CSE113: Parallel Programming Feb. 7, 2022

- Topics: Workstealing
 - DOALL loops
 - Parallel Schedules:
 - Static
 - global worklist
 - local worklist



Announcements

- Midterm is out!
 - You have until next Monday at midnight to do it.
 - Do not discuss with your classmates
 - Do not google specific questions or ask on online forums
 - Ask any clarifying questions as a private post on piazza
 - Late tests will not be accepted
 - You can ask me or Reese about the midterm, not Tim or Sanya
- Homework 3 is out
 - You have what you need for part 1
 - We'll be doing the material for part 2 today
 - You can share results starting in 1 week

Announcements

- Grades for HW 1 will be released after class today
- Halfway through the quarter!

Today's Quiz

• Due tomorrow by midnight. Please do it!

Input/output queues use atomic increments and decrements to protect against threads that are trying to concurrently enqueue and dequeue

⊖ True

⊖ False

A circular buffer is:

 \bigcirc A special type of memory that is organized in circular pattersTrue

○ Part of the C++ standard library

 \bigcirc a useful data representation for fixed-length queues

The performance of an application using a producer-consumer queue depends most on:

 \bigcirc The rate at which the producer enqueues elements

○ The rate at which the consumer dequeues elements

 $\bigcirc\,$ if the queue is implemented using a mutex or not

- Performance modeling:
 - Get timing for "perfect producer" and "perfect consumer"
 - i.e. time everything but delays from concurrent queue
 - Call these PP and PC
 - The fastest your application can possibly go is max(PP, PC).
- What is the relation between producer and consumer?
 - If the consumer goes much slower than the producer?
 - If the producer goes much slower than the consumer?

Review

- 1 enq, 1 deq
 - enq'er cannot deq
 - deq'er cannot enq
- Example: printf:
 - your program equeues values to print
 - the terminal process dequeues values and prints them

Synchronous Producer Consumer Queues



Async Producer Consumer Queues

• Start with a fixed size array

we will assume modular arithmetic:

if x = (SIZE - 1) then x + 1 == 0;



indexes will circulate in order and wrap around

conceptually it is a circle

• Start with a fixed size array

Two variables to keep track of where to deq and enq:

head and tail:

enq to the head, deq from the tail



indexes will circulate in order and wrap around

conceptually it is a circle

• Start with a fixed size array

Two variables to keep track of where to deq and enq:

head and tail

Empty queue is when head == tail

Full queue is when head == tail?

conceptually it is a circle



indexes will circulate in order and wrap around

but then

empty?

• Start with a fixed size array

Two variables to keep track of where to deq and enq:

head and tail

Empty queue is when head == tail

Full queue is when head == tail?

conceptually it is a circle



indexes will circulate in order and wrap around

• Start with a fixed size array

Two variables to keep track of where to deq and enq:

head and tail

Empty queue is when head == tail

Full queue is when head + 1 == tail

conceptually it is a circle



indexes will circulate in order and wrap around



Other questions:

Do these need to be atomic RMWs?

class ProdConsQueue {
private:
<pre>atomic_int head;</pre>
<pre>atomic_int tail;</pre>
<pre>int buffer[SIZE];</pre>
public:
<pre>void enq(int x) {</pre>
// wait for their to be room
// store value at head
// increment head
}
int deq() {
// wait while queue is empty
// get value at tail
// increment tail
}
}

On to new material!

Schedule

- Workstealing
 - DOALL Loops
 - Parallel Schedules
 - Static schedule
 - Global worklist
 - Local worklists

Schedule

- Workstealing
 - DOALL Loops
 - Parallel Schedules
 - Static schedule
 - Global worklist
 - Local worklists

```
adds two arrays
for (int i = 0; i < SIZE; i++) {
    a[i] = b[i] + c[i];
}</pre>
```

adds elements with neighbors

```
for (int i = 0; i < SIZE; i++) {
    a[i] += a[i+1]
}</pre>
```

are they the same if you traverse them backwards?

adds two arrays

```
for (int i = 0; i < SIZE; i++) {
    a[i] = b[i] + c[i];
}</pre>
```

```
for (int i = SIZE-1; i >= 0; i--) {
    a[i] = b[i] + c[i];
}
```

adds elements with neighbors

```
for (int i = 0; i < SIZE; i++) {
    a[i] += a[i+1]
}</pre>
```

```
for (int i = SIZE-1; i >= 0; i--) {
    a[i] += a[i+1]
}
```

are they the same if you traverse them backwards?

adds two arrays

```
for (int i = 0; i < SIZE; i++) {
    a[i] = b[i] + c[i];
}</pre>
```

```
for (int i = SIZE-1; i >= 0; i--) {
    a[i] = b[i] + c[i];
}
```

adds elements with neighbors

```
for (int i = 0; i < SIZE; i++) {
    a[i] += a[i+1]
}</pre>
```

```
for (int i = SIZE-1; i >= 0; i--) {
    a[i] += a[i+1]
}
```

No!

what about a random order?

adds two arrays

```
for (int i = 0; i < SIZE; i++) {
    a[i] = b[i] + c[i];
}</pre>
```

```
for (pick i randomly) {
    a[i] = b[i] + c[i];
}
```

adds elements with neighbors

```
for (int i = 0; i < SIZE; i++) {
    a[i] += a[i+1]
}</pre>
```

```
for (pick i randomly) {
    a[i] += a[i+1]
}
```

what about a random order?

adds two arrays

```
for (int i = 0; i < SIZE; i++) {
    a[i] = b[i] + c[i];
}</pre>
```

```
for (pick i randomly) {
    a[i] = b[i] + c[i];
}
```

adds elements with neighbors

```
for (int i = 0; i < SIZE; i++) {
    a[i] += a[i+1]
}</pre>
```

```
for (pick i randomly) {
    a[i] += a[i+1]
}
```

```
for (int i = 0; i < SIZE; i++) {
    a[i] = b[i] + c[i];
}</pre>
```

These are **DOALL** loops:

- Loop iterations are independent
- You can do them in ANY order and get the same results

```
for (int i = 0; i < SIZE; i++) {
    a[i] = b[i] + c[i];
}</pre>
```

These are **DOALL** loops:

- Loop iterations are independent
- You can do them in ANY order and get the same results
- Most importantly: you can do the iterations in parallel!
- Assign each thread a set of indices to compute

- Given a nest of For loops, can we make the outer-most loop parallel?
 - Safely
 - Efficiently

- We will consider a special type of for loop, common in scientific applications:
 - Operates on N dimensional arrays (only side-effects are array writes)
 - Array bases are disjoint and constant
 - Bounds, indexes are a function of loop variables, input variables and constants
 - Loops Increment by 1

```
for (int i = 0; i < dim1; i++) {
  for (int j = 0; j < dim3; j++) {
    for (int k = 0; k < dim2; k++) {
        a[i][j] += b[i][k] * c[k][j];
        example
    }
}</pre>
```

- We will consider a special type of for loop, common in scientific applications:
 - Operates on N dimensional arrays (only side-effects are array writes)
 - Array bases are disjoint and constant
 - Bounds, indexes are a function of loop variables, input variables and constants
 - Loops Increment by 1

- We will consider a special type of for loop, common in scientific applications:
 - Operates on N dimensional arrays (only side-effects are array writes)
 - Array bases are disjoint and constant
 - Bounds, indexes are a function of loop variables, input variables and constants
 - Loops Increment by 1

- We will consider a special type of for loop, common in scientific applications:
 - Operates on N dimensional arrays (only side-effects are array writes)
 - Array bases are disjoint and constant
 - Bounds, indexes are a function of loop variables, input variables and constants
 - Loops Increment by 1

- We will consider a special type of for loop, common in scientific applications:
 - Operates on N dimensional arrays (only side-effects are array writes)
 - Array bases are disjoint and constant
 - Bounds, indexes are a function of loop variables, input variables and constants
 - Loops Increment by 1

```
for (int i = 2; i < 100; i+=3) {
    a[i] = c[i + 128];
}</pre>
```

```
for (int j = 0; j < 33; j+=1) {
    a[3*j+2] = c[(3*j+2) + 128];
}</pre>
```

- Given a nest of *candidate For* loops, determine if we can we make the outer-most loop parallel?
 - Safely
 - efficiently
- Criteria: every iteration of the outer-most loop must be *independent*
 - The loop can execute in any order, and produce the same result

Safety Criteria

- How do we check this?
 - If the property doesn't hold then there exists 2 iterations, such that if they are re-ordered, it causes different outcomes for the loop.
 - Write-Write conflicts: two distinct iterations write different values to the same location
 - **Read-Write conflicts**: two distinct iterations where one iteration reads from the location written to by another iteration.

Safety Criteria

- Criteria: every iteration of the outer-most loop must be *independent*
- the loop must produce the same result for any order of the iterations

```
for (i = 0; i < size; i++) {
    a[index(i)] = loop(i);
}</pre>
```
- Criteria: every iteration of the outer-most loop must be *independent*
- the loop must produce the same result for any order of the iterations

```
for (i = 0; i < size; i++) {
    a[index(i)] = loop(i);
}</pre>
```

index calculation based on the loop variable

- Criteria: every iteration of the outer-most loop must be *independent*
- the loop must produce the same result for any order of the iterations

```
for (i = 0; i < size; i++) {
    a[index(i)] = loop(i);
}</pre>
```

index calculation based on the loop variable Computation to store in the memory location

- Criteria: every iteration of the outer-most loop must be *independent*
- the loop must produce the same result for any order of the iterations

```
for (i = 0; i < size; i++) {
    a[index(i)] = loop(i);
}</pre>
```

Write-write conflicts:

for two distinct iteration variables: $i_x != i_y$ Check: $index(i_x) != index(i_y)$

- Criteria: every iteration of the outer-most loop must be *independent*
- the loop must produce the same result for any order of the iterations

```
for (i = 0; i < size; i++) {
    a[index(i)] = loop(i);
}</pre>
```

Write-write conflicts:

for two distinct iteration variables: $i_x != i_y$ Check: $index(i_x) != index(i_y)$

```
Why?
Because if
index(i<sub>x</sub>) == index(i<sub>y</sub>)
then:
a[index(i<sub>x</sub>)] will equal
either loop(i<sub>x</sub>) or loop(i<sub>y</sub>)
depending on the order
```

• Criteria: every iteration of the outer-most loop must be *independent*

Read-write conflicts:

```
for two distinct iteration variables:
i<sub>x</sub> != i<sub>y</sub>
Check:
write_index(i<sub>x</sub>) != read_index(i<sub>y</sub>)
```

• Criteria: every iteration of the outer-most loop must be *independent*

Read-write conflicts:

for two distinct iteration variables:

i_x != i_y Check: write_index(i_x) != read_index(i_y)

Why?

if i_x iteration happens first, then iteration i_y reads an updated value.

if i_y happens first, then it reads the original value

```
Examples:
```

```
for (i = 0; i < 128; i++) {
    a[i]= a[i]*2;
}</pre>
```

Examples:

```
for (i = 0; i < 128; i++) {
    a[i]= a[i]*2;
}
for (i = 0; i < 128; i++) {
    a[i]= a[0]*2;
}</pre>
```

Examples:

}

```
for (i = 0; i < 128; i++) {
    a[i]= a[i]*2;
}
for (i = 0; i < 128; i++) {
    a[i]= a[0]*2;</pre>
```

```
for (i = 1; i < 128; i++) {
    a[i]= a[0]*2;
}</pre>
```

```
Examples:
```

```
for (i = 0; i < 128; i++) {
  a[i]= a[i]*2;
}
for (i = 0; i < 128; i++) {
  a[i]= a[0]*2;
}
                                      }
for (i = 0; i < 128; i++) {
  a[i%64]= a[i]*2;
}
```

```
for (i = 1; i < 128; i++) {
    a[i]= a[0]*2;
}</pre>
```

Examples:

```
for (i = 0; i < 128; i++) {
  a[i]= a[i]*2;
}
for (i = 0; i < 128; i++) {
  a[i]= a[0]*2;
}
for (i = 0; i < 128; i++) {
  a[i%64]= a[i]*2;
}
```

```
for (i = 1; i < 128; i++) {
    a[i]= a[0]*2;
}
for (i = 0; i < 128; i++) {
    a[i%64]= a[i+64]*2;
}</pre>
```

Schedule

• DOALL Loops

- Static
- Global Worklists
- Local Worklists

• Consider the following program:

There are 3 arrays: a, b, c. We want to compute

```
for (int i = 0; i < SIZE; i++) {
    c[i] = a[i] + b[i];
}</pre>
```

Is this a DOALL loop?

• Consider the following program:

There are 3 arrays: a, b, c. We want to compute

```
for (int i = 0; i < SIZE; i++) {
    c[i] = a[i] + b[i];
}</pre>
```











• Which one is more efficient?

- Which one is more efficient?
- These are called Parallel Schedules for DOALL Loops
- We will discuss several of them today.

Schedule

- DOALL Loops
- Parallel Schedules:
 - Static
 - Global Worklists
 - Local Worklists

• Works well when loop iterations take similar amounts of time

```
void foo() {
    ...
    for (int x = 0; x < SIZE; x++) {
        // Each iteration takes roughly
        // equal time
     }
    ...
}</pre>
```

0	1	2	3	4	5	6	7	SIZE -1

• Works well when loop iterations take similar amounts of time

```
void foo() {
    for (int x = 0; x < SIZE; x++) {
        // Each iteration takes roughly
        // equal time
     }
        say SIZE / NUM_THREADS = 4
}</pre>
```

0	1	2	3	4	5	6	7		SIZE -1
---	---	---	---	---	---	---	---	--	---------

• Works well when loop iterations take similar amounts of time

```
void foo() {
. . .
  for (int x = 0; x < SIZE; x++) {
  // Each iteration takes roughly
  // equal time
  }
                                                   say SIZE / NUM_THREADS = 4
. . .
}
     Thread 0
                            Thread 1
                                                  Thread N
            2
                 3
                            5
                                                   SIZE -1
 0
       1
                       4
                                  6
                                       7
```

• Works well when loop iterations take similar amounts of time

```
void foo() {
    ...
    for (int x = 0; x < SIZE; x++) {
        // Each iteration takes roughly
        // equal time
     }
    ...
}</pre>
```

make a new function with the for loop inside. Pass all needed variables as arguments. Take an extra argument for a thread id

• Works well when loop iterations take similar amounts of time



make a new function with the for loop inside. Pass all needed variables as arguments. Take an extra argument for a thread id

Works well when loop iterations take similar amounts of time

```
void foo() {
    for (int x = 0; x < SIZE; x++) {
        // Each iteration takes roughly
        // equal time
    }
    ...
}</pre>
void parallel_loop(..., int tid, int num_threads)
{
        int chunk_size = SIZE / NUM_THREADS;
        for (int x = 0; x < SIZE; x++) {
            // work based on x
        }
    }
}
```

determine chunk size in new function

Works well when loop iterations take similar amounts of time

Works well when loop iterations take similar amounts of time

```
void foo() {
    ...
    for (int t = 0; t < NUM_THREADS; t++) {
        spawn(parallel_loop(..., t, NUM_THREADS))
        int
        join();
    ...
    }
}</pre>
```

```
void parallel_loop(..., int tid, int num_threads)
{
```

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   // work based on x
}
```

You will need to adapt the thread spawn, join to C++

Spawn threads

• Example, 2 threads/cores, array of size 8

0	1	2	3	4	5	6	7
---	---	---	---	---	---	---	---

thread 1

```
chunk size = 4
```

```
0: start = ? 1: start = ?
```

0: end = ? 1: end = ?

thread 0

void parallel_loop(..., int tid, int num_threads) {

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   // work based on x
}
```

}

• Example, 2 threads/cores, array of size 8

0	1	2	3	4	5	6	7
---	---	---	---	---	---	---	---

```
chunk_size = 4
0: start = 0 1: start = 4
0: end = 4 1: end = 8
thread 0
thread 1
```

```
void parallel_loop(..., int tid, int num_threads)
{
```

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   // work based on x
}
```

}

End example

• Example, 2 threads/cores, array of size 9

0	1	2	3	4	5	6	7	8	
chunk size = ?									

thread 1

0: end = ? 1: end = ?

thread 0

void parallel_loop(..., int tid, int num_threads)
{

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   // work based on x
}
```

• Example, 2 threads/cores, array of size 9

0	1	2	3	4	5	6	7	8

4

void parallel_loop(..., int tid, int num_threads)
{

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   // work based on x
}
```

• Example, 2 threads/cores, array of size 9

0	1	2	3	4	5	6	7	8

4

void parallel_loop(..., int tid, int num_threads)
{

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
if (tid == num_threads - 1) {
  end = SIZE;
}
for (int x = start; x < end; x++) {
  // work based on x
}
```
last thread gets more work

• Example, 2 threads/cores, array of size 9

0	1	2	3	4	5	6	7	8

4

```
void parallel_loop(..., int tid, int num_threads)
{
```

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
if (tid == num_threads - 1) {
  end = SIZE;
}
for (int x = start; x < end; x++) {
  // work based on x
}
```

• Example, 2 threads/cores, array of size 9

4

0	1	2	3	4	5	6	7	8

chunk_size = 4

$$0: \text{ start} = 0$$
 $1: \text{ start} = 1$
 $0: \text{ end} = 4$ $1: \text{ end} = 9$
thread 0
thread 1

last thread gets more work

What is the worst case?

```
void parallel_loop(..., int tid, int num_threads)
{
```

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
if (tid == num_threads - 1) {
  end = SIZE;
}
for (int x = start; x < end; x++) {
  // work based on x
}
```

End example

• Example, 2 threads/cores, array of size 9

0	1	2	3	4	5	6	7	8

```
chunk_size = 4
0: start = 0 1: start = 4
0: end = 4 1: end = 8
thread 0
thread 1
```

ceiling division, this will distribute uneven work in the last thread to all other threads

```
void parallel_loop(..., int tid, int num_threads)
{
    int chunk_size =
      (SIZE+(NUM_THREADS-1))/NUM_THREADS;
    int start = chunk_size * tid;
    int end = start + chunk_size;
    for (int x = start; x < end; x++) {
        // work based on x
    }
}</pre>
```

• Example, 2 threads/cores, array of size 9



• Example, 2 threads/cores, array of size 9



Example, 2 threads/cores, array of size 9

0	1	2	3	4	5	6	7	8

```
chunk_size = 5
0: start = 0 1: start = 5
0: end = 5 1: end = 9
thread 0
thread 1
```

most threads do equal amounts of work, last thread may do less.

Which one is better/worse? Max slowdown for last thread does all the extra work?

Max slowdown for ceiling?

```
void parallel_loop(..., int tid, int num_threads)
{
    int chunk_size =
    (SIZE+(NUM_THREADS-1))/NUM_THREADS;
    int start = chunk_size * tid;
    int end =
    min(start+chunk_size, SIZE)
for (int x = start; x < end; x++) {
      // work based on x
    }
}</pre>
```

End example

Schedule

- DOALL Loops
- Parallel Schedules:
 - Static
 - Global Worklists
 - Local Worklists

- Tasks are not balanced
- Appears in lots of emerging workloads

- Tasks are not balanced
- Appears in lots of emerging workloads

social network analytics where threads are parallel across users





- Tasks are not balanced
- Appears in lots of emerging workloads





- Independent iterations have different amount of work to compute
- Threads with longer tasks take longer to compute.
- Threads with shorter tasks are under utilized.

```
for (x = 0; x < SIZE; x++) {
   for (y = x; y < SIZE; y++) {
      a[x,y] = b[x,y] + c[x,y];
   }
}</pre>
```

irregular (or unbalanced) parallelism: each x iteration performs different amount of work.

- Calculate imbalance cost if x is chunked:
 - Thread 1 takes iterations 0 SIZE/2
 - Thread 2 takes iterations SIZE/2 SIZE

```
for (x = 0; x < SIZE; x++) {
  for (y = x; y < SIZE; y++) {
     a[x,y] = b[x,y] + c[x,y];
  }
}</pre>
```

- Calculate imbalance cost if x is chunked:
 - Thread 1 takes iterations 0 SIZE/2
 - Thread 2 takes iterations SIZE/2 SIZE

```
Calculate how much total work:
```

total_work =
$$\sum_{n=0}^{SIZE} n$$

```
for (x = 0; x < SIZE; x++) {
   for (y = x; y < SIZE; y++) {
      a[x,y] = b[x,y] + c[x,y];
   }
}</pre>
```

- Calculate imbalance cost if x is chunked:
 - Thread 1 takes iterations 0 SIZE/2
 - Thread 2 takes iterations SIZE/2 SIZE

```
Calculate how much total work:
```

total_work =
$$\sum_{n=0}^{SIZE} n$$

Calculate work done by first thread:

```
for (x = 0; x < SIZE; x++) {
   for (y = x; y < SIZE; y++) {
      a[x,y] = b[x,y] + c[x,y];
   }
}</pre>
```

$$t1_work = \sum_{n=0}^{SIZE/2} n$$

- Calculate imbalance cost if x is chunked:
 - Thread 1 takes iterations 0 SIZE/2
 - Thread 2 takes iterations SIZE/2 SIZE

```
total_work = \sum_{i=1}^{SIZE} n
```

Calculate work done by first thread:

Calculate how much total work:

```
for (x = 0; x < SIZE; x++) {
   for (y = x; y < SIZE; y++) {
      a[x,y] = b[x,y] + c[x,y];
   }
}</pre>
```

$$t1_work = \sum_{n=0}^{SIZE/2} n$$

Calculate work work done by second thread:

t2_work = total_work - t1_work

Calculate how much total work:

Example: SIZE = 64

total_work = 2016 t2_work = 496 t1_work = 1520

t1 does ~3x more work than t2

Only provides ~1.3x speedup

Potential solution: Have T1 do only ¼ of the iterations Gives a better speedup of 1.77x

Not a feasible solution because often times load imbalance is not given by a static equation on loop bounds!

total_work =
$$\sum_{n=0}^{SIZE} n$$

Calculate work done by first thread:

$$t1_work = \sum_{n=0}^{SIZE/2} n$$

Calculate work work done by second thread:

t2_work = total_work - t1_work

Work stealing

• Tasks are dynamically assigned to threads.

• Pros

- Simple to implement
- Cons:
 - High contention on global counter
 - Potentially bad memory locality.

• Global worklist: threads take tasks (iterations) dynamically

0	1	2	3	4	5	6	7		SIZE -1
---	---	---	---	---	---	---	---	--	---------

cannot color initially!



• Global worklist: threads take tasks (iterations) dynamically



• Global worklist: threads take tasks (iterations) dynamically

2	3	4	5	6	7		SIZE -1
---	---	---	---	---	---	--	---------



• Global worklist: threads take tasks (iterations) dynamically

2	3	4	5	6	7		SIZE -1
---	---	---	---	---	---	--	---------



• Global worklist: threads take tasks (iterations) dynamically



finished tasks

• Global worklist: threads take tasks (iterations) dynamically

3	4	5	6	7		SIZE -1
---	---	---	---	---	--	---------

1



• Global worklist: threads take tasks (iterations) dynamically

3	4	5	6	7		SIZE -1
---	---	---	---	---	--	---------



• Global worklist: threads take tasks (iterations) dynamically



• Global worklist: threads take tasks (iterations) dynamically

4	5	6	7		SIZE -1
---	---	---	---	--	---------





• Global worklist: threads take tasks (iterations) dynamically

4	5	6	7		SIZE -1
---	---	---	---	--	---------



• Global worklist: threads take tasks (iterations) dynamically



• Global worklist: threads take tasks (iterations) dynamically

5	6	7		SIZE -1
---	---	---	--	---------



End example

• How to implement

```
void foo() {
    ...
    for (x = 0; x < SIZE; x++) {
        // dynamic work based on x
    }
    ...
}</pre>
```

• How to implement

void foo() {
 ...
 for (x = 0; x < SIZE; x++) {
 // dynamic work based on x
 }
 ...
}</pre>

void parallel_loop(...) {

```
for (x = 0; x < SIZE; x++) {
    // dynamic work based on x
}</pre>
```

Replicate code in a new function. Pass all needed variables as arguments.

How to implement

void foo() {
 ...
 for (x = 0; x < SIZE; x++) {
 // dynamic work based on x
 // ...
}</pre>

atomic_int x(0); void parallel_loop(...) {
 for (x = 0; x < SIZE; x++) {
 // dynamic work based on x
 }
}</pre>

move loop variable to be a global atomic variable
• How to implement

void foo() {
 ...
 for (x = 0; x < SIZE; x++) {
 // dynamic work based on x
 }
}
atomic_int x(0);
void parallel_loop(...) {
 for (int local_x = ??
 local_x < SIZE;
 local_x = ??) {
 // dynamic work based on x
 }
}</pre>

change loop bounds in new function to use a local variable using global variable.

• How to implement

These must be atomic updates!

void foo() {



change loop bounds in new function to use a local variable using global variable.

```
• How to implement
```

```
void foo() {
    ...
    for (t = 0; x < THREADS; t++) {
        spawn(parallel_loop);
    }
    join();
    ...
}</pre>
```

```
atomic_int x(0);
void parallel_loop(...) {
```

```
// dynamic work based on x
```

Spawn threads in original function and join them afterwards

You will have to change to C++ syntax for the homework!

• Global worklist: threads take tasks (iterations) dynamically

0	1	2	3	4	5	6	7	SIZE -1

thread 1

thread 0

```
x: 0
0 - local_x - UNDEF
1 - local_x - UNDEF
```

• Global worklist: threads take tasks (iterations) dynamically



• Global worklist: threads take tasks (iterations) dynamically

2	3	4	5	6	7	SIZE -1

```
// dynamic work based on x
```

}

}



• Global worklist: threads take tasks (iterations) dynamically

}

2	3	4	5	6	7	SIZE -1

0

thread 0

thread 1

• Global worklist: threads take tasks (iterations) dynamically



• Global worklist: threads take tasks (iterations) dynamically

3	4	5	6	7	SIZE -1

}



• Global worklist: threads take tasks (iterations) dynamically

					1	
3	4	5	6	7		SIZE -1

}

}





• Global worklist: threads take tasks (iterations) dynamically



• Global worklist: threads take tasks (iterations) dynamically

4	5	6	7	SIZE -1



• Global worklist: threads take tasks (iterations) dynamically

	·	·		
4	5	6	7	SIZE -1

}



3

thread 1

• Global worklist: threads take tasks (iterations) dynamically



• Global worklist: threads take tasks (iterations) dynamically

5	6	7	SIZE -1

```
      4
      3

      atomic_int x(0);
void parallel_loop(...) {

      for (int local_x = atomic_fetch_add(&x,1);
local_x < SIZE;
local_x = atomic_fetch_add(&x,1)) {

      thread 0
      // dynamic work based on x
}
```

Schedule

• DOALL Loops

• Parallel Schedules:

- Static
- Global Worklists
- Local Worklists

- More difficult to implement: typically requires concurrent datastructures
- low contention on local data-structures
- potentially better cache locality

• local worklists: divide tasks into different worklists for each thread

0	1	2	3
---	---	---	---

thread 0 thread 1















• local worklists: divide tasks into different worklists for each thread

Queue 0





• How to implement:

```
void foo() {
    ...
    for (x = 0; x < SIZE; x++) {
        // dynamic work based on x
    }
    ...
}</pre>
```

• How to implement:

void foo() {
 ...
 for (x = 0; x < SIZE; x++) {
 // dynamic work based on x
 }
 ...
 }
}</pre>
void parallel_loop(..., int tid) {
 for (x = 0; x < SIZE; x++) {
 // dynamic work based on x
 }
 ...
}

Make a new function, taking any variables used in loop body as args. Additionally take in a thread id

• How to implement:

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
...
for (x = 0; x < SIZE; x++) {
    // dynamic work based on x
}
</pre>
```

```
void parallel_loop(..., int tid) {
```

```
for (x = 0; x < SIZE; x++) {
   // dynamic work based on x
}</pre>
```

Make a global array of concurrent queues

• How to implement:

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
...
_____for (x = 0; x < SIZE; x++) {
____// dynamic work based on x
_____}
...
</pre>
```

```
void parallel_loop(..., int tid) {
```

```
for (x = 0; x < SIZE; x++) {
   // dynamic work based on x
}</pre>
```

What type of queues?

• How to implement:

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
void parallel_loop(..., int tid) {
```

```
for (x = 0; x < SIZE; x++) {
   // dynamic work based on x
}</pre>
```

What type of queues? We're going to use InputOutput Queues!

Make a global array of concurrent queues









concurrent_queues cq[NUM_THREADS];
void foo() {

. . .

. . .

First we need to initialize the queues
```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
• • •
```

. . .

// Spawn threads to initialize
// join initializing threads

```
void parallel_enq(..., int tid, int num_threads)
{
```

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   cq[tid].enq(x);
}
```

Just like the static schedule, except we are enqueuing

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
•••
```

. . .

// Spawn threads to initialize
// join initializing threads

Make sure to account for boundary conditions!

```
void parallel_enq(..., int tid, int num_threads)
{
```

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   cq[tid].enq(x);
}
```

Just like the static schedule, except we are enqueuing

• How to implement in a compiler:



```
Make sure to account for boundary conditions!
```

```
void parallel_enq(..., int tid, int num_threads)
{
```

```
int chunk_size = SIZE / NUM_THREADS;
int start = chunk_size * tid;
int end = start + chunk_size;
for (int x = start; x < end; x++) {
   cq[tid].enq(x);
}
```

Just like the static schedule, except we are enqueuing

```
concurrent_queues cq[NUM_THREADS]; voi
void foo() {
    ...
```

```
// initialize queues
```

// join threads

. . .

```
// launch loop function
```

```
void parallel_loop(..., int tid, int num_threads) {
  for (x = 0; x < SIZE; x++) {
    // dynamic work based on x
  }
}</pre>
```

How do we modify the parallel loop?

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
...
// initialize
```

. . .

```
// initialize queues
// initialize dueues
```

```
// join threads
```

```
// launch loop function
```

```
void parallel_loop(..., int tid, int num_threads) {
    int task = 0;
    for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    {
      // dynamic work based on task
    }
}
```

loop until the queue is empty

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
// initialize queues
```

// join threads

. . .

```
// launch loop function
```

```
void parallel_loop(..., int tid, int num_threads) {
    int task = 0;
    for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    {
      // dynamic work based on task
    }
}
```

loop until the queue is empty Are we finished?

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
// initialize queues
```

```
// join threads
```

. . .

```
// launch loop function
```

```
atomic_int finished_threads(0);
void parallel_loop(..., int tid, int num_threads) {
    int task = 0;
    for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    {
        // dynamic work based on task
    }
    atomic_fetch_add(&finished_threads,1);
}
```

Track how many threads are finished

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
• • •
```

. . .

```
// initialize queues
```

// join threads

```
// launch loop function
```

```
atomic_int finished_threads(0);
void parallel_loop(..., int tid, int num_threads) {
```

```
int task = 0;
for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
{
    // dynamic work based on task
}
atomic_fetch_add(&finished_threads,1);
while (finished_threads.load() != num_threads) {
```

While there are threads that are still working

}

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
• • •
```

. . .

```
// initialize queues
```

// join threads

```
// launch loop function
```

```
atomic_int finished_threads(0);
void parallel_loop(..., int tid, int num_threads) {
```

```
int task = 0;
for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
{
    // dynamic work based on task
}
atomic_fetch_add(&finished_threads,1);
while (finished_threads.load() != num_threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
  }
}
```

pick a random target and steal a task

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
• • •
```

. . .

```
// initialize queues
```

// join threads

```
// launch loop function
// join loop threads
```

```
atomic_int finished_threads(0);
void parallel_loop(..., int tid, int num_threads) {
```

```
int task = 0;
for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
{
    // dynamic work based on task
}
atomic_fetch_add(&finished_threads,1);
while (finished_threads.load() != num_threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
    if (task != -1) {
        // perform task
    }
}
```

```
concurrent_queues cq[NUM_THREADS];
void foo() {
```

```
• • •
```

. . .

```
// initialize queues
```

// join threads

// launch loop function
// join loop threads

```
join the threads
```

```
atomic_int finished_threads(0);
void parallel_loop(..., int tid, int num_threads) {
```

```
int task = 0;
for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
{
    // dynamic work based on task
}
atomic_fetch_add(&finished_threads,1);
while (finished_threads.load() != num_threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
    if (task != -1) {
        // perform task
    }
}
```



```
atomic_int finished_threads(0);
void parallel_loop(..., int tid, int num_threads) {
```

```
int task = 0;
for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
{
    // dynamic work based on task
}
atomic_fetch_add(&finished_threads,1);
while (finished_threads.load() != num_threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
    if (task != -1) {
        // perform task
    }
}
```



```
atomic_int finished_threads(0);
void parallel_loop(..., int tid, int num_threads) {
```

```
int task = 0;
for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
{
    // dynamic work based on task
}
atomic_fetch_add(&finished_threads,1);
while (finished_threads.load() != num_threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
    if (task != -1) {
        // perform task
    }
}
```

























```
atomic_int finished threads(0);
void parallel loop(..., int tid, int num threads) {
  int task = 0;
  for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    // dynamic work based on task
  atomic fetch add(&finished threads,1);
 while (finished_threads.load() != num_threads) {
    int target = // pick a thread to steal from
   int task = cq[target].deq();
    if (task != -1) {
       // perform task
    }
  }
```



IOQueue 0

```
IOQueue 1
```



```
void parallel_loop(..., int tid, int num_threads) {
    int task = 0;
    for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    {
        // dynamic work based on task
    }
    atomic_fetch_add(&finished_threads,1);
    while (finished_threads.load() != num_threads) {
        int target = // pick a thread to steal from
        int task = cq[target].deq();
        if (task != -1) {
            // perform task
        }
    }
}
```

atomic int finished threads(0);



IOQueue 0

thread 0



```
1
thread 1
```

```
atomic int finished threads(0);
void parallel loop(..., int tid, int num_threads) {
  int task = 0;
  for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    // dynamic work based on task
  atomic fetch add(&finished threads,1);
  while (finished threads.load() != num threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
    if (task != -1) {
       // perform task
    }
  }
}
```





finished_threads: 2

IOQueue 0

IOQueue 1

```
thread 0
```



```
atomic int finished threads(0);
void parallel loop(..., int tid, int num_threads) {
  int task = 0;
  for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    // dynamic work based on task
  atomic fetch add(&finished threads,1);
 while (finished threads.load() != num_threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
    if (task != -1) {
       // perform task
    }
  }
```

```
finished_threads: 2
```

IOQueue 0

IOQueue 1

thread 0



```
atomic int finished threads(0);
void parallel loop(..., int tid, int num_threads) {
  int task = 0;
  for (x = cq[tid].deq(); x != -1; x = cq[tid].deq())
    // dynamic work based on task
  atomic fetch add(&finished threads,1);
  while (finished threads.load() != num threads) {
    int target = // pick a thread to steal from
    int task = cq[target].deq();
    if (task != -1) {
       // perform task
    }
  }
```

Thanks!

- Work on midterm!
 - Ask on piazza if you have questions or comments
- Homework 3 is out
- Do the quiz!