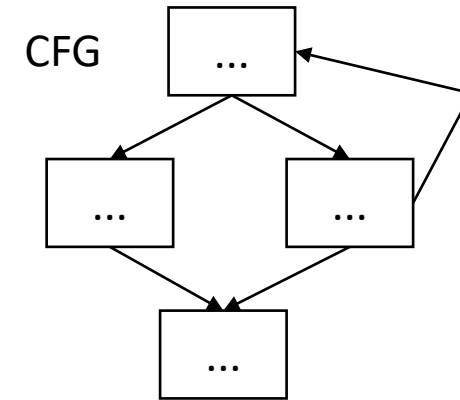
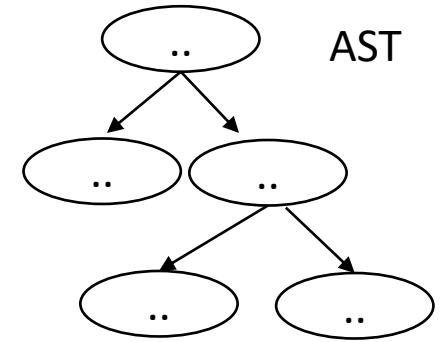


CSE110A: Compilers

May 10, 2023

Topics:

- *ASTs*
 - *parse trees into ASTs*
 - *type checking*



3 address code

```
store i32 0, ptr %2
%3 = load i32, ptr %1
%4 = add nsw i32 %3, 1,
store i32 %4, ptr %1
%5 = load i32, ptr %2
```

Announcements

- Midterm is over!
 - We will start grading early next week and aim to scores released by next Friday

Announcements

- HW 3
 - Released on Monday
 - 7 days to do the homework (due May 15)
 - Implementing a recursive descent parser for your grammar
 - Implementing a symbol table
 - You will need to implement a scanner
 - Use your HW 1 solution or the exact match scanner (given in HW 1)
 - You will need a grammar
 - Either use your grammar or the provided grammar (Rithik provided on Piazza)
- Come to office hours or post on Piazza for help!

Announcements

- HW 3 clarification: You do not need to return anything from your parser!
 - If the input program satisfies the grammar then you return without issue
 - If it does not, then you throw an exception
 - Scanner exception if you cannot create a token
 - Parser exception if the input violates the grammar
 - Symbol table exception if a variable is used outside of a scope it is declared
- HW 4 will be creating an IR inside your parser.

Quiz

Quiz

Parse tree is an Abstract Syntax Tree

☐ True

☐ False

Quiz

If you are writing a compiler on M languages for N target architectures. How many compilers will you need to write with and without the help of Intermediate Representation?

☐ M, N

☐ $MN, M+N$

☐ $M+N, MN$

☐ MN, NM

☐ M, NM

☐ $M, N + M$

Quiz

Loop unrolling will ____ loop overhead and ____ program code size

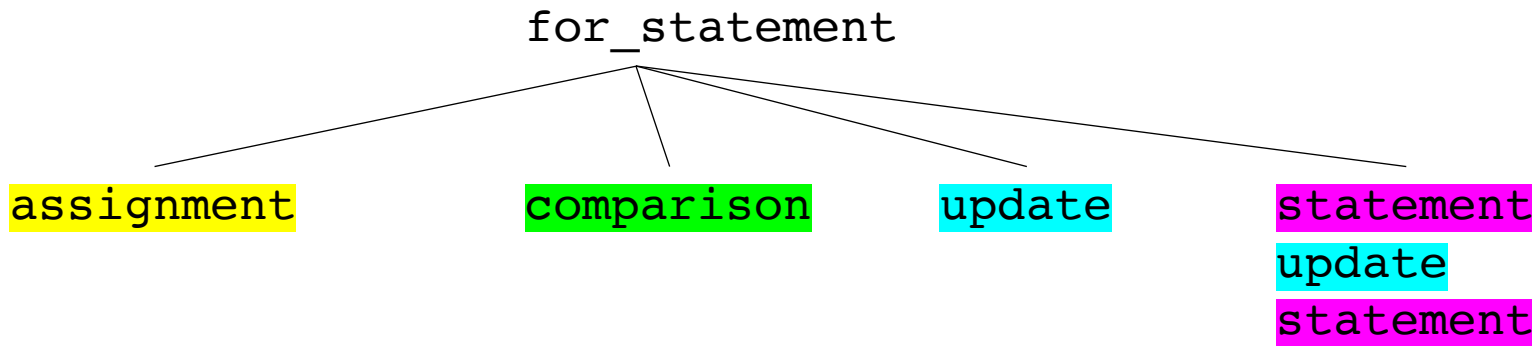
☐ increase, increase

☐ increase, reduce

☐ reduce, increase

☐ reduce, reduce

Example: loop unrolling



```
for (i = 0; i < 100; i = i + 1) {  
    x = x + 1;  
}
```

Check:

1. Find iteration variable by examining assignment, comparison and update.

2. found i

3. check that statement doesn't change i.

4. check that comparison goes around an even number of times.

Perform optimization

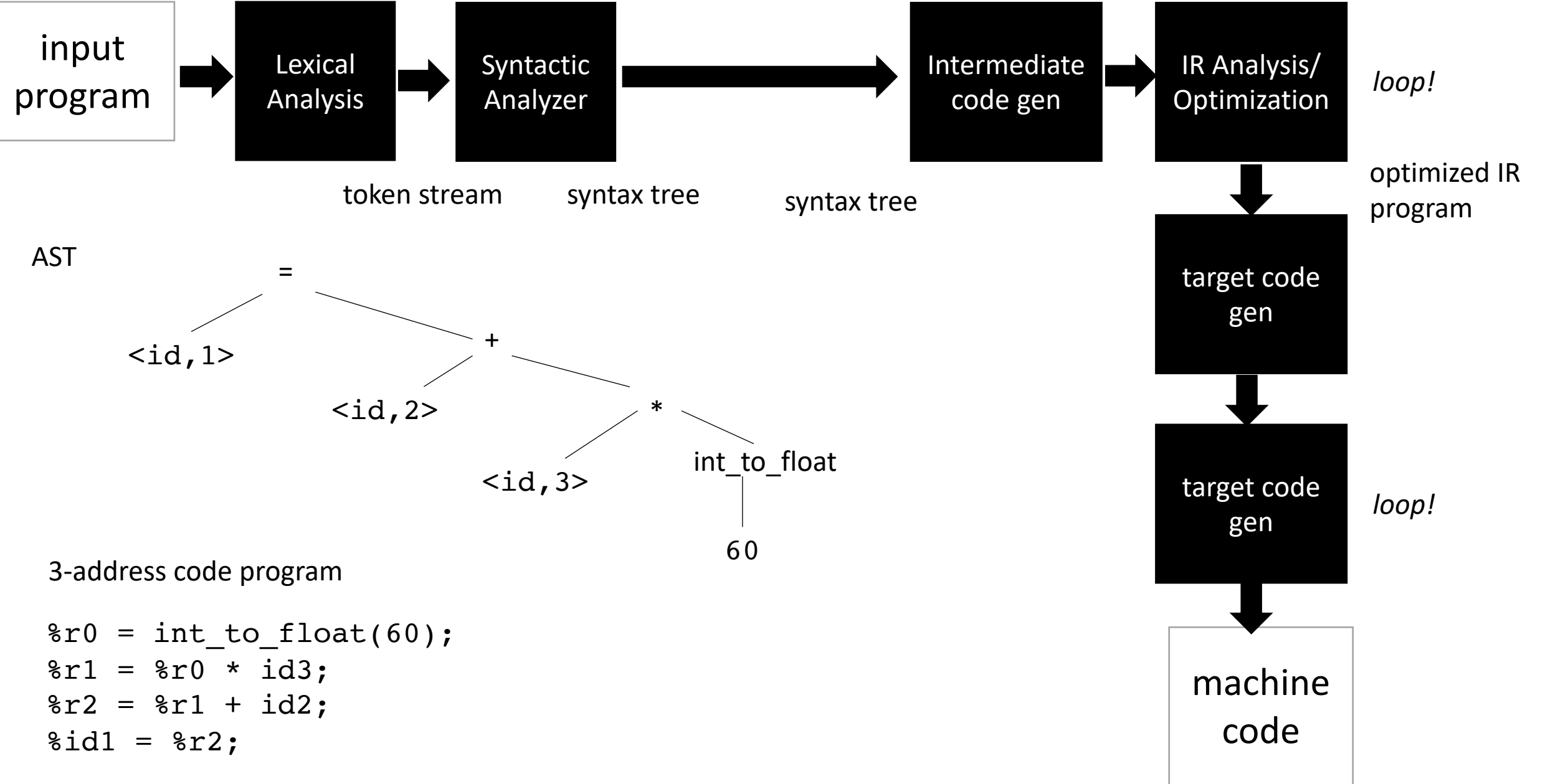
copy statement and put an update before it

Quiz

Name a few Intermediate Representations you have seen in real life

Review

```
position = initial + rate * 60;
```



Intermediate representations

- Several forms:
 - tree - abstract syntax tree
 - graphs - control flow graph
 - linear program - 3 address code
- Often times the program is represented as a hybrid
 - graphs where nodes are a linear program
 - linear program where expressions are ASTs
- Progression:
 - start close to a parse tree
 - move closer to an ISA

Our first IR: abstract syntax tree

- One step away from parse trees
- Great representation for expressions
- Natural representation to apply type checking

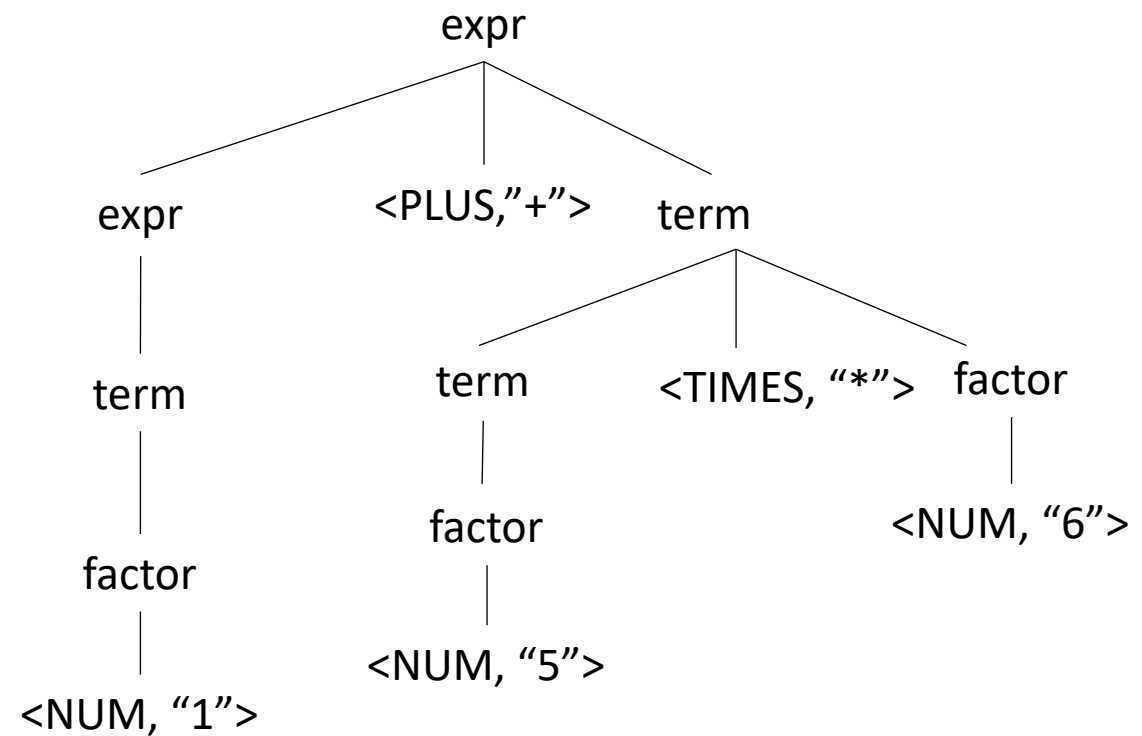
What is an AST?

Parse trees are defined **entirely** by the grammar

- **Tokens**
- **Production rules**

Parse trees are often not explicitly constructed. We use them to visualize the parsing computation

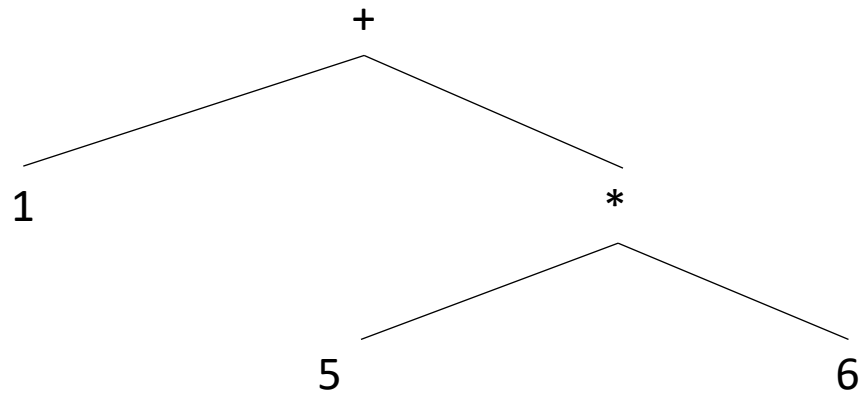
input: 1+5*6



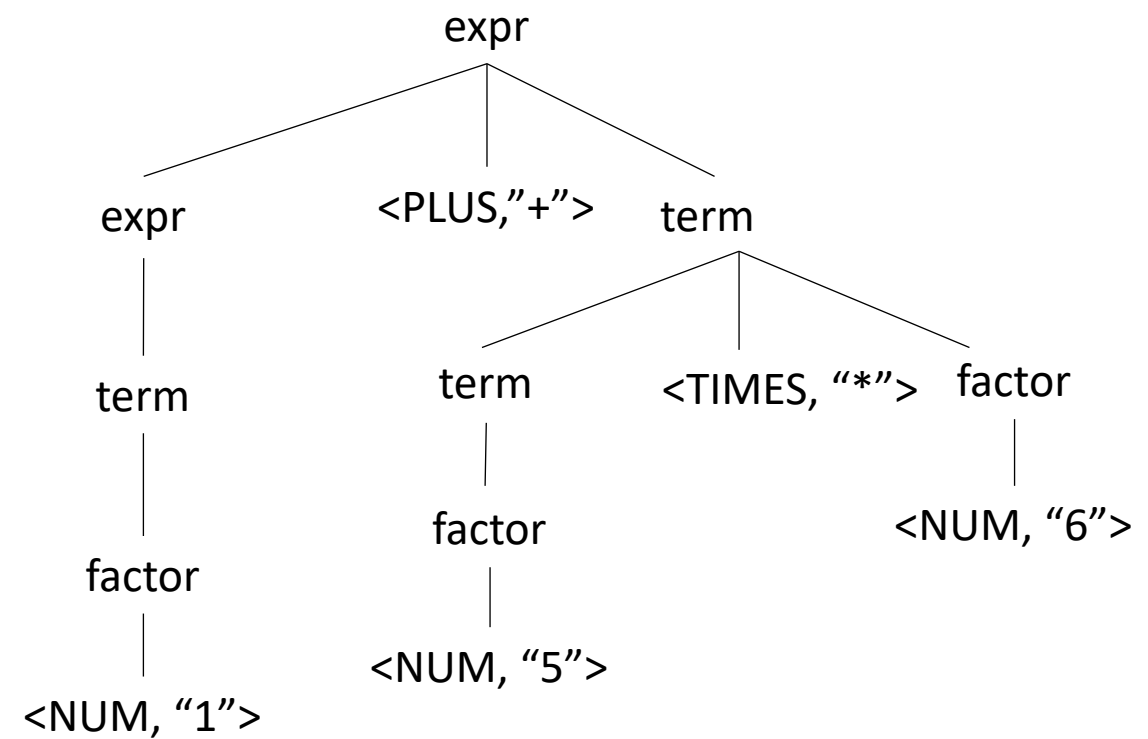
input: 1+5*6

What is an AST?

What are some differences?



AST

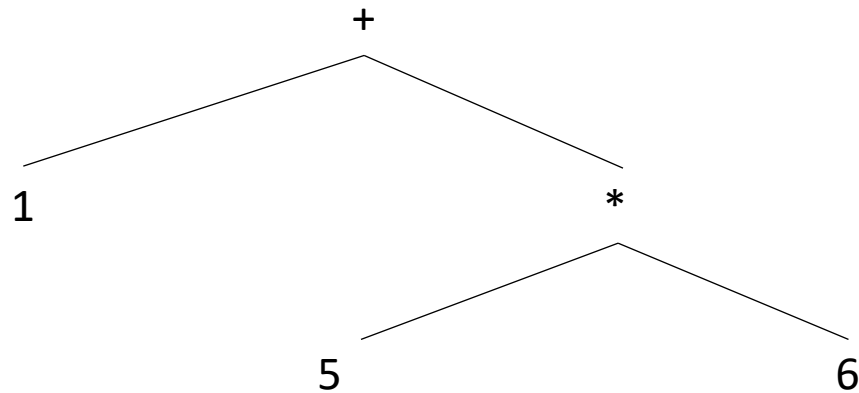


Parse Tree

input: 1+5*6

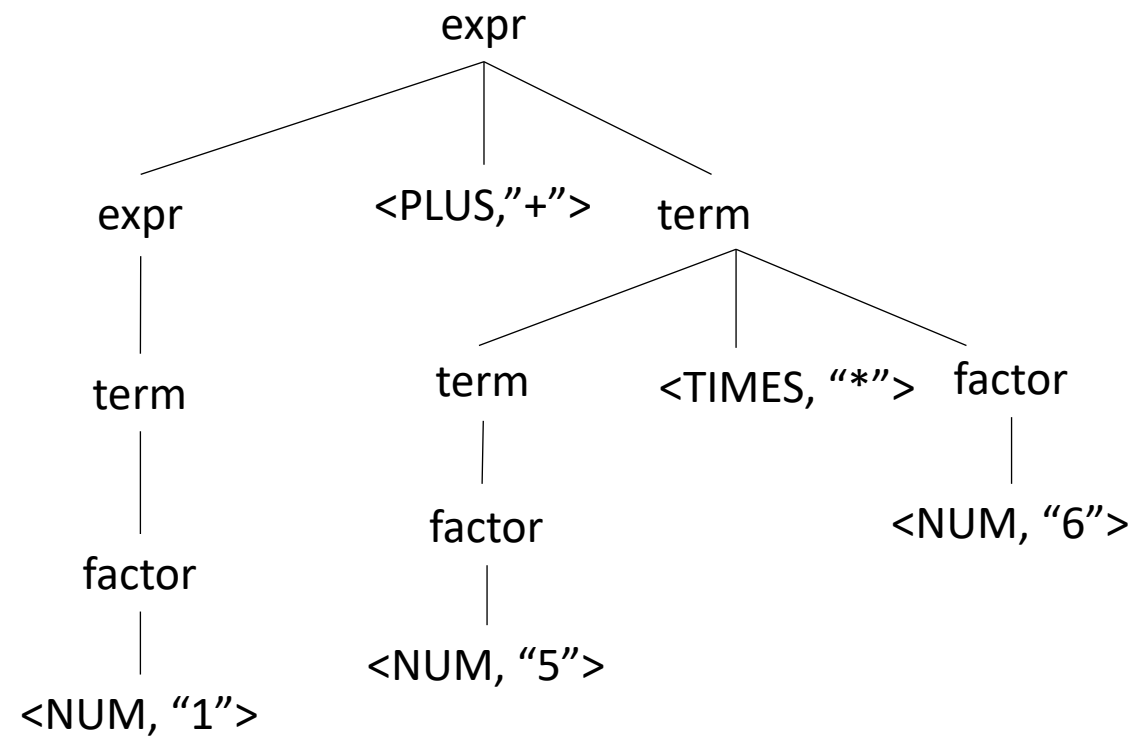
What is an AST?

What are some differences?



AST

- decoupled from the grammar
- leaves are data, not lexemes
- nodes are operators, not non-terminals

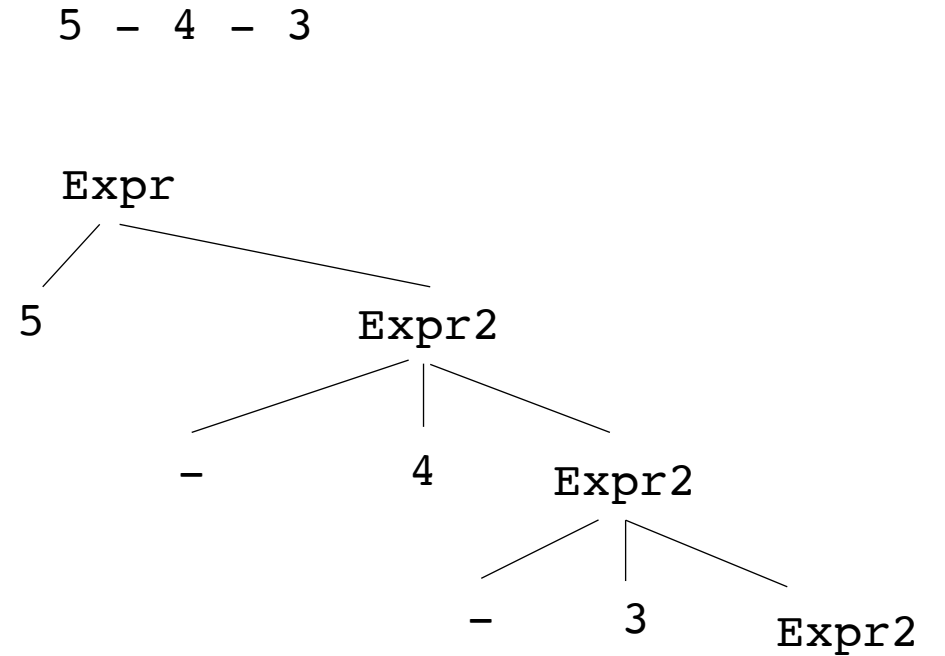


Parse Tree

Creating an AST from predictive grammar

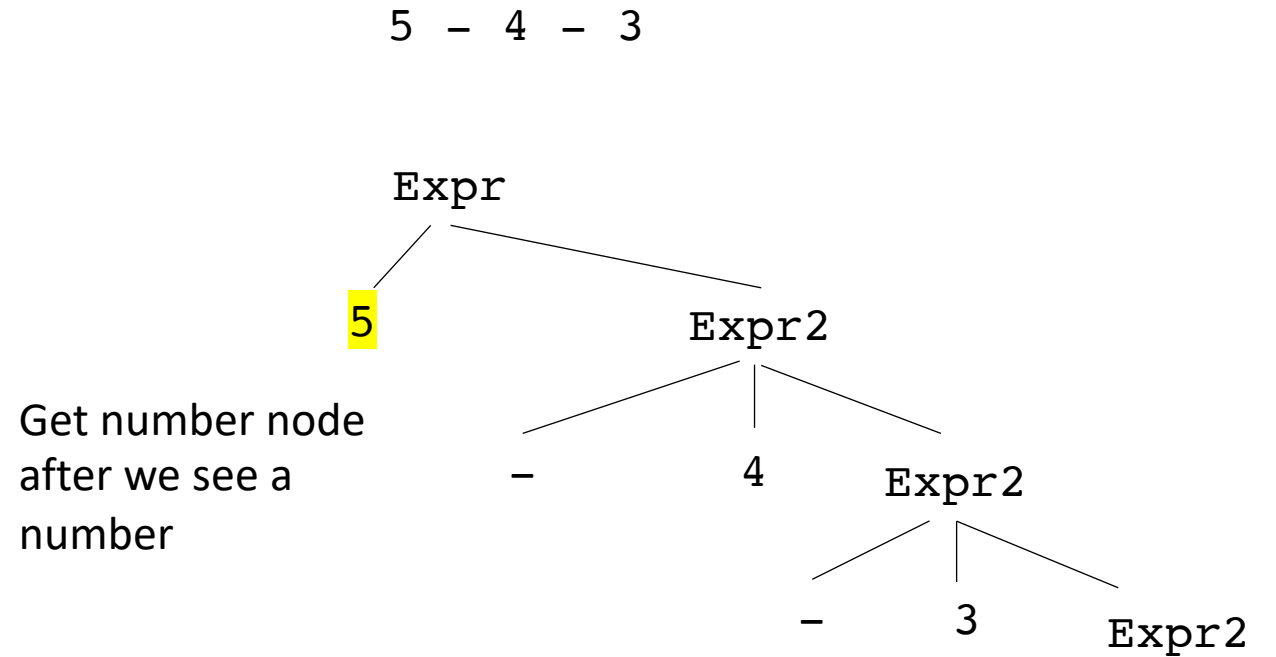
```
Expr  ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```

Keep in mind that because we wrote our own parser, we can inject code at any point during the parse.



Creating an AST from predictive grammar

<code>Expr</code>	<code>::=</code>	<code>NUM</code>	<code>Expr2</code>
<code>Expr2</code>	<code>::=</code>	<code>MINUS</code>	<code>NUM</code> <code>Expr2</code>
			<code>""</code>

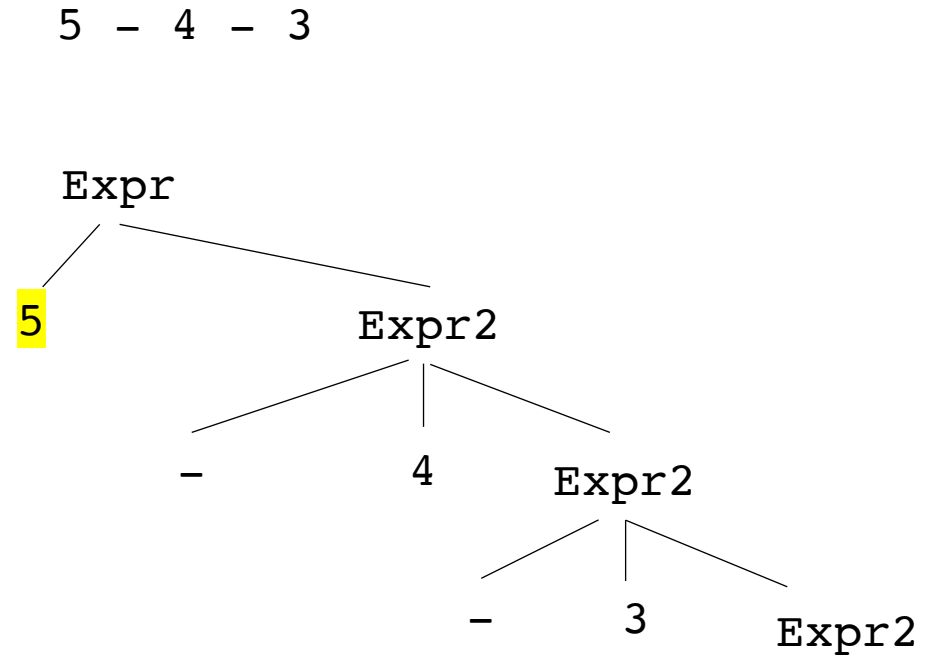


AST<5>

Creating an AST from predictive grammar

Expr	::=	NUM	Expr2
Expr2	::=	MINUS	NUM Expr2
		" "	

Pass the node
down



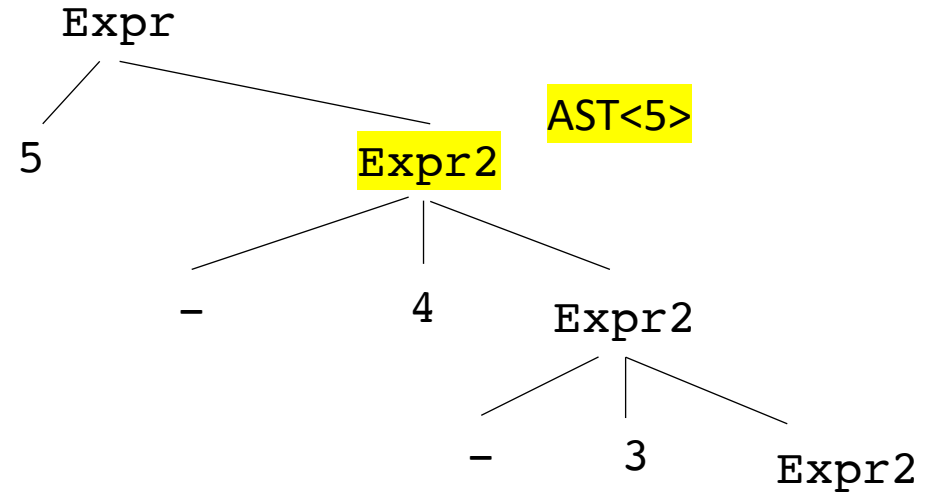
AST<5>

Creating an AST from predictive grammar

```
Expr  ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```

5 - 4 - 3

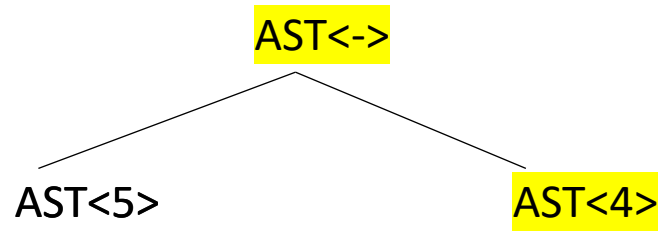
Pass the node
down



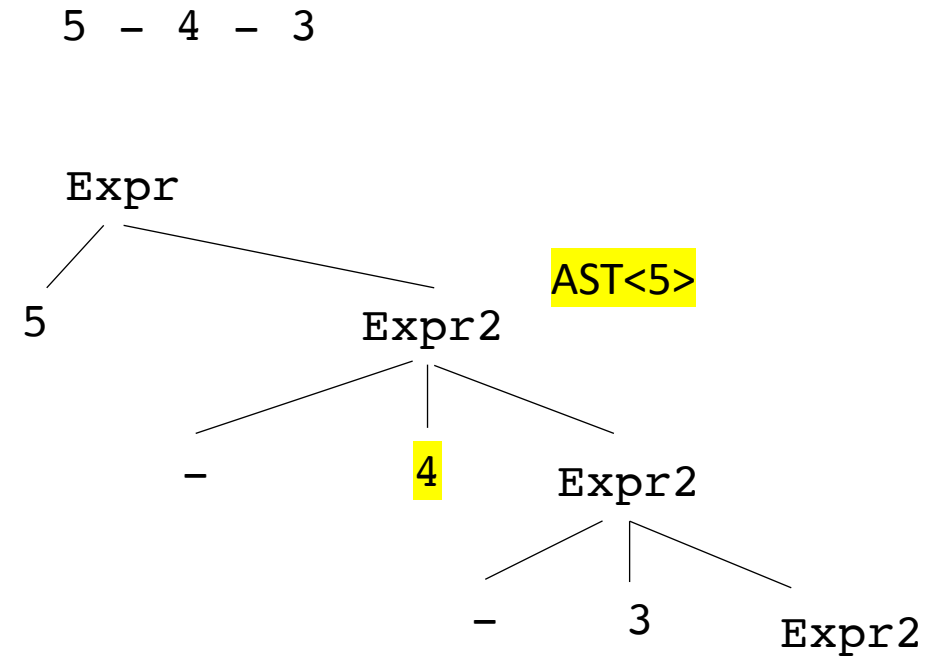
AST<5>

Creating an AST from predictive grammar

```
Expr  ::= NUM Expr2
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      | ""
```

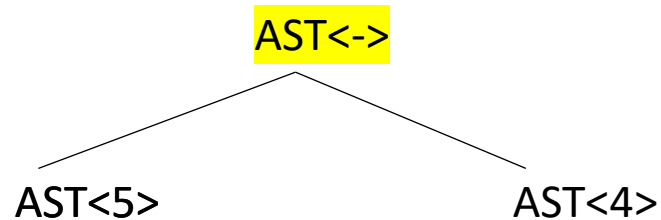


In Expr2, after 4 is
parsed, create a
number node and
a minus node

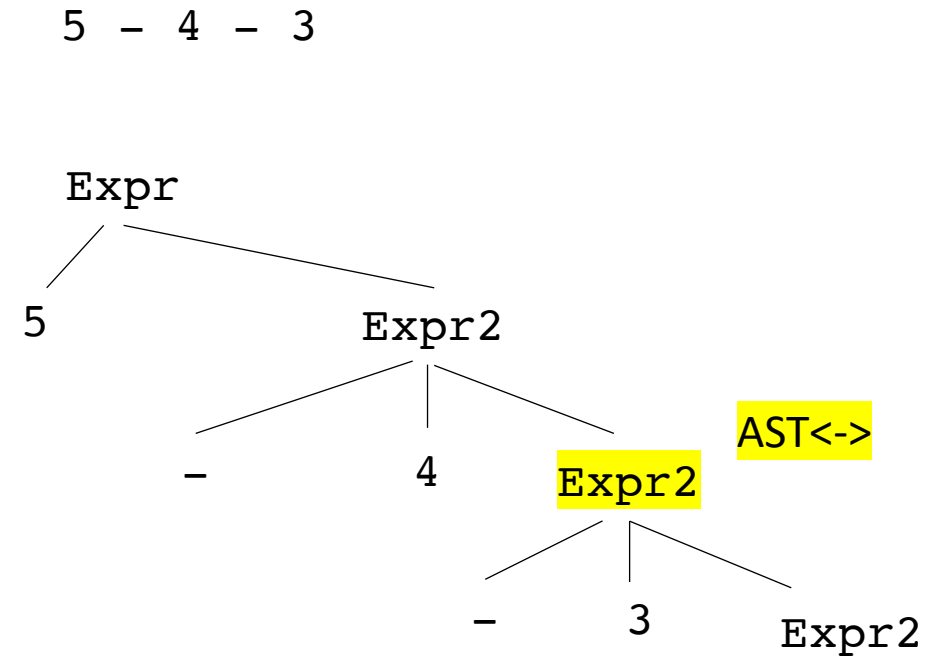


Creating an AST from predictive grammar

```
Expr  ::= NUM Expr2
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      | ""
```

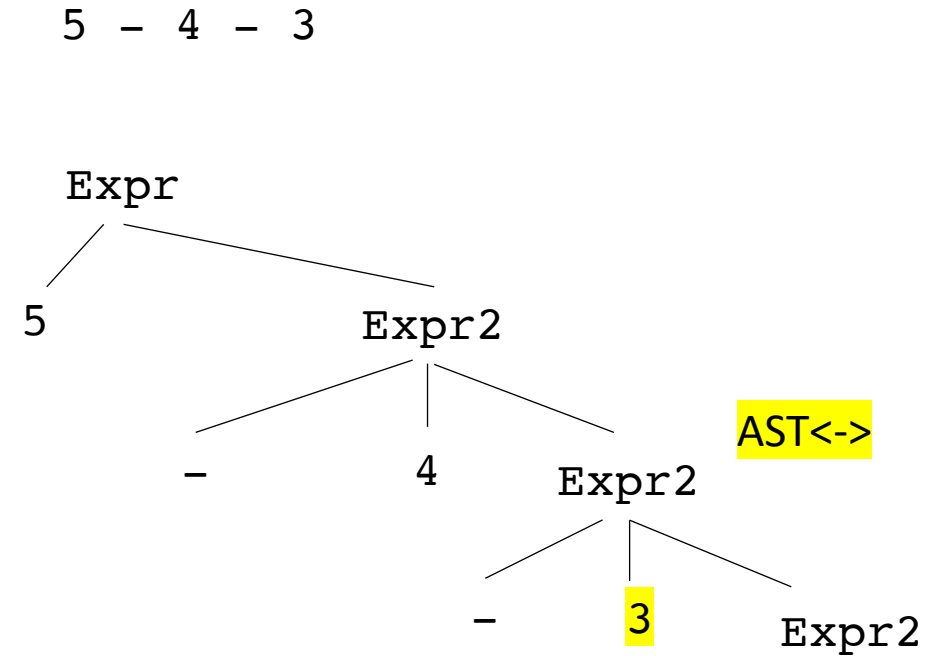
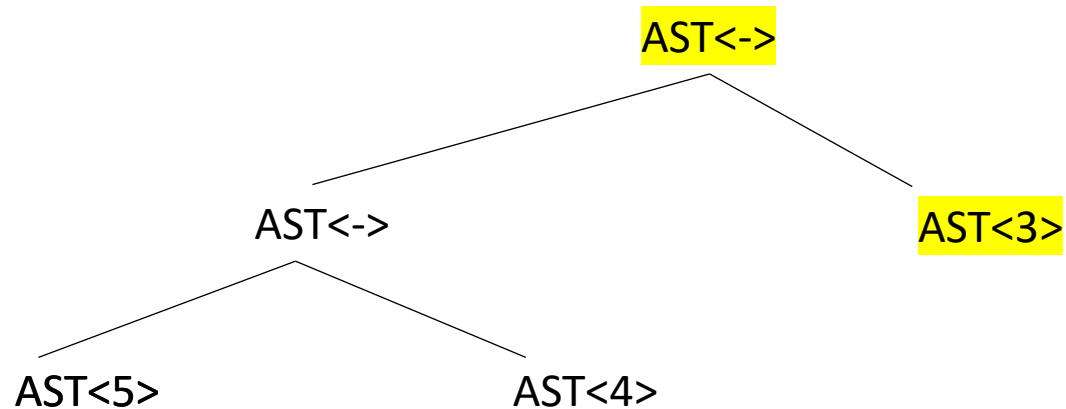


pass the new node
down



Creating an AST from predictive grammar

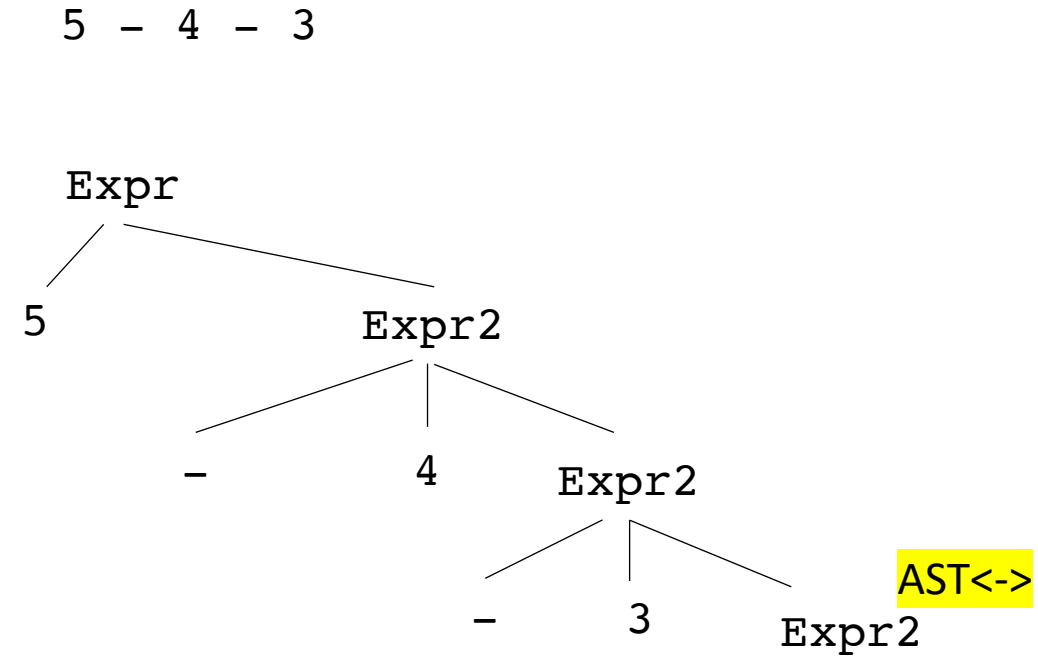
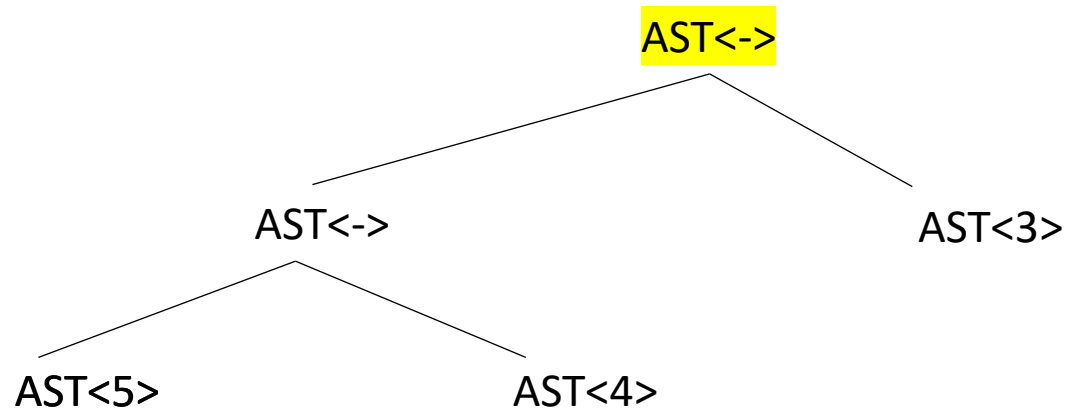
```
Expr  ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```



In Expr2, after 3 is
parsed, create a
number node and
a minus node

Creating an AST from predictive grammar

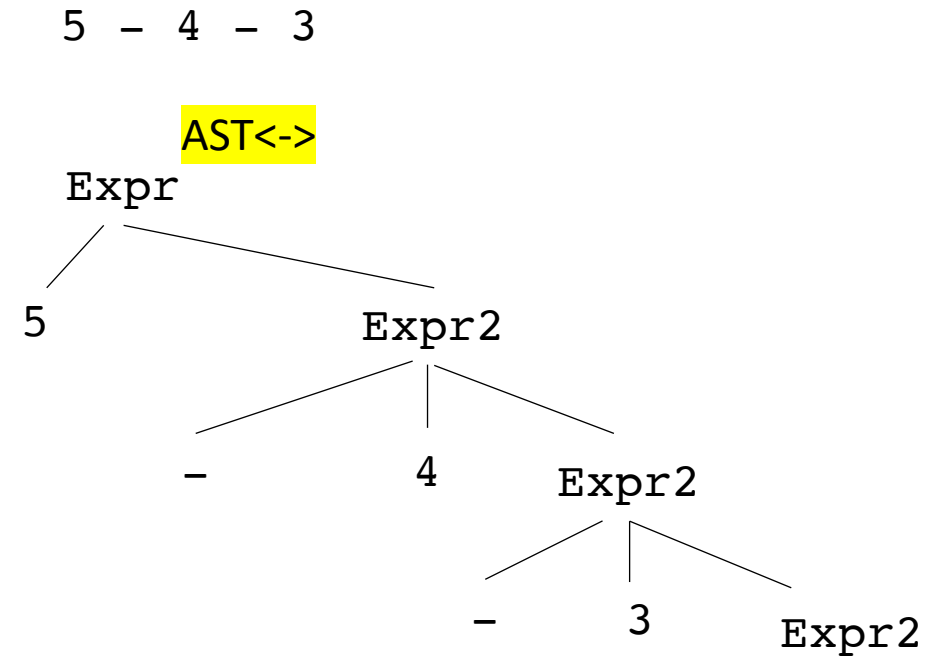
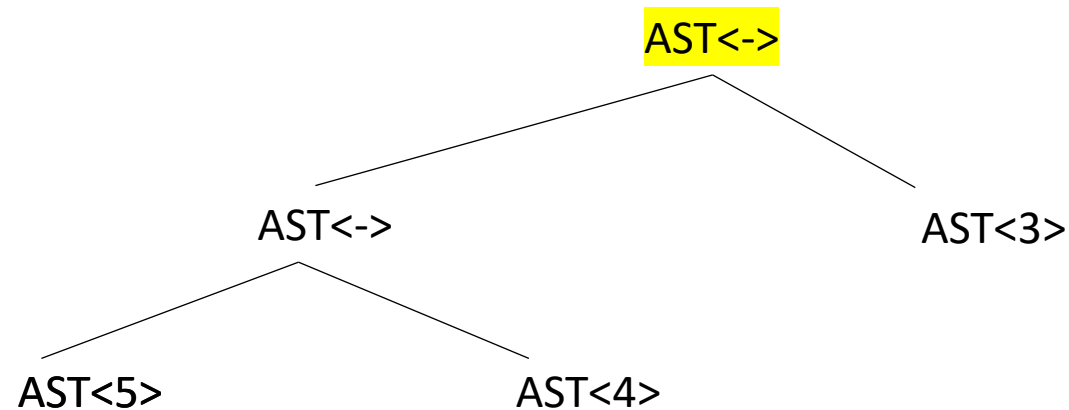
```
Expr  ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```



pass down the new
node

Creating an AST from predictive grammar

Expr	::=	NUM	Expr2
Expr2	::=	MINUS	NUM Expr2
		" "	



return the node
when there is
nothing left to
parse

Creating an AST from predictive grammar

```
Expr    ::= NUM Expr2
Expr2   ::= MINUS NUM Expr2
        |      ""
```

```
def parse_expr(self):
    #lexemes second field is the value
    value = self.next_word.value
    node = ASTNumNode(value)
    self.eat(Token.NUM)
    return self.parse_expr2(node)
```

Creating an AST from predictive grammar

```
Expr    ::= NUM Expr2
Expr2   ::= MINUS NUM Expr2
        |      ""
```

```
def parse_expr(self):
    #lexemes second field is the value
    value = self.next_word.value
    node = ASTNumNode(value)
    self.eat(Token.NUM)
    return self.parse_expr2(node)
```

```
def parse_expr2(self, lhs_node):
    # ... for applying the first production rule
    self.eat(Token.MINUS)
    value = self.next_word.value
    rhs_node = ASTNumNode(value)
    self.eat(Token.NUM)
    node = ASTMinusNode(lhs_node, rhs_node)
    return self.parse_expr2(node)
```

Creating an AST from predictive grammar

```
Expr    ::= NUM Expr2
Expr2   ::= MINUS NUM Expr2
        |      ""
```

```
def parse_expr(self):
    #lexemes second field is the value
    value = self.next_word.value
    node = ASTNumNode(value)
    self.eat(Token.NUM)
    return self.parse_expr2(node)
```

```
def parse_expr2(self, lhs_node):
    # ... for applying the second production rule
    return lhs_node
```

Creating an AST from predictive grammar

```
Expr    ::= Term Expr2
Expr2   ::= MINUS Term Expr2
        |      ""
```

In a more realistic grammar, you might have more layers: e.g. a **Term**

how to adapt?

```
def parse_expr(self):
    #lexemes second field is the value
    value = self.next_word.value
    node = ASTNumNode(value)
    self.eat(Token.NUM)
    return self.parse_expr2(node)
```

```
def parse_expr2(self, lhs_node):
    # ... for applying the first production rule
    self.eat(Token.MINUS)
    value = self.next_word.value
    rhs_node = ASTNumNode(value)
    self.eat(Token.NUM)
    node = ASTMinusNode(lhs_node, rhs_node)
    return self.parse_expr2(node)
```

Creating an AST from predictive grammar

```
Expr    ::= Term Expr2
Expr2   ::= MINUS Term Expr2
        |      ""
```

```
def parse_expr(self):
    node = self.parse_term()
    return self.parse_expr2(node)
```

In a more realistic grammar, you might have more layers: e.g. a **Term**

how to adapt?

```
def parse_expr2(self, lhs_node):
    # ... for applying the first production rule
    self.eat(Token.MINUS)
    rhs_node = self.parse_term()
    node = ASTMinusNode(lhs_node, rhs_node)
    return self.parse_expr2(node)
```

The `parse_term` will figure out how to get you an AST node for that term.

Example

- Python AST

```
import ast
```

```
print(ast.dump(ast.parse('5-4-2')))
```


Example

- Python AST

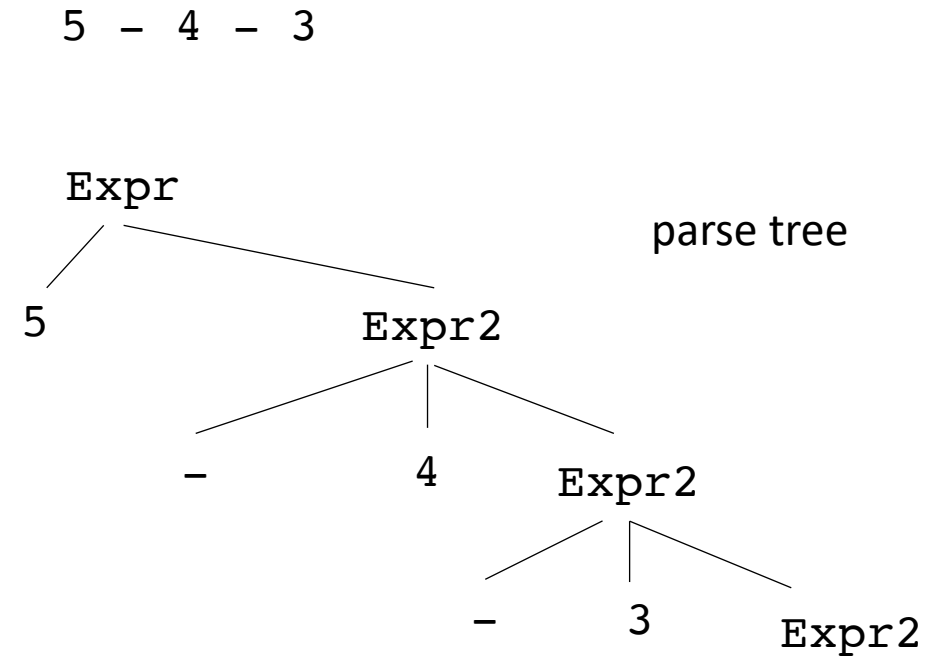
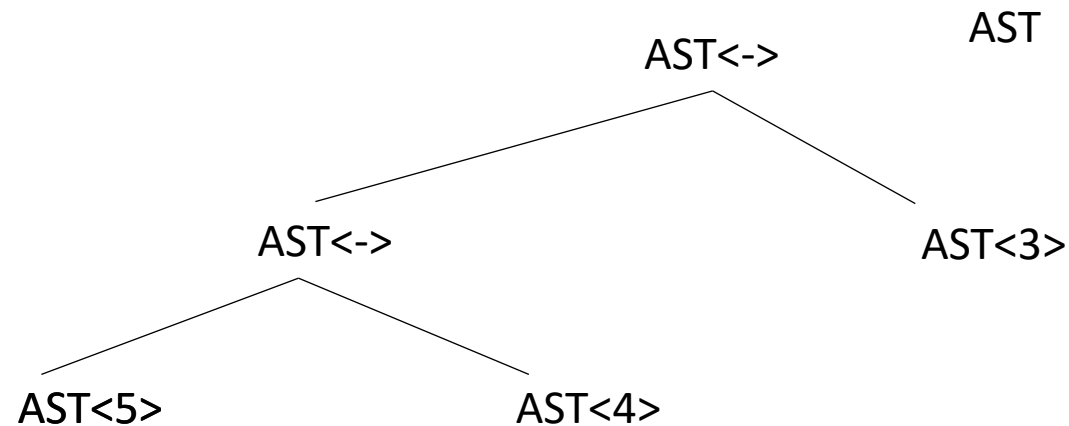
```
import ast
```

```
print(ast.dump(ast.parse('5-4-2')))
```

```
Expr(value=BinOp(left=BinOp(left=Num(n=5), op=Sub(), right=Num(n=4)), op=Sub(), right=Num(n=2)))
```

Evaluate an AST by doing a post order traversal

Expr	::=	NUM	Expr2
Expr2	::=	MINUS	NUM Expr2
			""



Parse trees cannot always be evaluated in post-order. An AST should always be

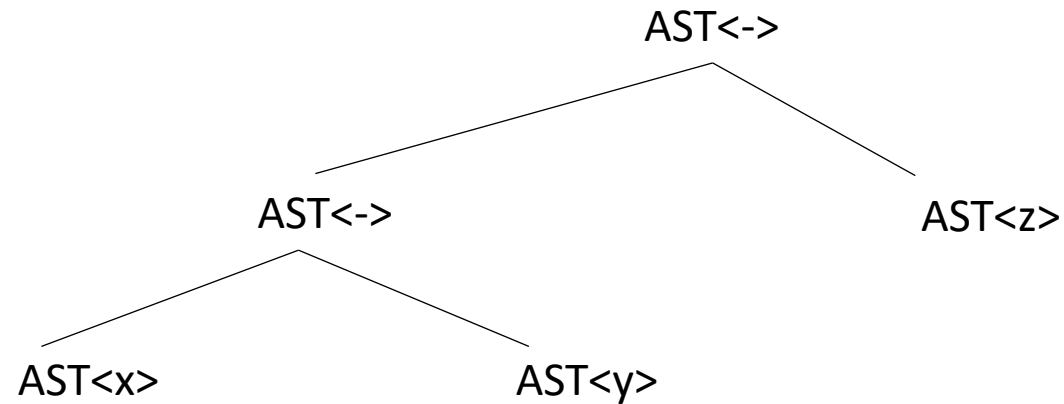
Evaluate an AST by doing a post order traversal

```
Expr ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```

What if you cannot evaluate it?

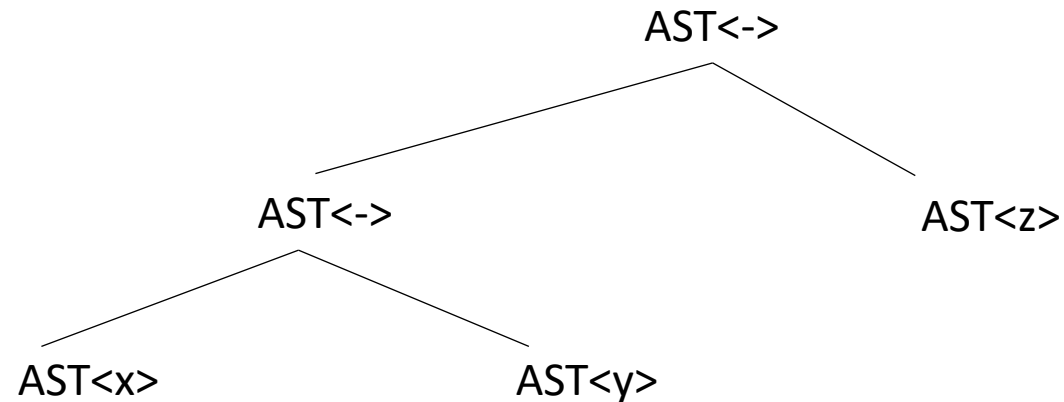
What else might you do?

x - y - z



Evaluate an AST by doing a post order traversal

```
Expr ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```



What if you cannot evaluate it?

What else might you do?

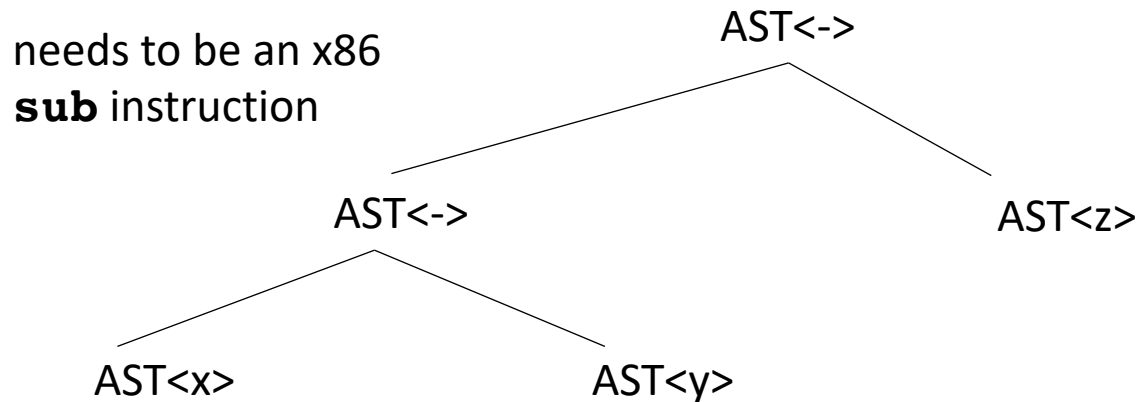
```
int x;
int y;
float z;
float w;
w = x - y - z
```

How does this change things?

Evaluate an AST by doing a post order traversal

```
Expr ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```

needs to be an x86
subss instruction



What if you cannot evaluate it?
What else might you do?

```
int x;
int y;
float z;
float w;
w = x - y - z
```

How does this change things?

Is this all?

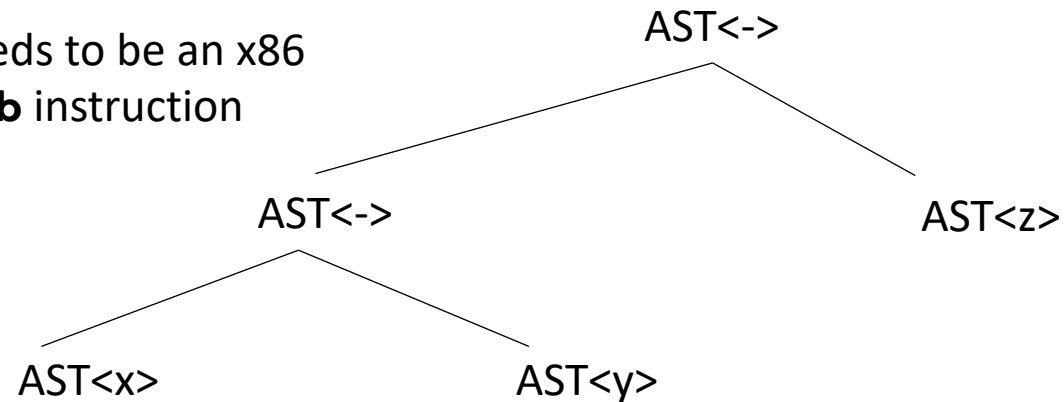
Evaluate an AST by doing a post order traversal

Expr	::=	NUM	Expr2
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			""

```
int x;  
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```

needs to be an x86
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needs to be an x86
sub instruction



Lets do some experiments.

What should 5 - 5.0 be?

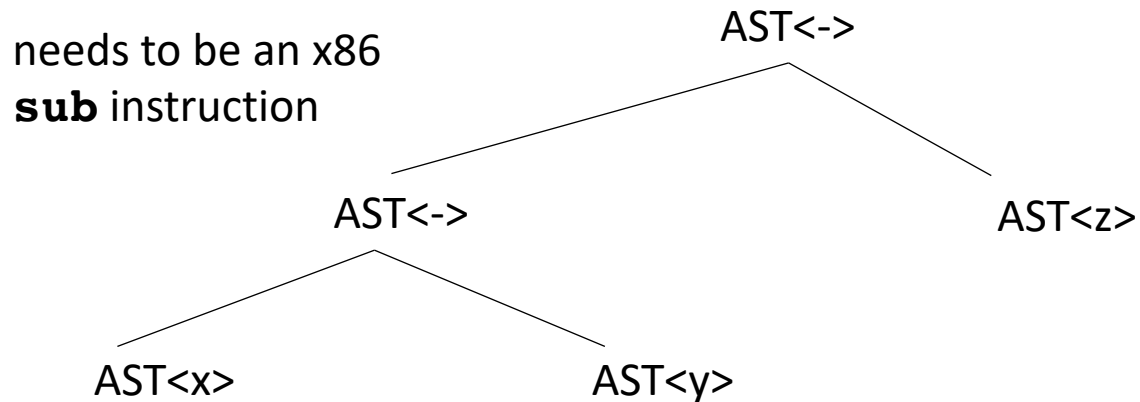
Is this all?

Evaluate an AST by doing a post order traversal

Expr	::=	NUM	Expr2
Expr2	::=	MINUS	NUM Expr2
			""

```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

needs to be an x86
subss instruction



Is this all?

Lets do some experiments.

What should 5 - 5.0 be?

but

subss r1 r2

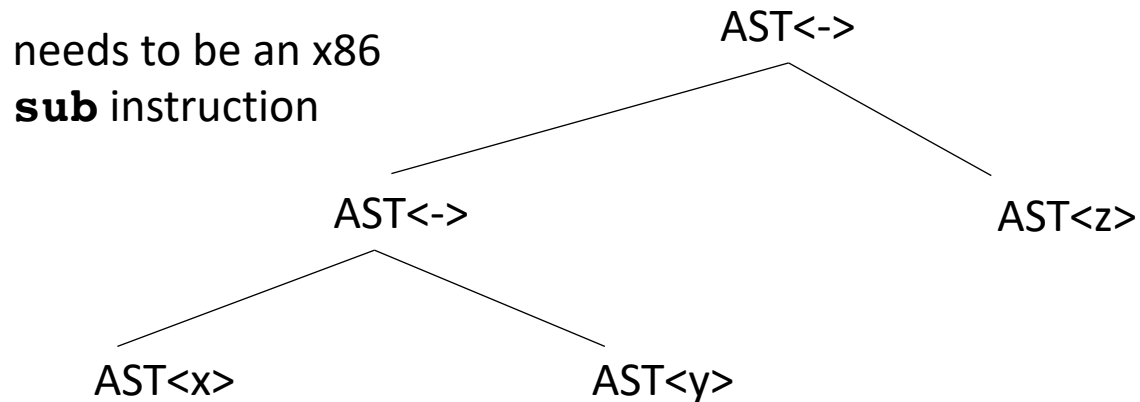
interprets both registers
as floats

Evaluate an AST by doing a post order traversal

```
Expr ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```

```
int x;
int y;
float z;
float w;
w = x - y - z
```

needs to be an x86
subss instruction



But the binary of 5 is 0b101
the float value of 0b101 is 7.00649232162e-45

We cannot just subtract them!

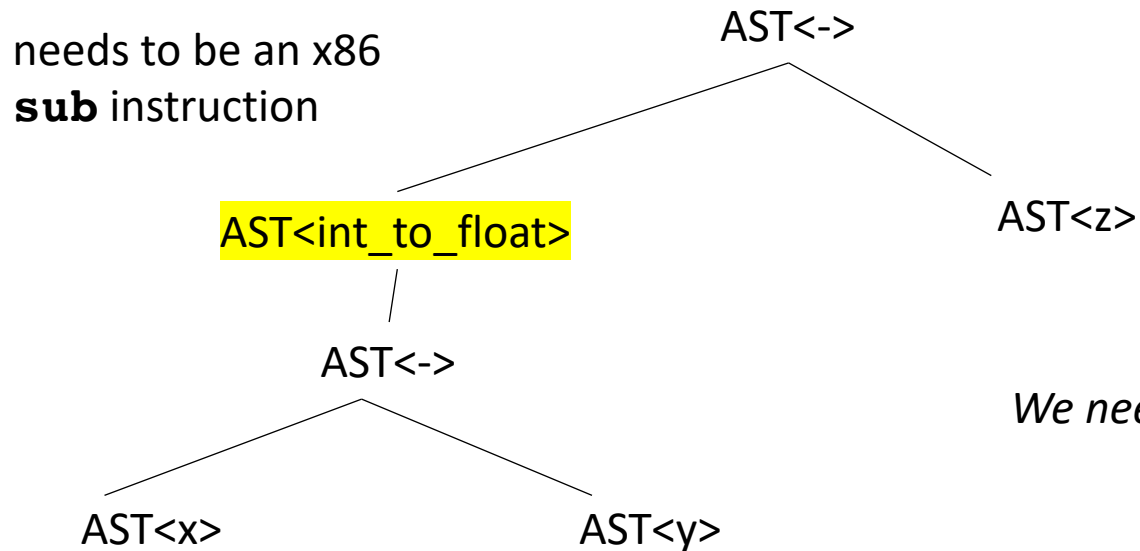
Is this all?

Evaluate an AST by doing a post order traversal

```
Expr ::= NUM Expr2
Expr2 ::= MINUS NUM Expr2
      | ""
```

```
int x;
int y;
float z;
float w;
w = x - y - z
```

needs to be an x86
subss instruction



We need to make sure our operands are in the right format!

Type systems

- Given a language a type system defines:
 - The primitive (base) types in the language
 - How the types can be converted to other types
 - implicitly or explicitly
 - How the user can define new types

Type checking

- Check a program to ensure that it adheres to the type system

Especially interesting for compilers as a program given in the type system for the input language must be translated to a type system for lower-level program

Type systems

- Different types of Type Systems for languages:
 - **statically typed**: types can be determined at compile time
 - **dynamically typed**: types are determined at runtime
 - **untyped**: the language has no types
- What are examples of each?
- What are pros and cons of each?

Type systems

- Different types of Type Systems for languages:
 - **statically typed**: types can be determined at compile time
 - **dynamically typed**: types are determined at runtime
 - **untyped**: the language has no types
- What are examples of each?
- What are pros and cons of each?
- In this class, we will be:
 - Compiling a statically typed language (similar to C)
 - into an untyped language (similar to an ISA)
 - using a dynamically typed language (python)

Type systems

Considerations:

- common base types in a language:
 - ints
 - chars
 - strings
 - floats
 - bool
- How to combine types in expressions:
 - int and float?
 - int and char?
 - int and bool?

Type systems

Considerations:

- common base types in a language:

- ints
- chars
- strings
- floats
- bool

size of ints?

How does C do it?

How does Python do it?

Pros and cons?

- How to combine types in expressions:

- int and float?
- int and char?
- int and bool?

Type systems

Considerations:

- common base types in a language:

- ints
- chars
- strings
- floats
- bool

Are strings a base type? In C? In Python?

- How to combine types in expressions:

- int and float?
- int and char?
- int and bool?

Type systems

Considerations:

- common base types in a language:

- ints
- chars
- strings
- floats
- **bool**

How are bools handled? in C? in Python

- How to combine types in expressions:

- int and float?
- int and char?
- int and bool?

Type systems

Considerations:

- common base types in a language:
 - ints
 - chars
 - strings
 - floats
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Type systems

Considerations:

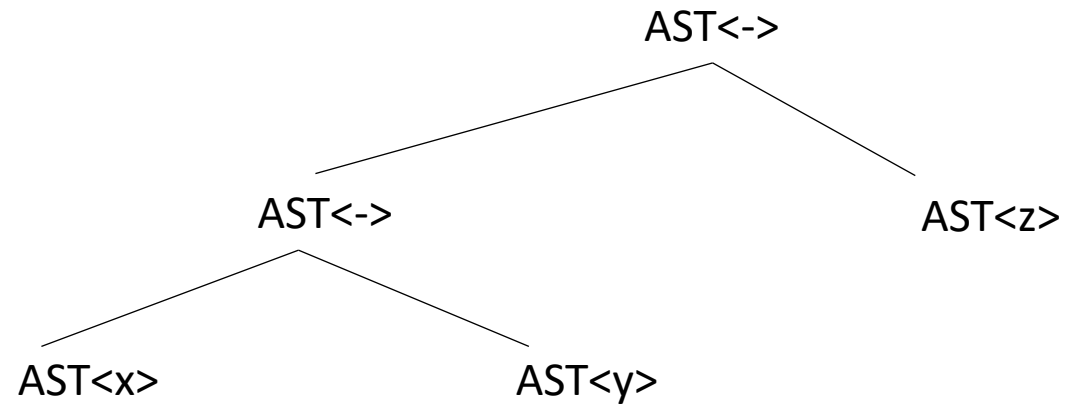
- common base types in a language:
 - ints
 - chars
 - strings
 - floats
 - bool
- How to combine types in expressions:
 - int and float?
 - int and char?
 - int and bool?

What do each of these do if they are +’ed together?

Type checking on an AST

```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

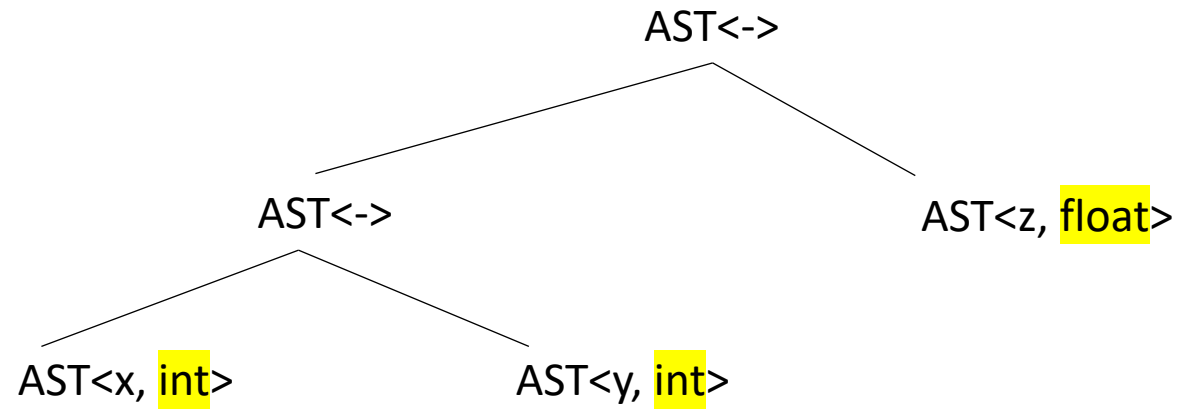
each node additionally gets a type



Type checking on an AST

```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

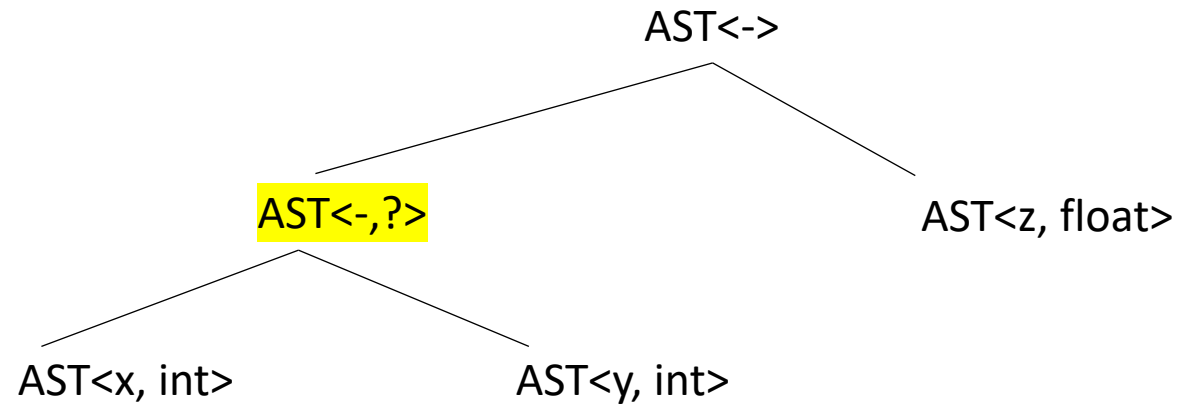
*each node additionally gets a type
we can get this from the symbol table for the leaves*



Type checking on an AST

```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

How do we get the type for this one?



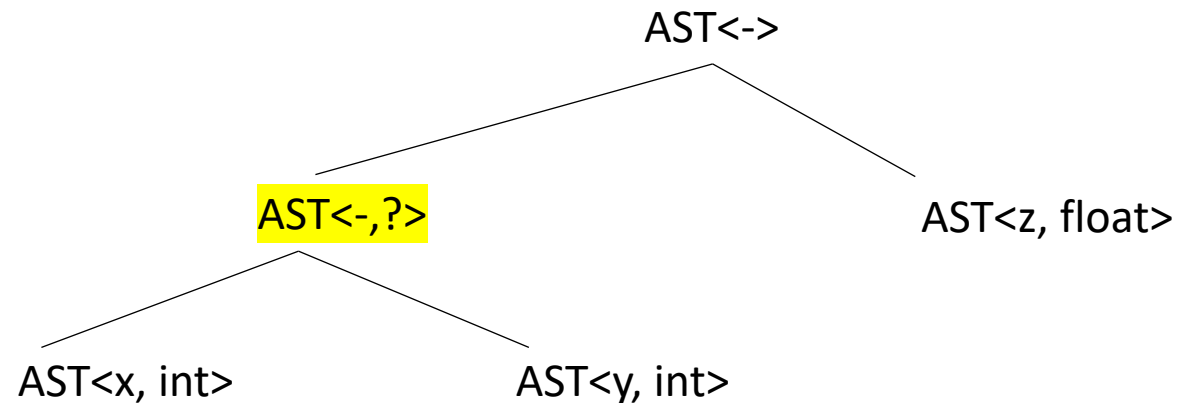
Type checking on an AST

```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

How do we get the type for this one?

combination rules for subtraction:

first	second	result
int	int	int
int	float	float
float	int	float
float	float	float



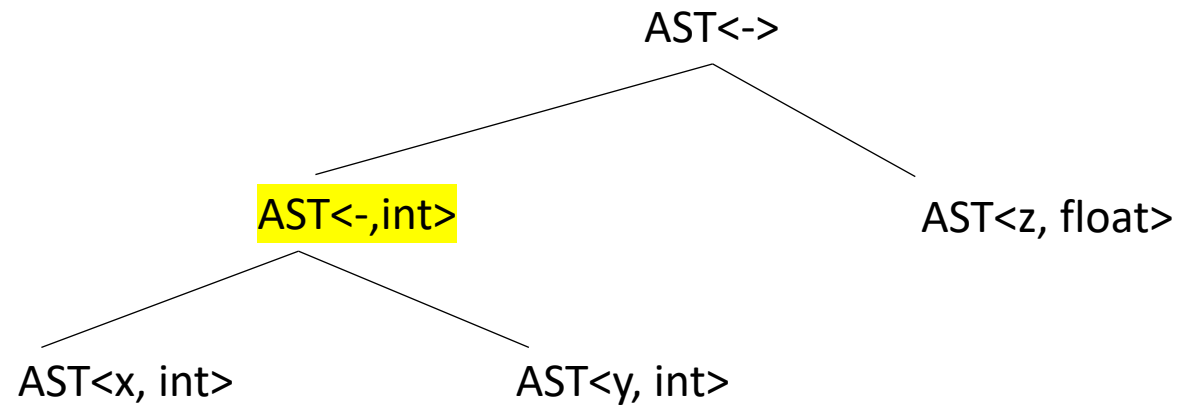
Type checking on an AST

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int x;  
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How do we get the type for this one?

inference rules for subtraction:

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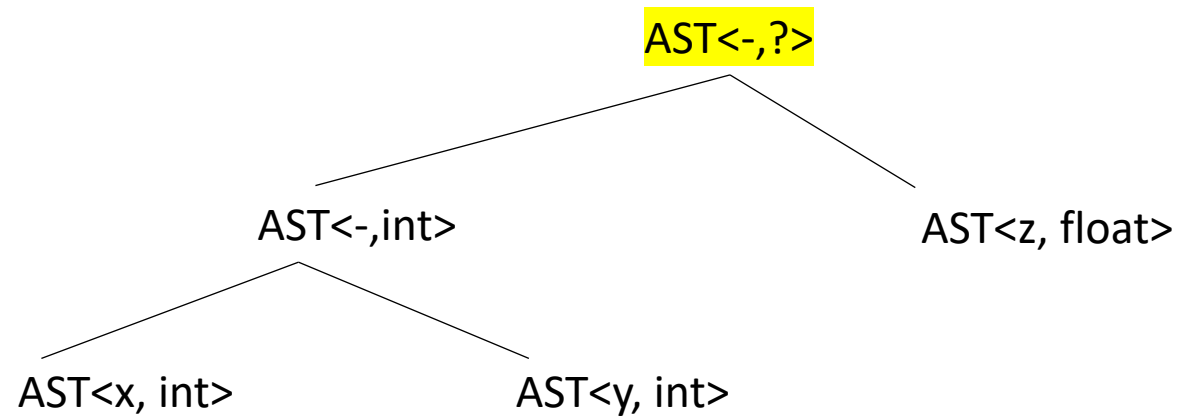
Type checking on an AST

```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
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How do we get the type for this one?

inference rules for subtraction:

first	second	result
int	int	int
int	float	float
float	int	float
float	float	float



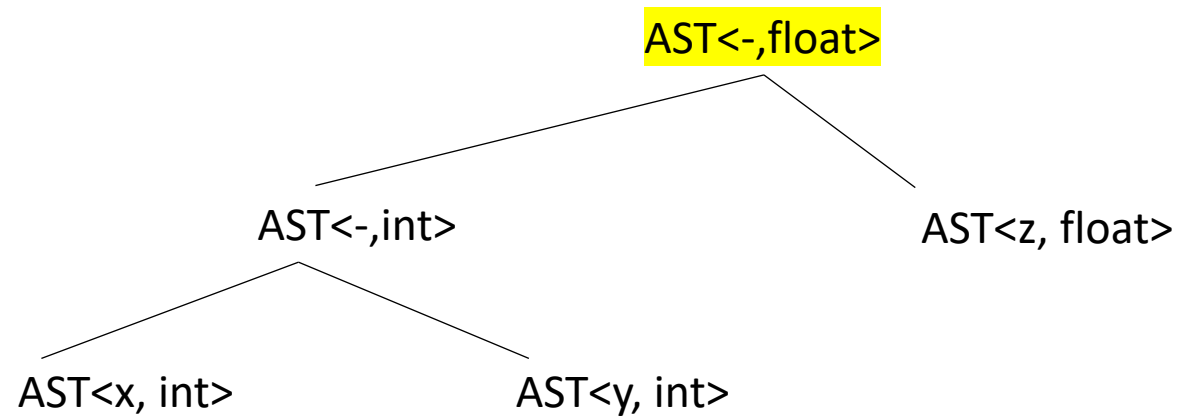
Type checking on an AST

```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

How do we get the type for this one?

inference rules for subtraction:

first	second	result
int	int	int
int	float	float
float	int	float
float	float	float



Type checking on an AST

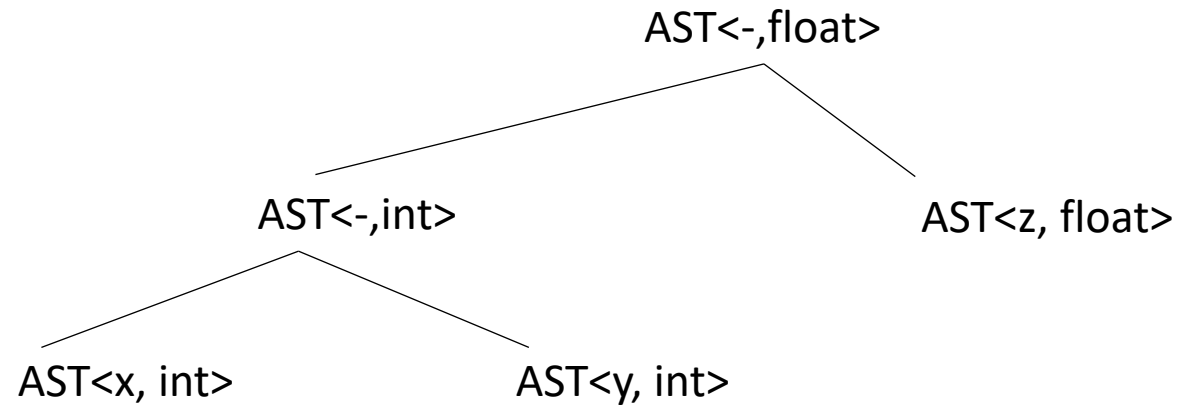
```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

How do we get the type for this one?

inference rules for subtraction:

first	second	result
int	int	int
int	float	float
float	int	float
float	float	float

what else?



Type checking on an AST

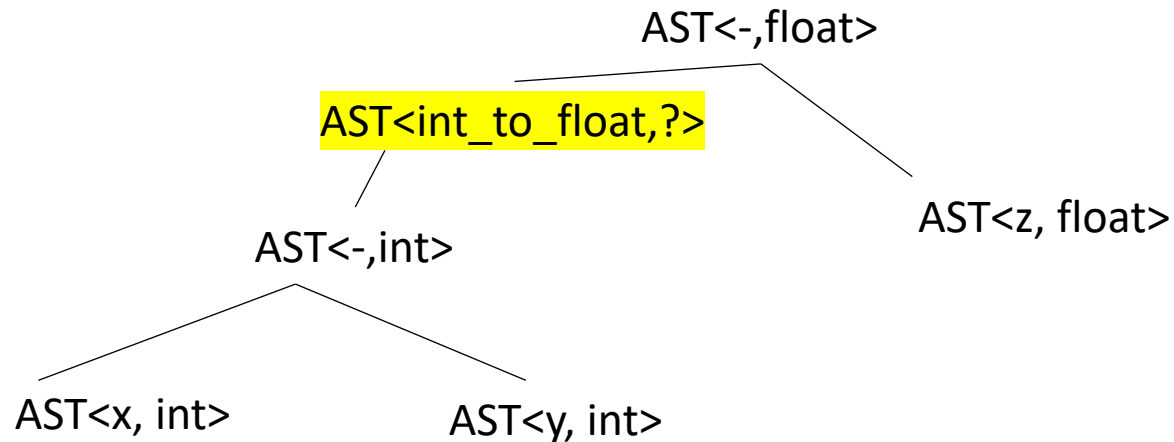
```
int x;  
int y;  
float z;  
float w;  
w = x - y - z
```

How do we get the type for this one?

inference rules for subtraction:

first	second	result
int	int	int
int	float	float
float	int	float
float	float	float

what else? need to convert the int to a float



```
class ASTNode():
    def __init__(self):
        pass
```

```
class ASTLeafNode(ASTNode):
    def __init__(self, value):
        self.value = value

class ASTNumNode(ASTLeafNode):
    def __init__(self, value):
        super().__init__(value)

class ASTIDNode(ASTLeafNode):
    def __init__(self, value):
        super().__init__(value)
```

```
class ASTBinOpNode(ASTNode):
    def __init__(self, l_child, r_child):
        self.l_child = l_child
        self.r_child = r_child

class ASTPlusNode(ASTBinOpNode):
    def __init__(self, l_child, r_child):
        super().__init__(l_child, r_child)

class ASTMultNode(ASTBinOpNode):
    def __init__(self, l_child, r_child):
        super().__init__(l_child, r_child)
```

Enum for types

```
from enum import Enum

class Types(Enum):
    INT = 1
    FLOAT = 2
```

Now we need to set the types for the leaf nodes

Our base AST Node needs a type

```
class ASTNode():
    def __init__(self):
        self.node_type = None
        pass

    def set_type(self, t):
        self.node_type = t

    def get_type(self):
        return self.node_type
```

Enum for types

```
from enum import Enum

class Types(Enum):
    INT = 1
    FLOAT = 2
```

Our base AST Node needs a type

```
class ASTNode():
    def __init__(self):
        self.node_type = None
        pass

    def set_type(self, t):
        self.node_type = t

    def get_type(self):
        return self.node_type
```

Now we need to set the types for the leaf nodes

```
class ASTNumNode(ASTLeafNode):
    def __init__(self, value):
        super().__init__(value)
        if is_int(value):
            self.set_type(Types.INT)
        else:
            self.set_type(Types.FLOAT)
```

Enum for types

```
from enum import Enum

class Types(Enum):
    INT = 1
    FLOAT = 2
```

Our base AST Node needs a type

```
class ASTNode():
    def __init__(self):
        self.node_type = None
        pass

    def set_type(self, t):
        self.node_type = t

    def get_type(self):
        return self.node_type
```

Now we need to set the types for the leaf nodes

```
class ASTNumNode(ASTLeafNode):
    def __init__(self, value):
        super().__init__(value)
        if is_int(value):
            self.set_type(Types.INT)
        else:
            self.set_type(Types.FLOAT)
```

```
class ASTIDNode(ASTLeafNode):
    def __init__(self, value, value_type):
        super().__init__(value)
        self.set_type(value_type)
```

Where can we get the value type for an ID?

Symbol Table

Say we are matched the statement:
`int x;`

- `SymbolTable ST;`

```

                                (TYPE, 'int') (ID, 'x')
declare_statement ::= TYPE ID SEMI
{
    eat(TYPE)
    id_name = self.to_match.value
    eat(ID)
    ST.insert(id_name, None)
    eat(SEMI)
}
```

*in previous lectures we didn't
record any information in the symbol
table*

Symbol Table

Say we are matched the statement:
`int x;`

- SymbolTable ST;

(TYPE, 'int') (ID, 'x')
declare_statement ::= TYPE ID SEMI

{

`value_type = self.to_match.value`

`eat(TYPE)`

`id_name = self.to_match.value`

`eat(ID)`

`ST.insert(id_name, value_type)`

`eat(SEMI)`

}

*in previous lectures we didn't
record any information in the symbol
table*

record the type in the symbol table

Enum for types

```
from enum import Enum

class Types(Enum):
    INT = 1
    FLOAT = 2
```

Our base AST Node needs a type

```
class ASTNode():
    def __init__(self):
        self.node_type = None
        pass

    def set_type(self, t):
        self.node_type = t

    def get_type(self):
        return self.node_type
```

Now we need to set the types for the leaf nodes

```
class ASTNumNode(ASTLeafNode):
    def __init__(self, value):
        super().__init__(value)
        if is_int(value):
            self.set_type(Types.INT)
        else:
            self.set_type(Types.FLOAT)
```

```
class ASTIDNode(ASTLeafNode):
    def __init__(self, value, value_type):
        super().__init__(value)
        self.set_type(value_type)
```

Where can we get the value type for an ID?

But that doesn't get us here...

add the type at parse time

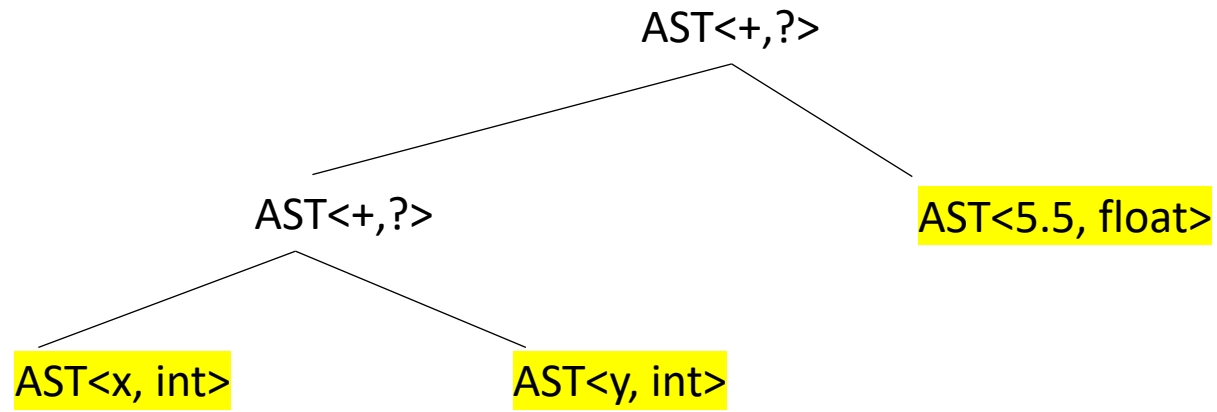
Unit ::= ID
NUM

```
def parse_unit(self, lhs_node):  
    # ... for applying the first production rule (ID)  
    value = self.next_word.value  
    # ... Check that value is in the symbol table  
    node = ASTIDNode(value, ST[value])  
    return node
```

Type inference

- We now have the types for the leaf nodes

```
int x;  
int y;  
float w;  
w = x + y + 5.5
```

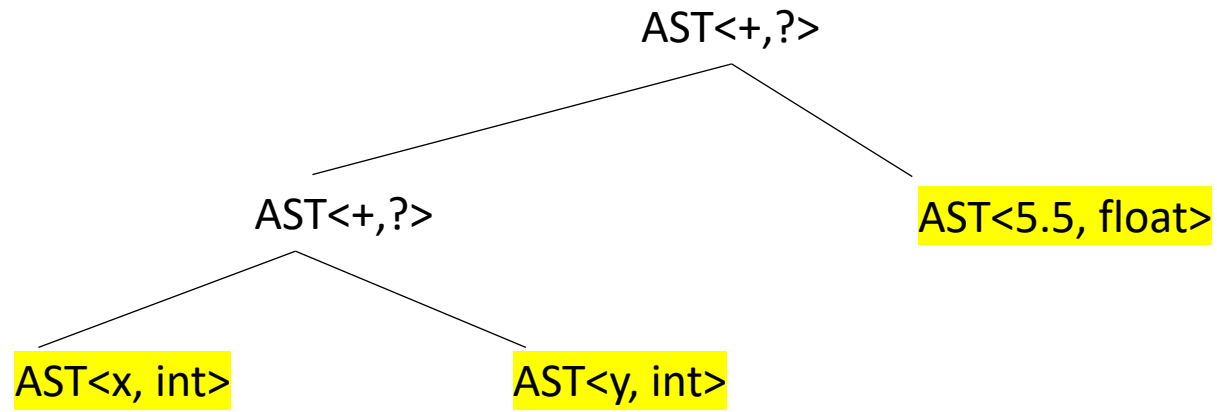


Type inference

- We now have the types for the leaf nodes

Next steps:

we do a post order traversal
on the AST and do a type inference



Type inference

def **type_inference**(n):

Given a node n: find its type and the types of any of its children

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

base case

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

```
        if n is a plus node:
            ...
```


Type inference

def type_inference(n): Given a node n: find its type and the types of any of its children

 case split on n:

 if n is a leaf node:
 return n.get_type()

 if n is a plus node: *lookup the rule for plus*
 return lookup type from table

inference rules for plus

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

def type_inference(n): Given a node n: find its type and the types of any of its children

 case split on n:

 if n is a leaf node:
 return n.get_type()

 if n is a plus node: *lookup the rule for plus*
 return lookup type from table

inference rules for plus

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

but we're missing a few things

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

*we need to make sure the
children have types!*

```
        if n is a plus node:
            do type inference on children
            return lookup type from table
```

inference rules for plus

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

we should record our type

```
        if n is a plus node:
            do type inference on children
            t = lookup type from table
            set n type to t
            return t
```

inference rules for plus

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

def type_inference(n): Given a node n: find its type and the types of any of its children

 case split on n:

 if n is a leaf node:
 return n.get_type()

 if n is a **plus node**:
 do type inference on children
 t = lookup type from table
 set n type to t
 return t

is this just for plus?

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

is this just for plus?

most language promote types, e.g. ints to float for expression operators

```
        if n is a plus node:
            do type inference on children
            t = lookup type from table
            set n type to t
            return t
```

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

is this just for plus?

most language promote types, e.g. ints to float for expression operators

```
        if n is a bin op node:
            do type inference on children
            t = lookup type from table
            set n type to t
            return t
```

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

```
def type_inference(n):  
  
    case split on n:  
  
        if n is a leaf node:  
            return n.get_type()  
  
        if n is a bin op node:  
            do type inference on children  
            t = lookup type from table  
            set n type to t  
            return t
```

What about for assignments?

```
int x;  
cout << (x = 5.5) << endl;
```

What does this return?

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

```
def type_inference(n):  
  
    case split on n:  
  
        if n is a leaf node:  
            return n.get_type()  
  
        if n is a bin op node:  
            do type inference on children  
            t = lookup type from table  
            set n type to t  
            return t
```

What about for assignments?

```
int x;  
cout << (x = 5.5) << endl;
```

What does this return?

left	right	result
int	int	int
int	float	int
float	int	float
float	float	float

whatever the left is

Type inference

```
def type_inference(n):
```

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

```
        if n is an assignment:
            ....
```

```
        if n is a bin op node:
            ...
```

What about for assignments?

```
int x;
cout << (x = 5.5) << endl;
```

What does this return?

left	right	result
int	int	int
int	float	int
float	int	float
float	float	float

whatever the left is

Type checking

- Checking for errors

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

we should record our type

```
        if n is a plus node:
            do type inference on children
            t = lookup type from table
            if t is None:
                throw type exception
            set n type to t
            return t
```

inference rules for plus

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float

Type inference

```
def type_inference(n):
```

Given a node n: find its type and the types of any of its children

```
    case split on n:
```

```
        if n is a leaf node:
            return n.get_type()
```

we should record our type

```
        if n is a plus node:
            do type inference on children
            t = lookup type from table
            if t is None:
                throw type exception
            set n type to t
            return t
```

inference rules for plus

left	right	result
int	int	int
int	float	float
float	int	float
float	float	float
string	int	None

like in Python

See everyone on Friday!

- We will discuss implementing type inference on Monday