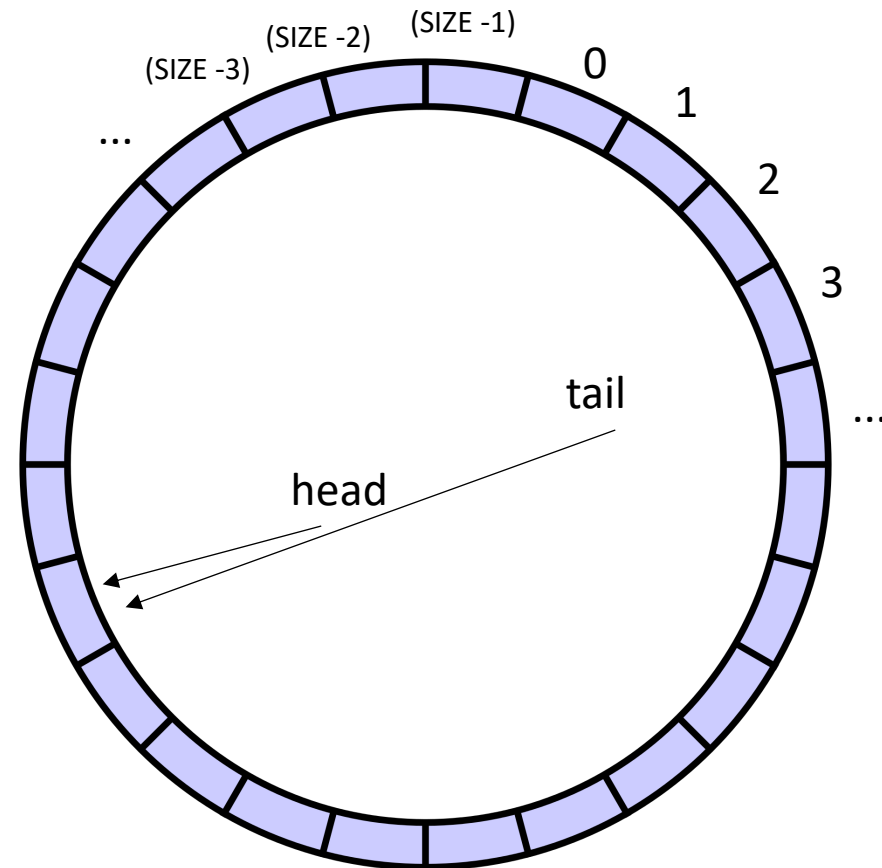
Two variables to keep track of where to deq and enq:

head and tail

Empty queue is when
 $head == tail$

Full queue is when
 $head == tail$?

conceptually it is a circle



indexes will circulate in order and wrap around

but then
how to tell
full queue from
empty?

Producer Consumer Queues

- Start with a fixed size array

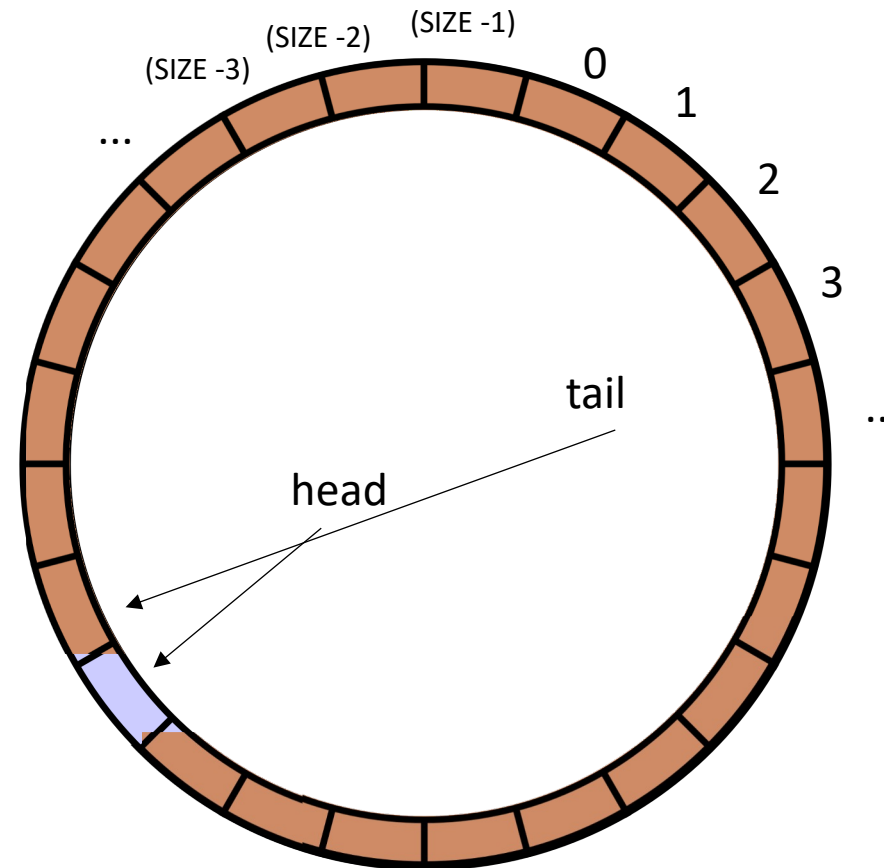
Two variables to keep track of where to deq and enq:

head and tail

Empty queue is when
 $\text{head} == \text{tail}$

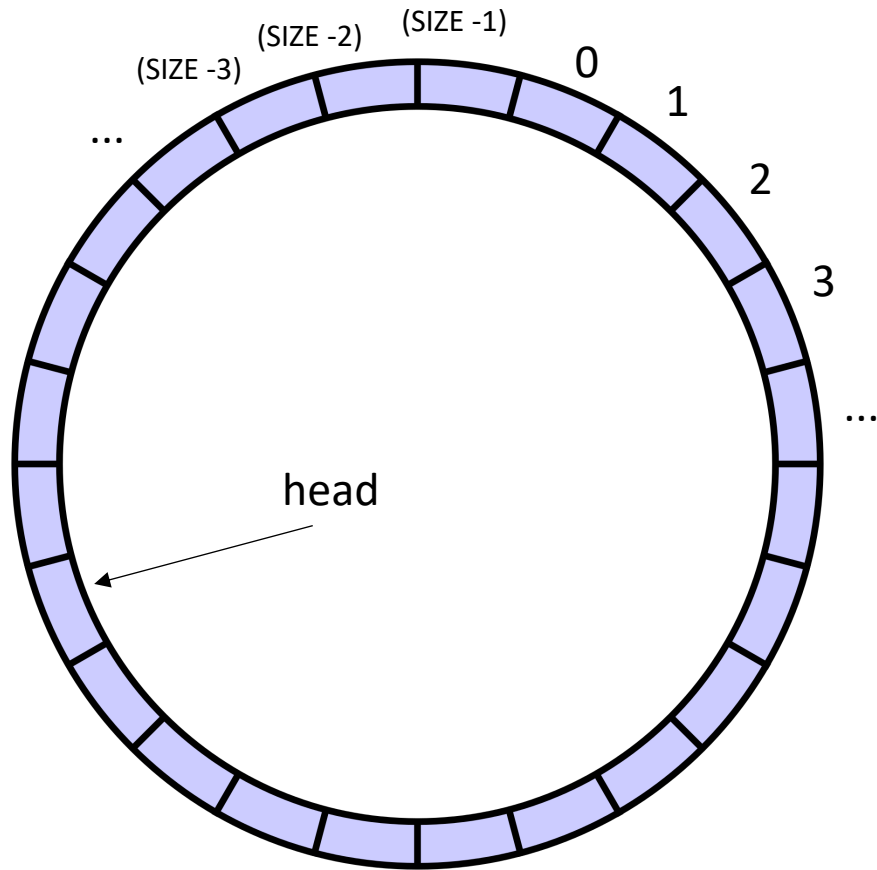
Full queue is when
 $\text{head} + 1 == \text{tail}$

conceptually it is a circle

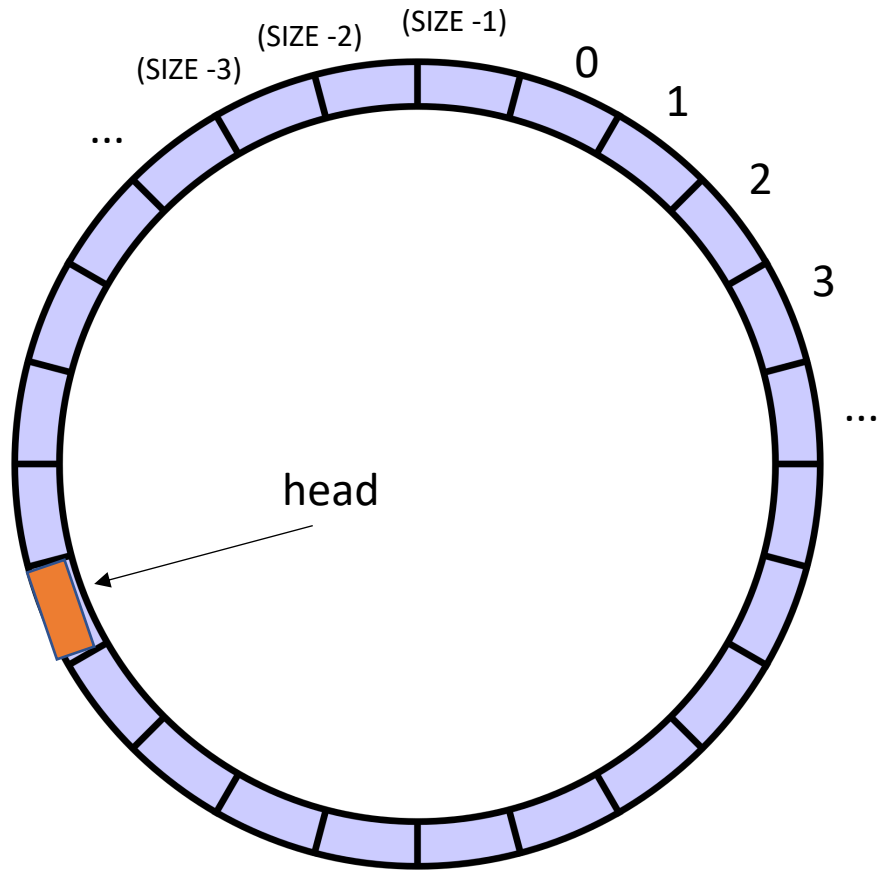


indexes will circulate in order and wrap around

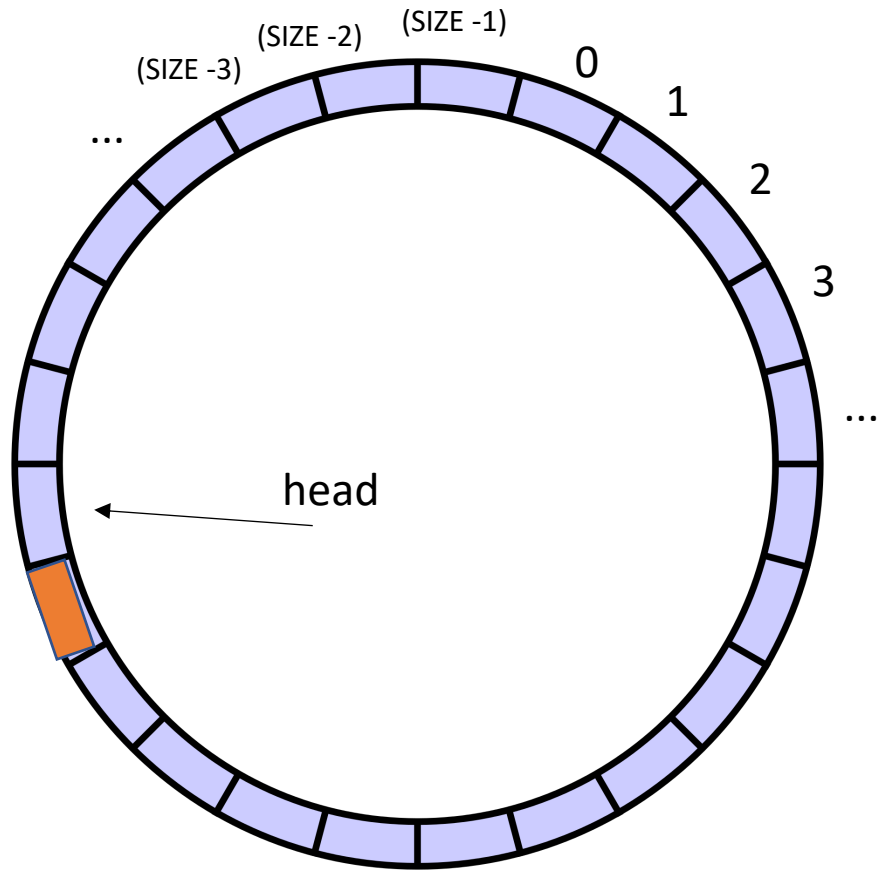
wasting one location, but its okay...



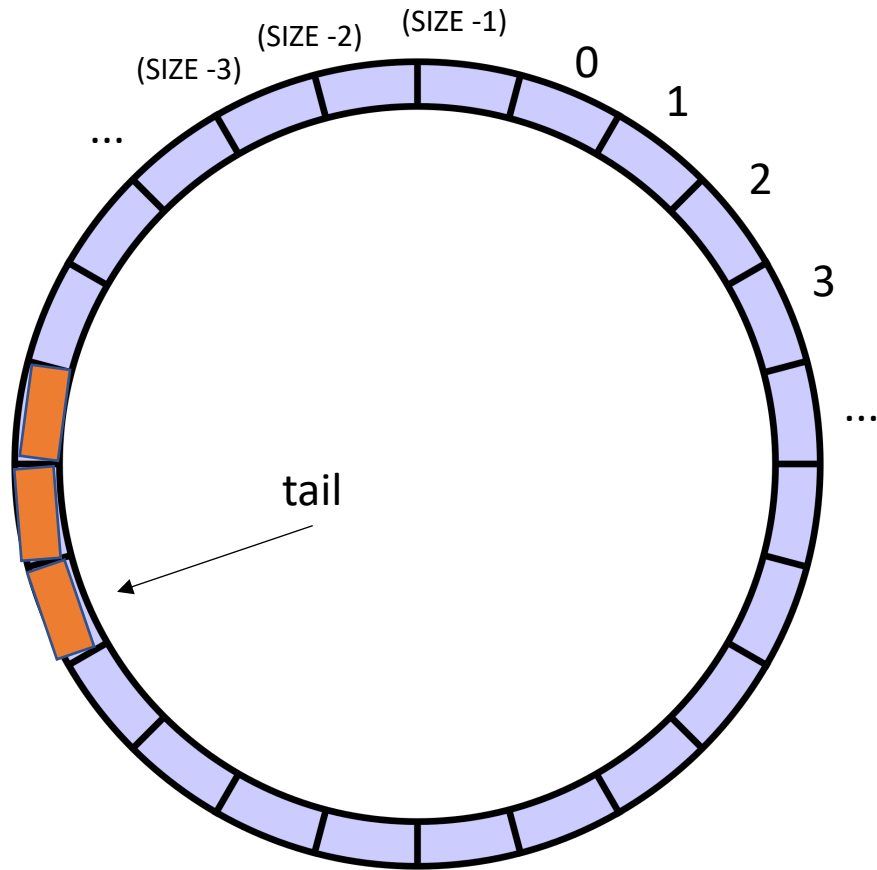
```
class ProdConsQueue {  
  private:  
    atomic_int head;  
    atomic_int tail;  
    int buffer[SIZE];  
  
  public:  
    void enq(int x) {  
      // store value at head  
      // increment head  
    }  
}
```



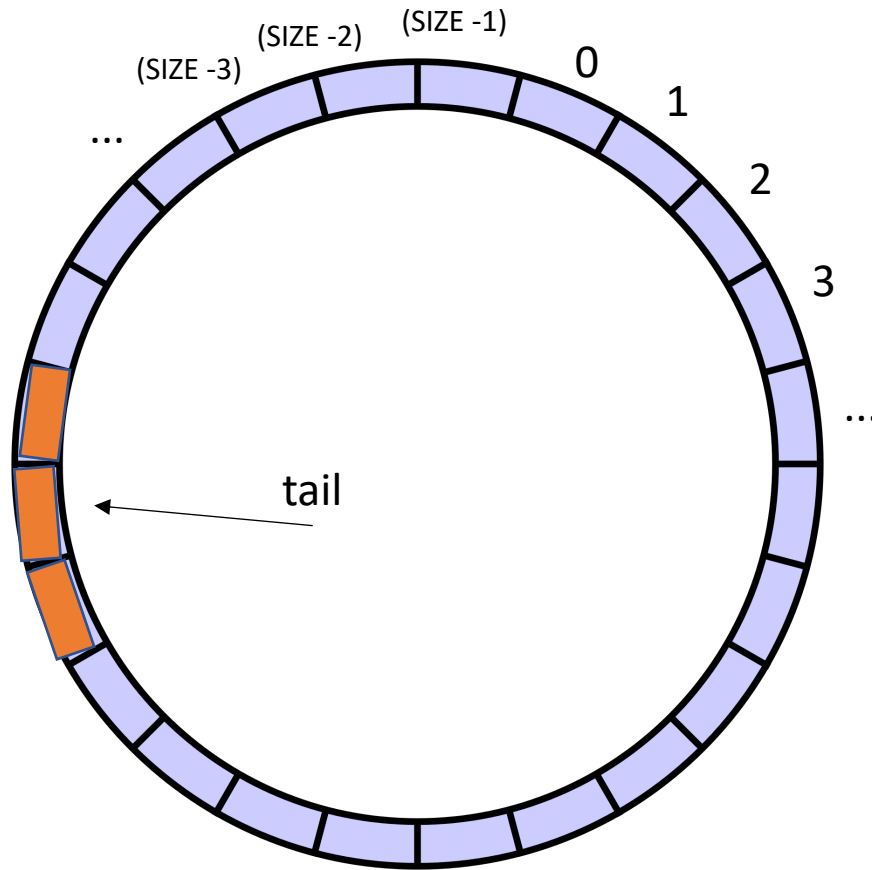
```
class ProdConsQueue {  
    private:  
        atomic_int head;  
        atomic_int tail;  
        int buffer[SIZE];  
  
    public:  
        void enq(int x) {  
            // store value at head  
            // increment head  
        }  
}
```



```
class ProdConsQueue {  
  private:  
    atomic_int head;  
    atomic_int tail;  
    int buffer[SIZE];  
  
  public:  
    void enq(int x) {  
      // store value at head  
      // increment head  
    }  
}
```

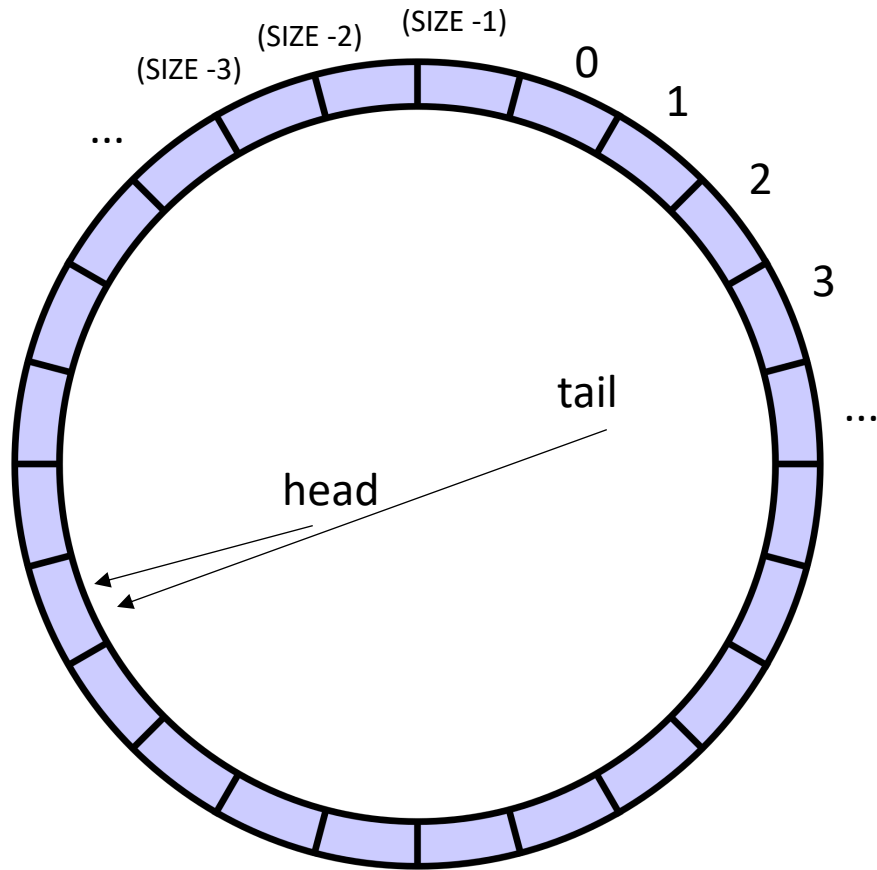


```
class ProdConsQueue {  
    private:  
        atomic_int head;  
        atomic_int tail;  
        int buffer[SIZE];  
  
    public:  
        void enq(int x) {  
            // store value at head  
            // increment head  
        }  
        int deq() {  
            // get value at tail  
            // increment tail  
        }  
}
```

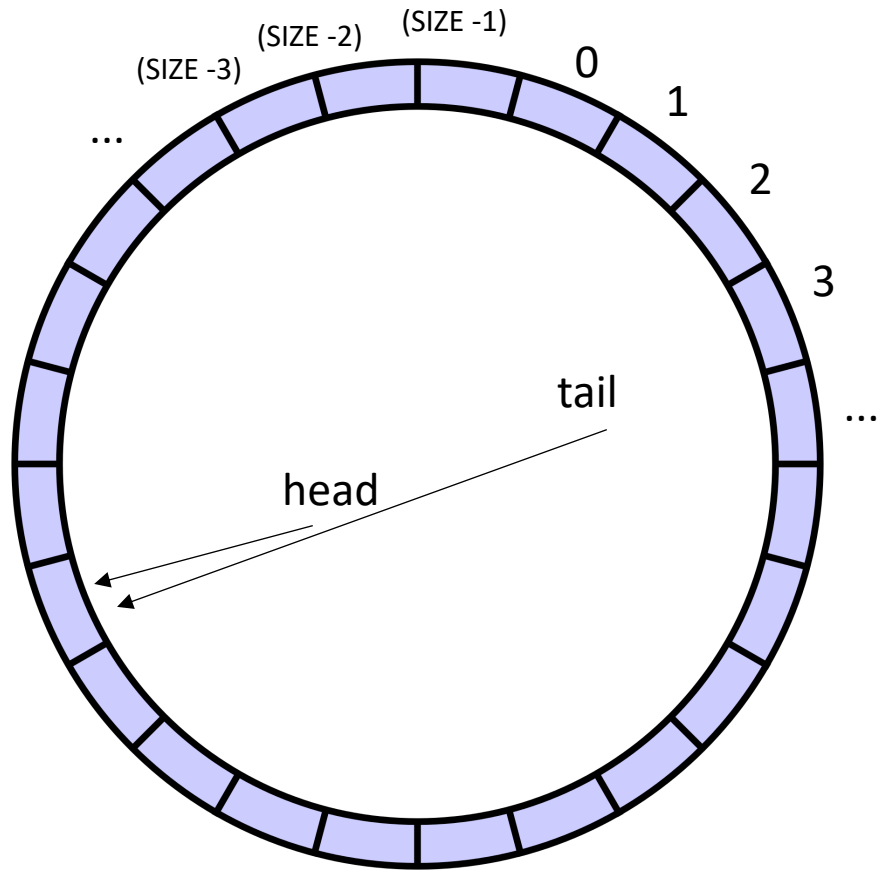
```
class ProdConsQueue {  
    private:  
        atomic_int head;  
        atomic_int tail;  
        int buffer[SIZE];  
  
    public:  
        void enq(int x) {  
            // store value at head  
            // increment head  
        }  
        int deq() {  
            // get value at tail  
            // increment tail  
        }  
}
```

This looks like the two threads don't even share head and tail! What is missing?

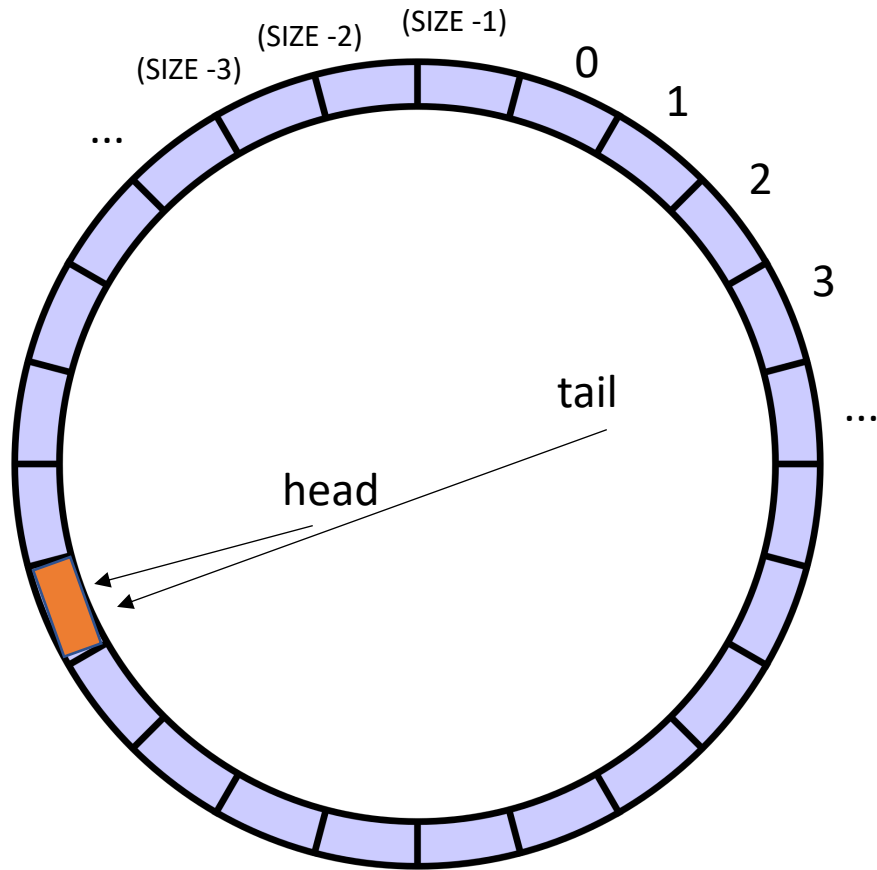


```
class ProdConsQueue {  
    private:  
        atomic_int head;  
        atomic_int tail;  
        int buffer[SIZE];  
  
    public:  
        void enq(int x) {  
            // store value at head  
            // increment head  
        }  
        int deq() {  
            // get value at tail  
            // increment tail  
        }  
}
```

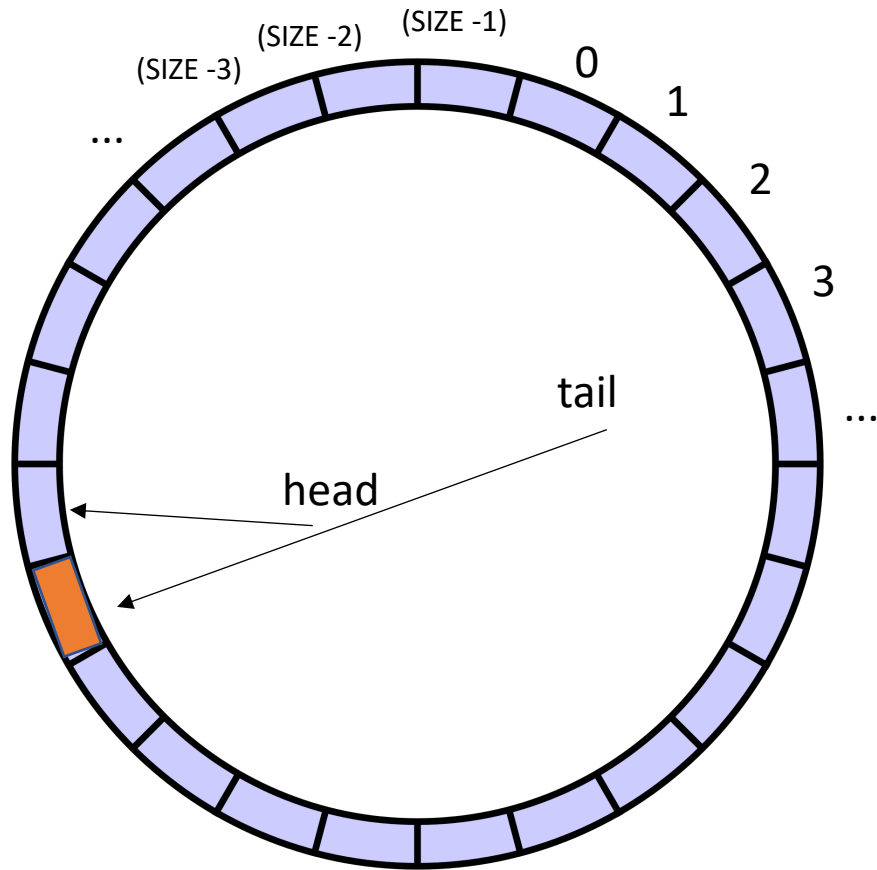
what happens if we try to dequeue here?



```
class ProdConsQueue {  
    private:  
        atomic_int head;  
        atomic_int tail;  
        int buffer[SIZE];  
  
    public:  
        void enq(int x) {  
            // store value at head  
            // increment head  
        }  
        int deq() {  
            // wait while queue is empty  
            // get value at tail  
            // increment tail  
        }  
}
```

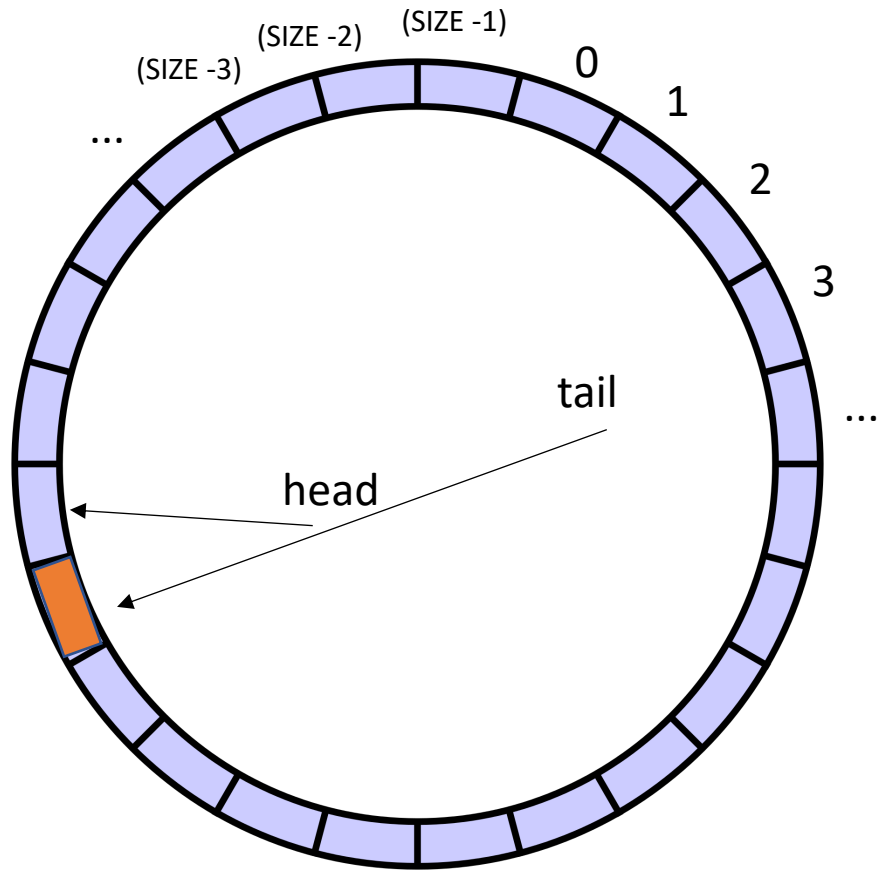


```
class ProdConsQueue {  
    private:  
        atomic_int head;  
        atomic_int tail;  
        int buffer[SIZE];  
  
    public:  
        void enq(int x) {  
            // store value at head  
            // increment head  
        }  
        int deq() {  
            // wait while queue is empty  
            // get value at tail  
            // increment tail  
        }  
}
```



```
class ProdConsQueue {
private:
    atomic_int head;
    atomic_int tail;
    int buffer[SIZE];

public:
    void enq(int x) {
        // store value at head
        // increment head
    }
    int deq() {
        // wait while queue is empty
        // get value at tail
        // increment tail
    }
}
```

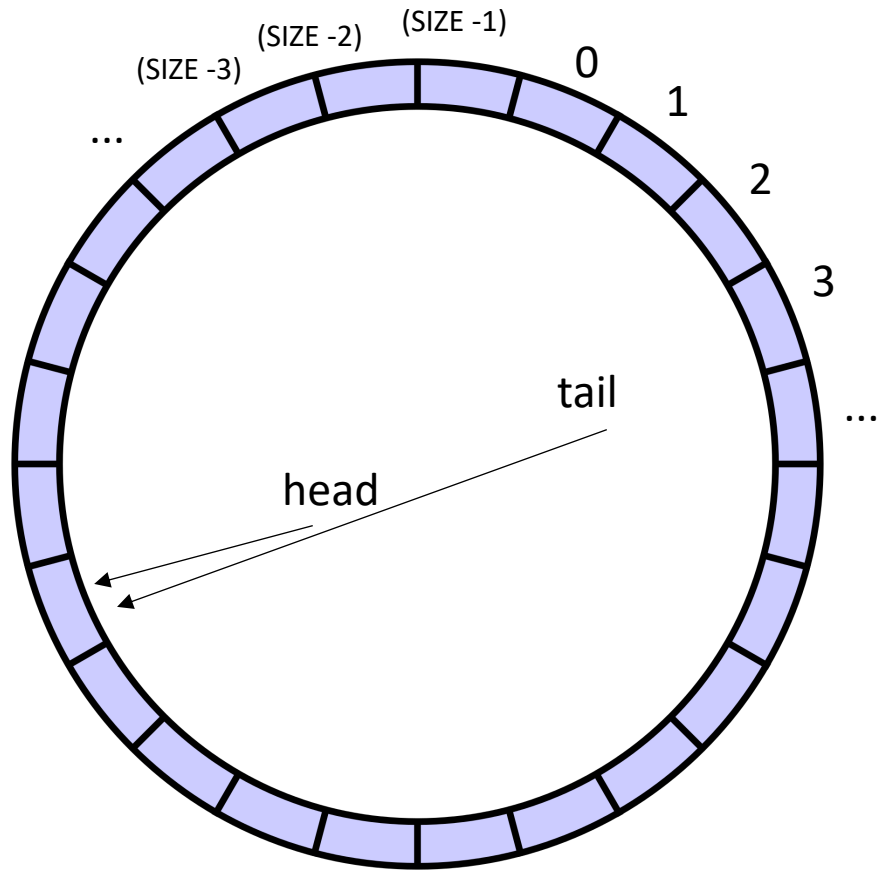


```

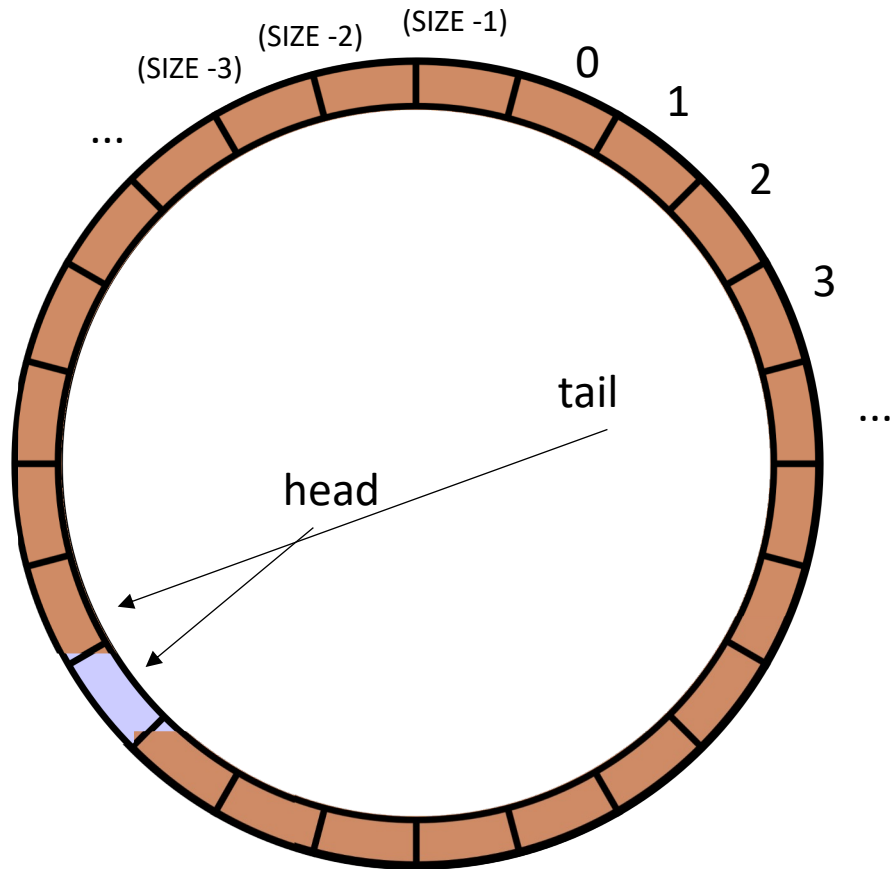
class ProdConsQueue {
private:
    atomic_int head;
    atomic_int tail;
    int buffer[SIZE];

public:
    void enq(int x) {
        // store value at head
        // increment head
    }
    int deq() {
        // wait while queue is empty
        // get value at tail
        // increment tail
    }
}

```



```
class ProdConsQueue {  
  private:  
    atomic_int head;  
    atomic_int tail;  
    int buffer[SIZE];  
  
  public:  
    void enq(int x) {  
      // store value at head  
      // increment head  
    }  
    int deq() {  
      // wait while queue is empty  
      // get value at tail  
      // increment tail  
    }  
}
```



```

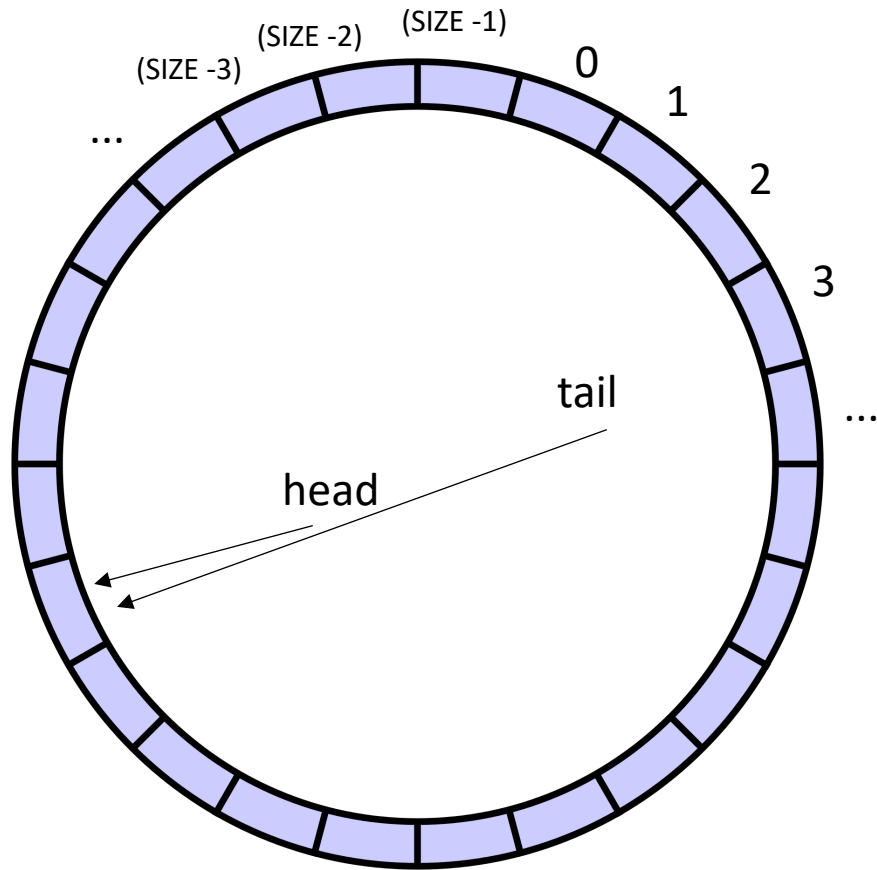
class ProdConsQueue {
private:
    atomic_int head;
    atomic_int tail;
    int buffer[SIZE];

public:
    void enq(int x) {
        // store value at head
        // increment head
    }
    int deq() {
        // wait while queue is empty
        // get value at tail
        // increment tail
    }
}

```

similarly for enqueue

but why can't we enqueue?



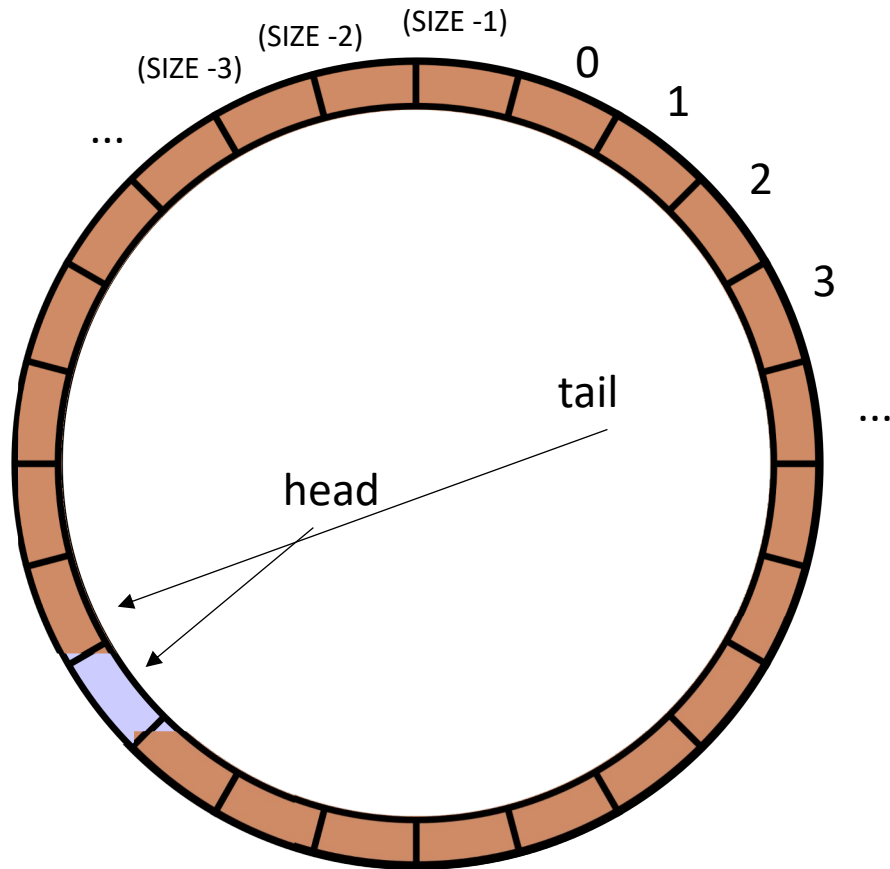
```

class ProdConsQueue {
private:
    atomic_int head;
    atomic_int tail;
    int buffer[SIZE];

public:
    void enq(int x) {
        // store value at head
        // increment head
    }
    int deq() {
        // wait while queue is empty
        // get value at tail
        // increment tail
    }
}

```

incrementing the head would make it empty!



we need to wait for there
to be room

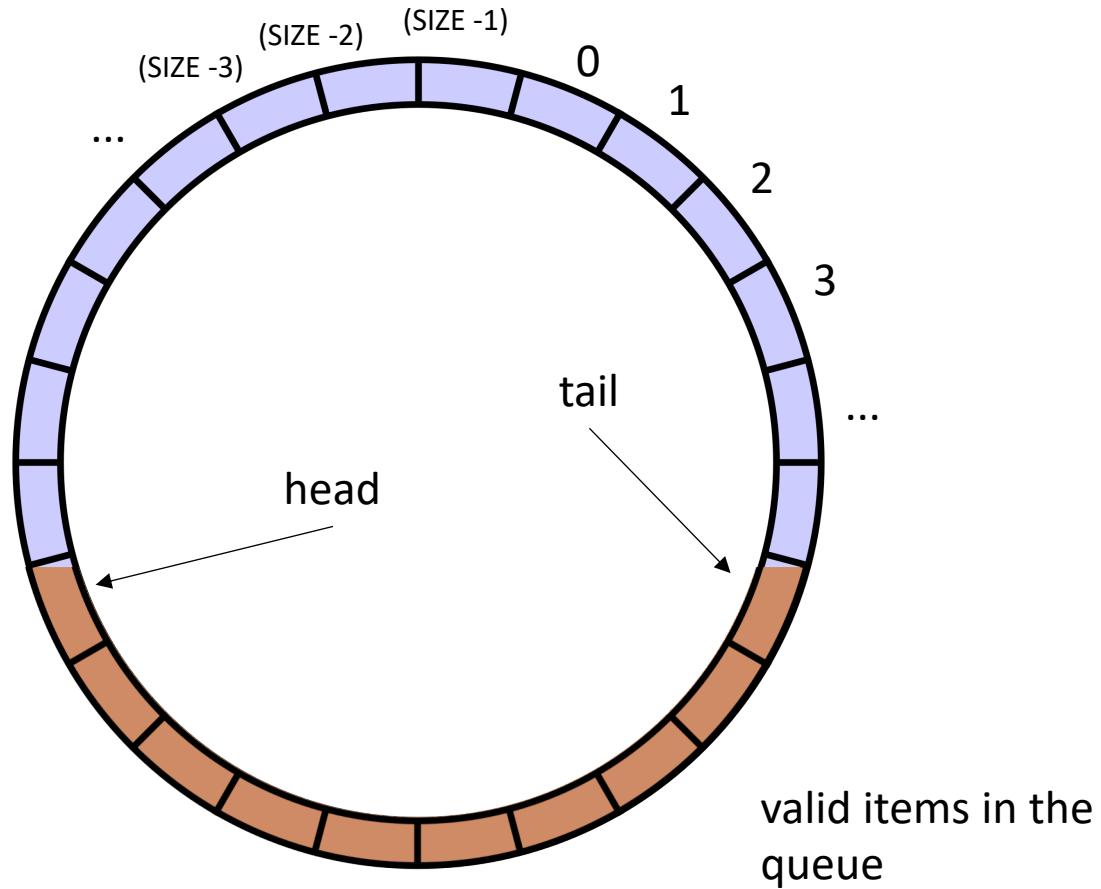
```

class ProdConsQueue {
private:
    atomic_int head;
    atomic_int tail;
    int buffer[SIZE];

public:
    void enq(int x) {
        // wait for there to be room
        // store value at head
        // increment head
    }
    int deq() {
        // wait while queue is empty
        // get value at tail
        // increment tail
    }
}

```

Other questions:

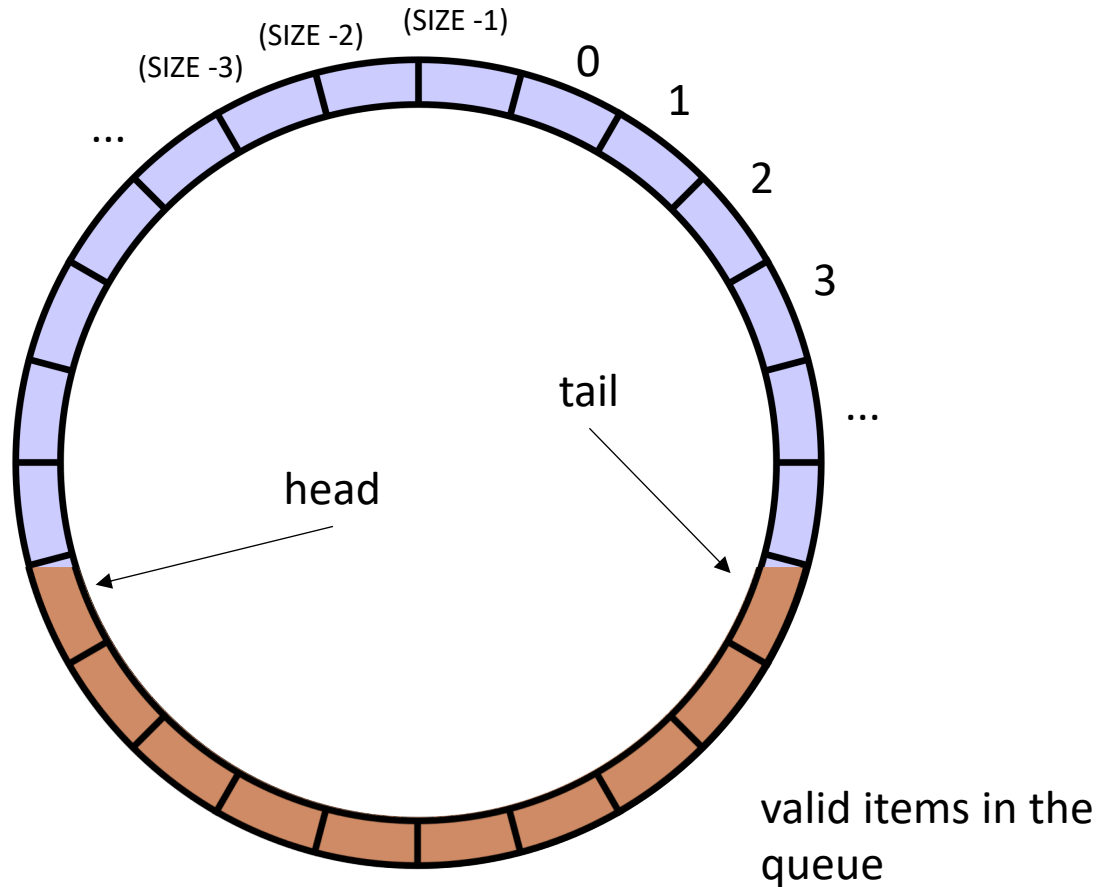


```
class ProdConsQueue {
private:
    atomic_int head;
    atomic_int tail;
    int buffer[SIZE];

public:
    void enq(int x) {
        // wait for their to be room
        // store value at head
        // increment head
    }
    int deq() {
        // wait while queue is empty
        // get value at tail
        // increment tail
    }
}
```

Other questions:

Do these need to be atomic RMWs?



```
class ProdConsQueue {
private:
    atomic_int head;
    atomic_int tail;
    int buffer[SIZE];

public:
    void enq(int x) {
        // wait for their to be room
        // store value at head
        // increment head
    }
    int deq() {
        // wait while queue is empty
        // get value at tail
        // increment tail
    }
}
```

Next week

- Work stealing and generalized concurrent objects
- Get HW 2 turned in today!
- HW 3 is out today. You can get started on Part 1
- Prepare for midterm on Monday